

## Graph Algorithms

- Graph traversal
- Topological sort
- Strongly connected components
- Single-source shortest path

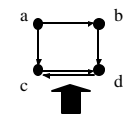
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## Graph and Digraph Traversal

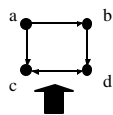
- Breadth-first
- Depth-first

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## Digraph Notation



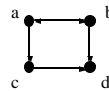
This is hard to draw.



So we'll use this to mean  $\langle c,d \rangle$  and  $\langle d,c \rangle$  are both edges of the digraph.

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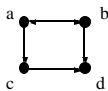
## Digraph Representation Adjacency Matrix



	a	b	c	d
a	0	1	1	0
b	1	0	0	1
c	0	0	0	1
d	0	0	0	0

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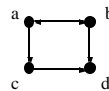
## Digraph Representation Adjacency Matrix



	in-edges			
out-edges	a	b	c	d
a	0	1	1	0
b	1	0	0	1
c	0	0	0	1
d	0	0	0	0

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## Digraph Representation Adjacency List

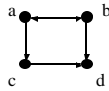


Vertex	Out-edges	In-edges (optional)
a	b→c	
b	a→d	a
c	d	a
d		b→c

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### Breadth-First(c)

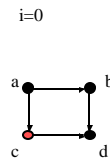
- For  $i=0,1,\dots$   
Visit each unvisited vertex that is distance  $i$  from  $c$



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### Breadth-First(c)

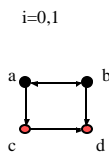
- For  $i=0,1,\dots$   
Visit each unvisited vertex that is distance  $i$  from  $c$



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### Breadth-First(c)

- For  $i=0,1,\dots$   
Visit each unvisited vertex that is distance  $i$  from  $c$



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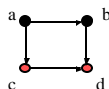
### Digraph notions

- Vertex  $y$  is *reachable* from  $x$  if there is a directed path in  $G$  from  $x$  to  $y$ .
- Note that  $x$  is reachable from  $x$  by a directed path of length 0.

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### Breadth-First(x)

- Breadth-First( $x$ ) visits all vertices *reachable* from  $x$ .



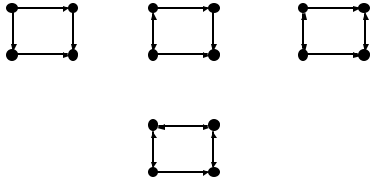
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### Digraph notions cont.

- Vertex  $y$  is *reachable* from  $x$  if there is a directed path in  $G$  from  $x$  to  $y$ .
- Vertices  $x$  and  $y$  are *strongly connected* if  $x$  is reachable from  $y$  and  $y$  is reachable from  $x$ .

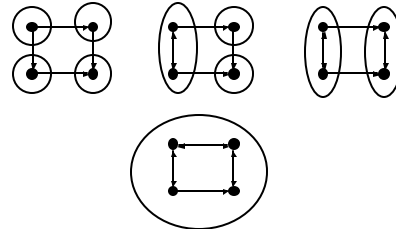
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### Strongly-connected vertices?



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### Strongly-connected vertices:



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### Digraph notions cont.

- Strongly-connected is an equivalence relation:
  - $x$  is strongly connect to  $x$ .
  - If  $x$  is strongly connected to  $y$  then  $y$  is strongly connected to  $x$ .
  - If  $x$  is strongly connected to  $y$  and  $y$  is strongly connected to  $z$  then  $x$  is strongly connected to  $z$ .
- The vertices are partitioned into equivalence classes.

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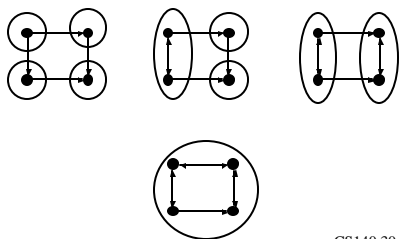
### Strongly Connected Component

A strongly connected components of  $G$  is a *maximal* subgraph of  $G$  in which each pair of vertices is strongly connected.

The strongly connected components are exactly the equivalence classes under the strongly connected relation

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### Strongly-connected components:



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### BFS Application

- Identify the strongly connected components of a digraph  $G$ .

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## DFS Application

- Identify the strongly connected components of a digraph G.

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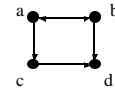
## Depth-First(c)

### Depth-First(x)

Mark x visited

For each edge  $\langle x,y \rangle$

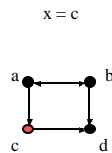
If y is unvisited then  
DFS(y)



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## Depth-First(c)

➡ **Depth-First(x)**  
Mark x visited  
For each edge  $\langle x,y \rangle$   
If y is unvisited then  
DFS(y)



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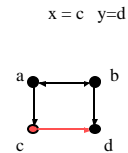
## Depth-First(c)

### Depth-First(x)

Mark x visited

For each edge  $\langle x,y \rangle$

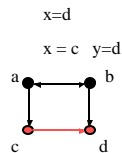
➡ If y is unvisited then  
DFS(y)



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## Depth-First(d)

➡ **Depth-First(x)**  
Mark x visited  
For each edge  $\langle x,y \rangle$   
If y is unvisited then  
DFS(y)



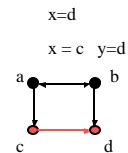
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## Depth-First(d)

### Depth-First(x)

Mark x visited

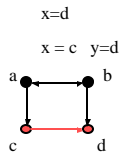
➡ For each edge  $\langle x,y \rangle$   
If y is unvisited then  
DFS(y)



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## Depth-First(d)

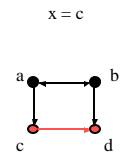
**Depth-First(x)**  
Mark x visited  
For each edge  $\langle x,y \rangle$   
If y is unvisited then  
DFS(y)  
➡ (return)



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## Depth-First(c)

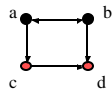
**Depth-First(x)**  
Mark x visited  
➡ For each edge  $\langle x,y \rangle$   
If y is unvisited then  
DFS(y)



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## Depth-First(c)

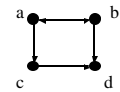
**Depth-First(x)**  
Mark x visited  
For each edge  $\langle x,y \rangle$   
If y is unvisited then  
DFS(y)



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## DFS(G)

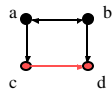
**DFS(G)**  
➡ While G has an  
unvisited vertex x:  
Depth-First(x)  
  
x=c



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## DFS(G)

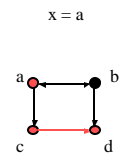
**DFS(G)**  
➡ While G has an  
unvisited vertex x:  
Depth-First(x)  
  
x=a



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## Depth-First(a)

**Depth-First(x)**  
➡ Mark x visited  
For each edge  $\langle x,y \rangle$   
If y is unvisited then  
DFS(y)



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### Depth-First(a)

**Depth-First(x)**  
 Mark x visited  
 → For each edge  $\langle x,y \rangle$   
 If y is unvisited then  
 DFS(y)

$x = a \quad y = b$

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### Depth-First(b)

**Depth-First(x)**  
 Mark x visited  
 → For each edge  $\langle x,y \rangle$   
 If y is unvisited then  
 DFS(y)

$x = b$   
 $x = a \quad y = b$

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### Depth-First(b)

**Depth-First(x)**  
 Mark x visited  
 → For each edge  $\langle x,y \rangle$   
 If y is unvisited then  
 DFS(y)

$x = b \quad y = a$   
 $x = a \quad y = b$

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### Depth-First(b)

**Depth-First(x)**  
 Mark x visited  
 → For each edge  $\langle x,y \rangle$   
 If y is unvisited then  
 DFS(y)

$x = b \quad y = d$   
 $x = a \quad y = b$

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### Depth-First(b)

**Depth-First(x)**  
 Mark x visited  
 For each edge  $\langle x,y \rangle$   
 If y is unvisited then  
 DFS(y)

$x = a \quad y = b$

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### Depth-First(a)

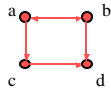
**Depth-First(x)**  
 Mark x visited  
 → For each edge  $\langle x,y \rangle$   
 If y is unvisited then  
 DFS(y)

$x = a \quad y = c$

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## DFS(G)

**DFS(G)**  
➔ While G has an  
unvisited vertex x:  
  Depth-First(x)



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## What is the running time of DFS?

- $O(m+n)$
- Every vertex is pushed onto the stack once and popped from the stack once.
- Each out-edge is inspected once.

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