

Matrix Chain Multiplication

- A is an $n \times m$ matrix
- B is an $m \times k$ matrix
- How many scalar multiplications are needed to compute AB?

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Matrix Chain Multiplication

- A is a 2×5 matrix
- B is a 5×1000 matrix
- C is a 1000×2 matrix.
- How many scalar multiplications are needed to compute ABC?

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Matrix Chain Multiplication

- Input: A list of $n+1$ integers p_1, p_2, \dots, p_{n+1}
- Output: The minimum number of scalar multiplications needed to compute $\prod_{i=1..n} A_i$ where A_i is a $p_i \times p_{i+1}$ matrix.
- Assume a standard matrix multiplication procedure is used; i.e. no Strassen-like improvements.

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Inductive Approach cont.

- Consider an input: $p_1, p_2, p_3, p_4, p_5, p_6$
- Imagine an optimal solution: 10773

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Inductive Approach cont.

- Consider an input: $p_1, p_2, p_3, p_4, p_5, p_6$
- Imagine an optimal way of multiplying matrices A_1, A_2, A_3, A_4, A_5 :

$$(A_1(A_2A_3)) (A_4A_5)$$

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Inductive Approach cont.

- It prescribes some last multiplication

$$(A_1(A_2A_3)) | (A_4A_5)$$

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Inductive Approach cont.

- It prescribes some last multiplication
 $(A_1(A_2A_3)) \mid (A_4A_5)$
- So $OPT(A_1, A_2, A_3, A_4, A_5) =$
 $OPT(A_1, A_2, A_3) + OPT(A_4, A_5) + p_1p_4p_5$

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Inductive Approach cont.

- We don't know where the top split occurs ... but clearly $OPT(A_1, A_2, A_3, A_4, A_5) =$

$$\min_{0 < k < 5} OPT(A_1, \dots, A_k) + OPT(A_{k+1}, \dots, A_5) + p_1p_{k+1}p_5$$

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Recursive Algorithm Running Time

- $T(n) = \sum_{0 < k < n} T(k) + T(n-k) + c$
- How bad is it?

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Dynamic Programming

- Use a table to store results
- What kind of results?
- $M(k, j) =$ Minimum number of multiplications to compute $\prod_{i=k \dots j} A_i$

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$M(k, j)$ for $k \leq j$

	1	2	3	4	5
1					
2					
3					
4					
5					

$M(3, 5)$ needs:
 $M(3, 3), M(4, 5),$
 $M(3, 4), M(5, 5)$

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$M(k, j)$ needs $M(i, m)$ where $m - i < j - k$

	1	2			
1					
2					
3					
4					
5					

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$M(k,j)$ needs $M(i,m)$
 where $m-i < j-k$

0				
	0			
		0		
			0	
				0

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Dynamic Programming Algorithm

- $M(k,k)=0$
- For j,k such that $j-k = 1, 2, \dots, n-1$
 $M(k,j) = \min_{i=k \dots j-1} M(k,i) + M(i+1,j) + p_k p_{i+1} p_j$
- Return $M(1,n)$

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Input: 2,3,1,5,4,8
 (A_1 is 2×3 , A_2 is 3×1 , ...)

0	6			
	0	15		
		0	20	
			0	160
				0

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MCM Algorithm

- Recursive Algorithm takes exponential time.
- Dynamic Programming takes _____.

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