

## Algorithm Design Techniques

- **Induction**
- Divide and Conquer
- Dynamic Programming
- Greedy
- Reduction

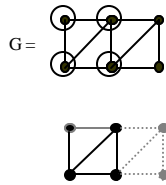
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## Induced Subgraphs

- Let  $G=(V,E)$  be a graph and let  $W$  be a subset of  $V$ .
- The subgraph of  $G$  induced by  $W$  is the graph has  
Vertex set:  $W$   
Edge set:  $\{(x,y) \in E \text{ such that } x \text{ and } y \text{ are in } W\}$

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## Example



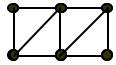
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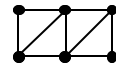
## Maximal Induced Subgraph

- Input: A graph  $G=(V,E)$  and an integer  $k$ .
- Output: A largest subgraph  $G'=(V',E')$  such that every vertex of  $G'$  has degree at least  $k$ .

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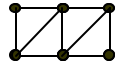
## Maximal Induced Subgraph Example

- Input:  $G=$    $K=2$

- Output: 

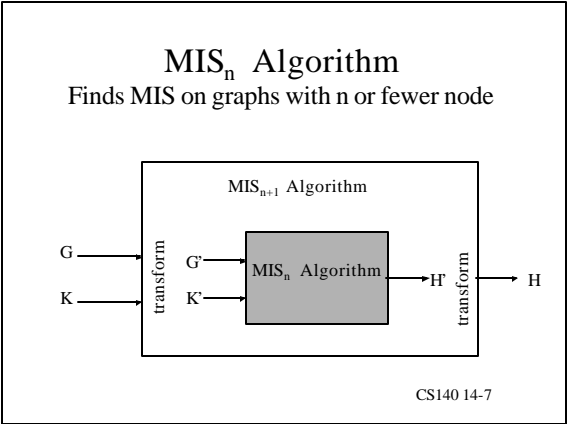
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## Maximal Induced Subgraph Example

- Input:  $G=$    $K=3$

- Output:  $\phi$

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### MIS<sub>n</sub> Algorithm

- Suppose  $x$  is a vertex of  $G$  with degree  $< k$
- If  $H$  is a solution in  $G$  then  $H$  exists in  $G - \{x\}$

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### Evaluating Polynomials

- Input: Set of integers  $\{a_0, a_1, \dots, a_{n-1}\}$  and an integer  $x$ .
- Output:  $P_n(x) = \sum_{i=0}^{n-1} a_i x^i$

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### Evaluating Polynomials

- Input: List of integers  $a_0, a_1, \dots, a_{n-1}$  and an integer  $x$ .
- Output:  $P_n(x) = \sum_{i=0}^{n-1} a_i x^i$

Example:

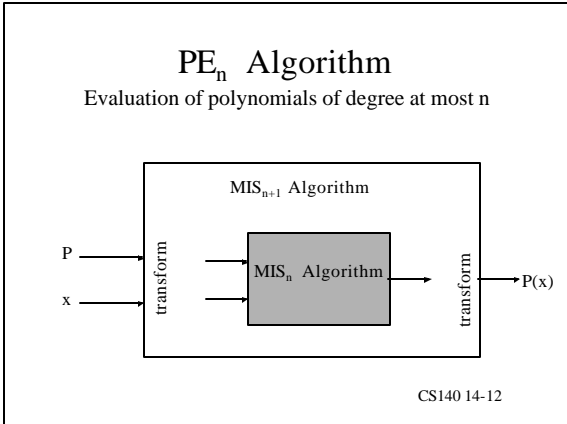
Input: 3,-1,0,2 and 2  
Output: 22  
Explanation:  $P(x)=3x^3-x^2+2$ ;  $P(2)=22$

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### Naïve Approach

- Number of multiplications:
- Number of additions:

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### $PE_n$ Algorithm

- $P(x) = \sum_{i=0, \dots, n} a_i x^i$   
 $= a_0 + x (\sum_{i=1, \dots, n} a_i x^{i-1})$

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### Inductive Approach

- Number of multiplications:
- Number of additions:

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