Outline

- K^{th} order statistics
  - Min
  - Min/Max
- Proof of upper and lower bounds
- Adversary argument

K-th Order Statistics

- Input: Set of integers S
- Output: *k-th* smallest integer in S

What if k=1?
What if k=n?
What if k=n/2?

FIND-MIN

- How many comparisons does it take to find the minimum in a set of integers?
- Answer: n-1

FIND-MIN

In worst case

- How many comparisons does it take to find the minimum in a set of integers?
- Answer: n-1

Upper Bound for FIND-MIN

Upper Bound Theorem: Finding the minimum in a set of n integers requires no more than n-1 comparisons.

Proof: Give algorithm

Lower Bound for FIND-MIN

Lower Bound Theorem: Finding the minimum in a set of integers requires at least n-1 comparisons.

Proof of Lower Bound:

- Consider an algorithm ***A*** on input of size ***n.***
- Let G be a graph with a vertex for each input integer. Initially G has no edges. When ***A*** compares two input values, we'll add an edge between the corresponding vertices of G.
- ***A*** cannot conclude until G has ________ edges.
- Thus ***A*** cannot conclude until it has made ______ comparisons.

Upper Bound for FIND-MIN/MAX

- Upper Bound Theorem: Finding the minimum and maximum in a set of n integers requires no more than $\lceil 3n/2 \rceil$-2 comparisons.
- Proof: Give an algorithm

Proof of Upper Bound:

- Algorithm for even n:
  - Make n/2 pair-wise comparisons
  - Find the maximum of the winners with n/2-1 comparisons
  - Find the minimum of the losers with n/2-1 comparisons
- Algorithm for odd n:
  - Run even algorithm on first n-1 integers
  - Compare the min and max to the last integer

Lower Bound for FIND-MIN/MAX

- Lower Bound Theorem: Finding the minimum and maximum in a set of n integers requires at least $\lceil 3n/2 \rceil$ - 2 comparisons.
- Proof: Adversary argument

Example of an adversary

You pick a number y between 1 and 100

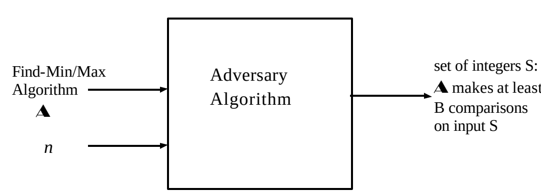
I have to guess y by posing queries of the form "Is it x?"

You answer "yes, x=y" or "no, y<x" or "no, y>x"

- How many queries can you force me to make?
- Prove it!

Adversary Argument to prove bound B

(for FIND MIN/MAX)



FIND-MIN/MAX Adversary - Accounting

- Adversary = interactive comparison oracle
- Accounting scheme: For x in S

$b_{MAX}(x)=$ 1 if the algorithm can rule out x as the largest integer; 0 otherwise

$b_{MIN}(x)=$ 1 if the algorithm can rule out x as the smallest integer; 0 otherwise

FIND-MIN/MAX Adversary- Strategy

- On query "Is x<y ?"
- Answer consistently with previous answers
- If yes and no both consistent then answer so as to minimize the changes in b_{MAX} and b_{MIN} variables

FIND-MIN/MAX Adversary - Analysis

Consider a query "Is x<y?"

- If NO: $b_{MIN}(x)\to 1$ and $b_{MAX}(y)\to 1$
- If YES: $b_{MAX}(x)\to 1$ and $b_{MIN}(y)\to 1$

Proof of Lower Bound:

- Claim: At most $\lfloor n/2 \rfloor$ queries can result in the change of two $b_{MIN/MAX}$ variables
- Claim: 2n-2 changes must occur before the algorithm concludes

$\Rightarrow \lfloor n/2 \rfloor + (2n-2) - 2(\lfloor n/2 \rfloor$ queries are necessary

Double 0's

- Input: n-bit vector of 0/1's
- Question: are there two adjacent 0's
- How many queries are needed?

Exercise

- What is a good adversary strategy?
- What is a good algorithm strategy

Upper Bound

Claim: Double 0's can be solved with f(n) queries where:

$f(n) = n-1$ if $n \equiv 1 \mod 3$
$= n$ otherwise

Lower Bound

Double 0's cannot be solved with fewer than g(n) queries where:

$G(n) = n-1$ if $n \equiv 1 \mod 3$
$= n$ otherwise

Proof

- Do as homework!