

Software

Python



How does Python function ?

Hmmm

4 Hmmm problems
+ 1 loop problem
due **Tues. 3/5**

Assembly Language
Machine Language

CS 5 this week



Fall break?!...

RAM

registers

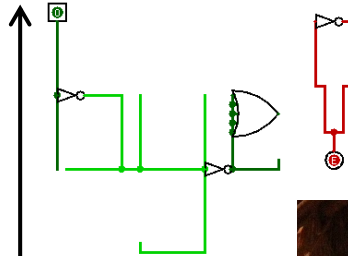
1-bit memory: flip-flops

arithmetic

bitwise functions

logic gates

transistors / switches




Hardware

I have a looming sense...



Fun with circuits?

 **CircuitVerse**

Project ▾ Circuit ▾ Tools ▾ Help ▾

ALU Fun

FA x 4-Bit Adder x mult x div x Complement x 4-Bit Sub x Neg x Add or Sub x +

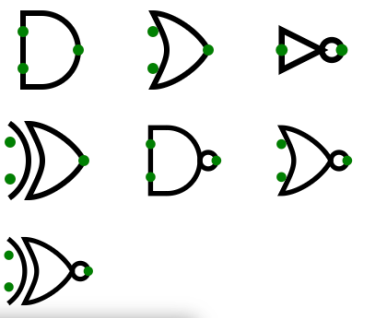
CIRCUIT ELEMENTS

Search..

Input ▾

Output ▾

Gates ▴




TESTBENCH

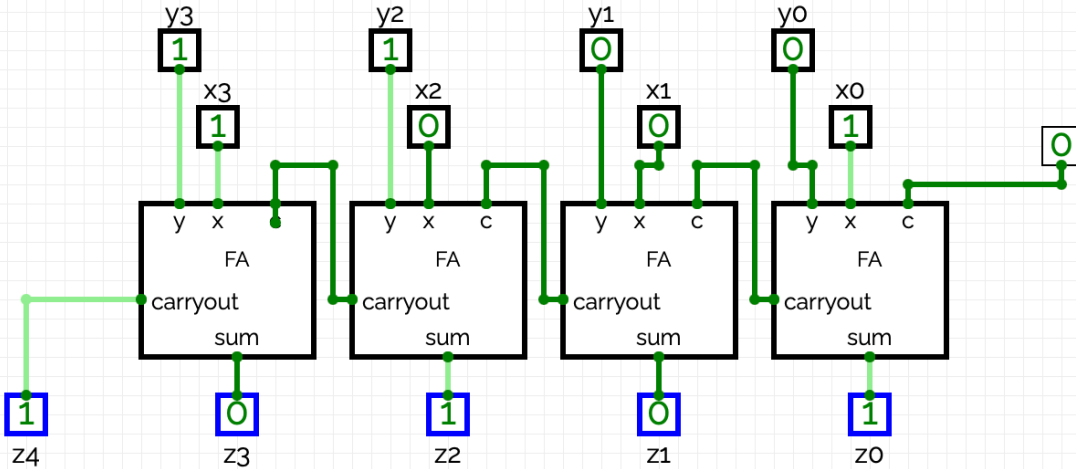
Sequential Elements ▾

Annotation ▾

Misc ▾

TIMING DIAGRAM





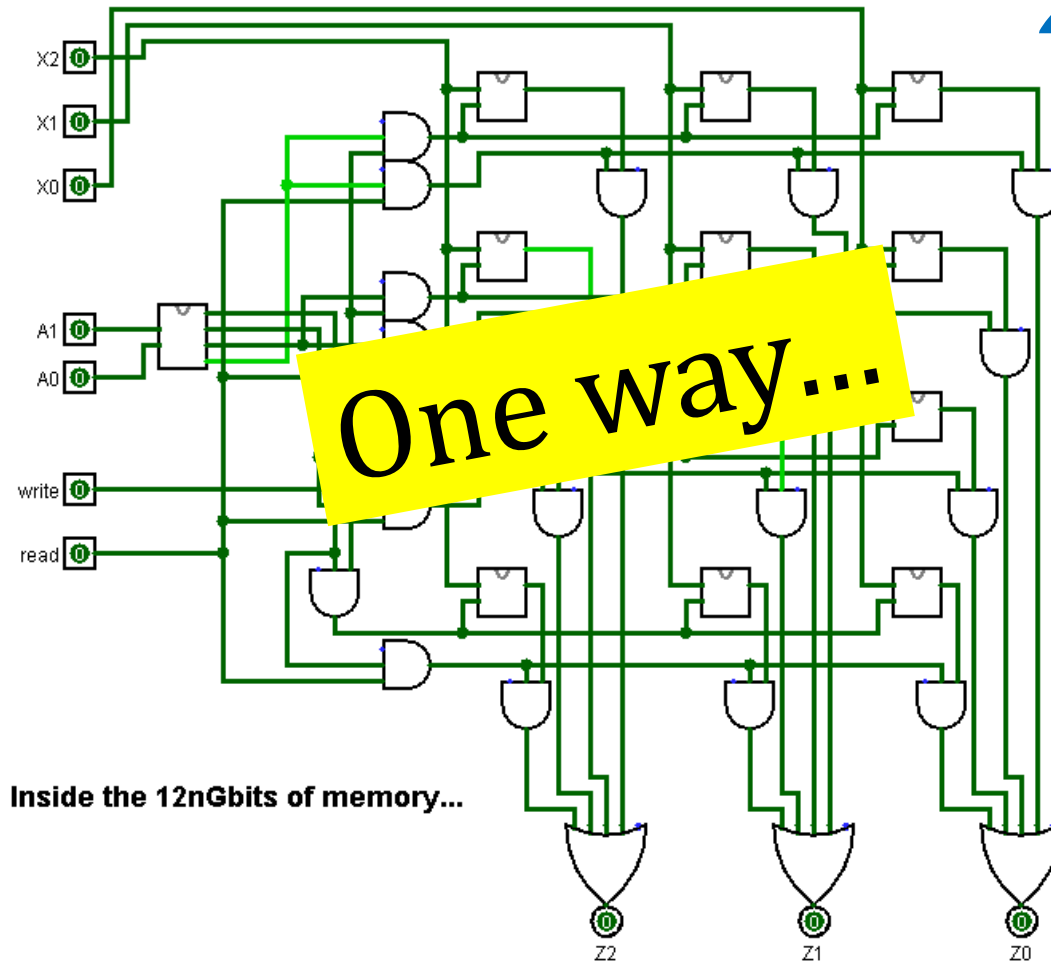
The diagram shows a 4-bit ripple-carry adder circuit. It consists of four Full Adder (FA) blocks connected in series. The carryout of one FA is connected to the carryin of the next. Inputs x_3, x_2, x_1, x_0 are connected to the x inputs of the FAs. Inputs y_3, y_2, y_1, y_0 are connected to the y inputs of the FAs. The carryin of the first FA is connected to input z_4 . The sum outputs are z_3, z_2, z_1, z_0 . The final carryout is connected to input 0.

| Input | Value |
|-------|-------|
| z_4 | 1 |
| z_3 | 0 |
| z_2 | 1 |
| z_1 | 0 |
| z_0 | 1 |

| Output | Value |
|----------------|-------|
| y_3 | 1 |
| x_3 | 1 |
| y_2 | 1 |
| x_2 | 0 |
| y_1 | 0 |
| x_1 | 0 |
| y_0 | 0 |
| x_0 | 1 |
| Final carryout | 0 |

Making memories...

*~1952-
2024*

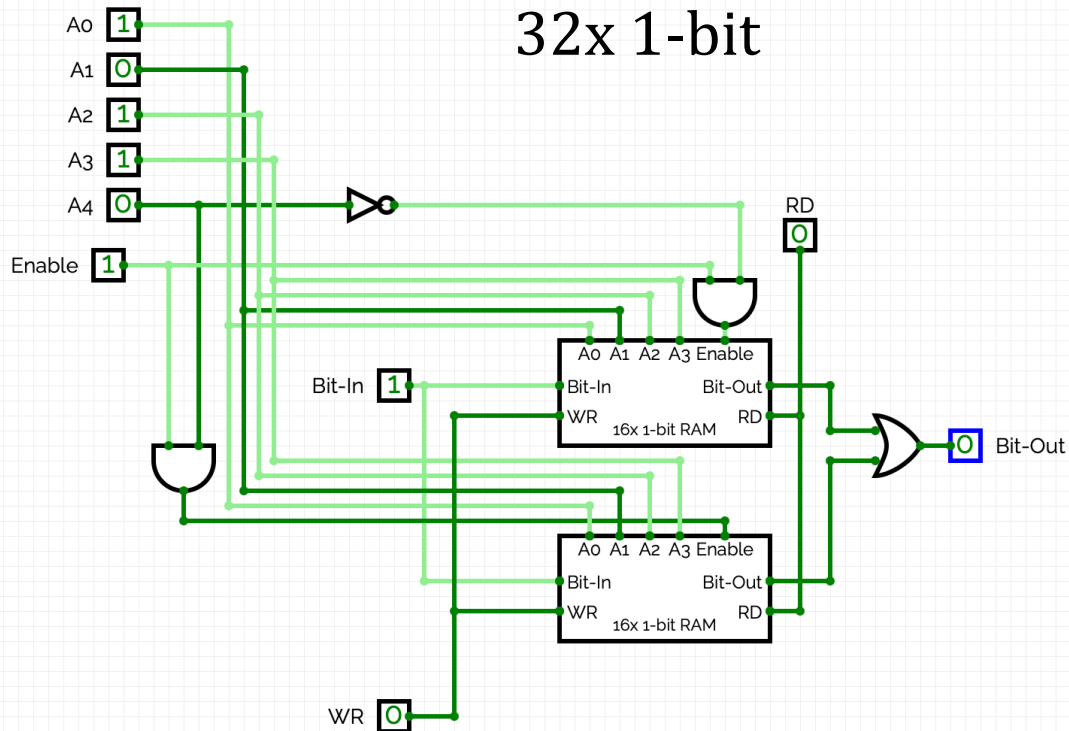


My head is spinning...

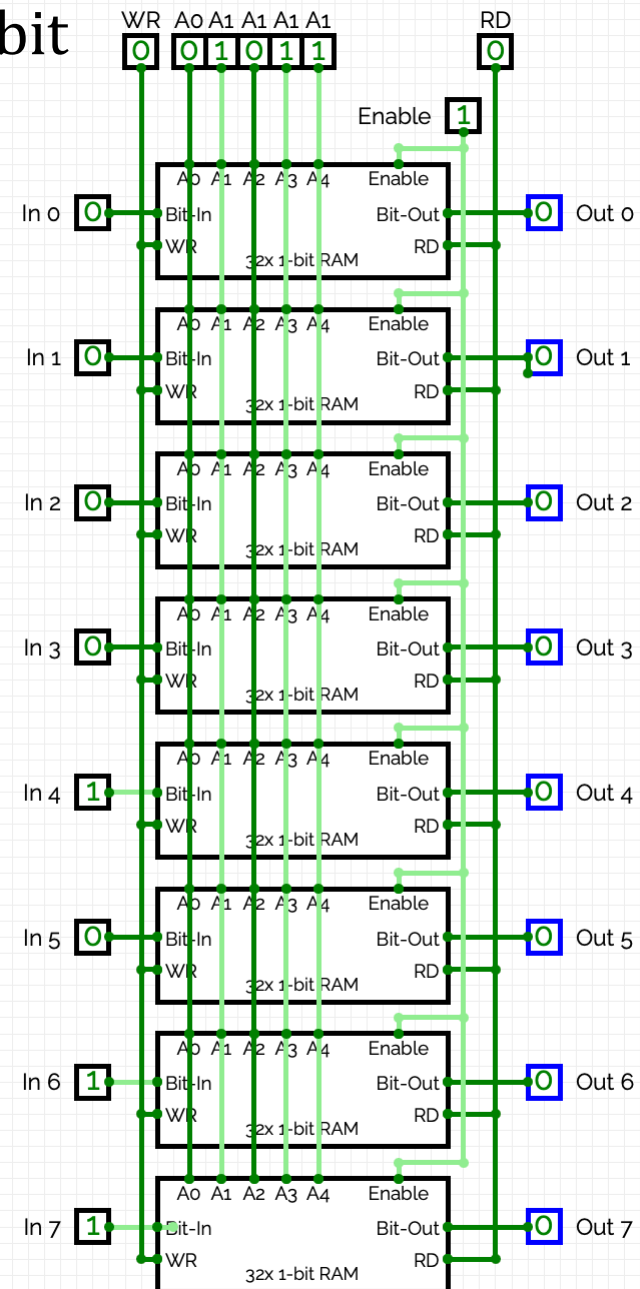
Circuits ~ Memory!

32 bytes of memory

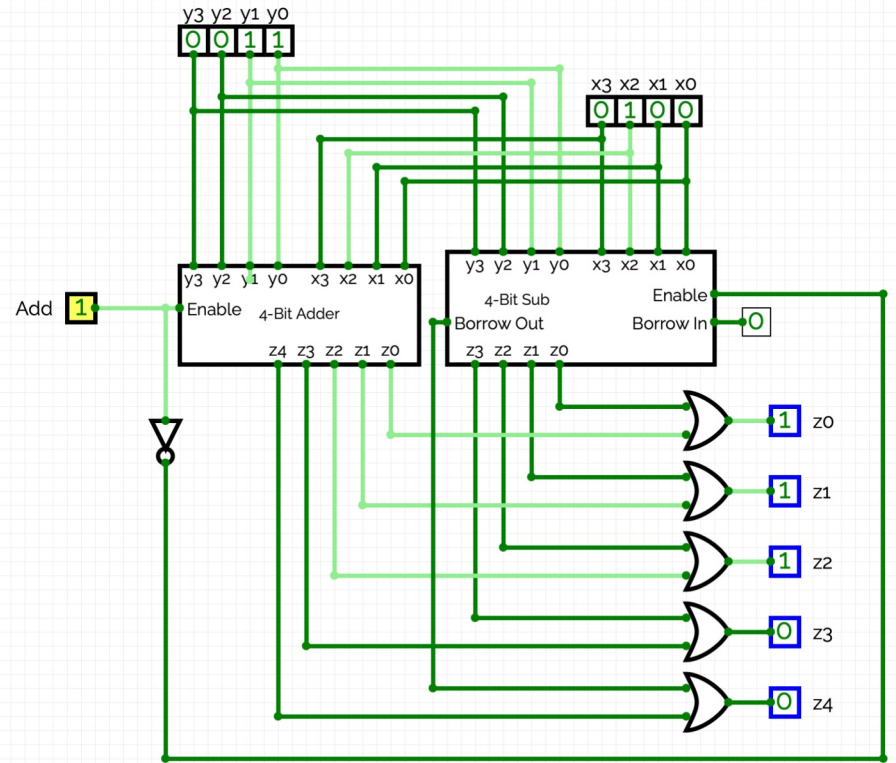
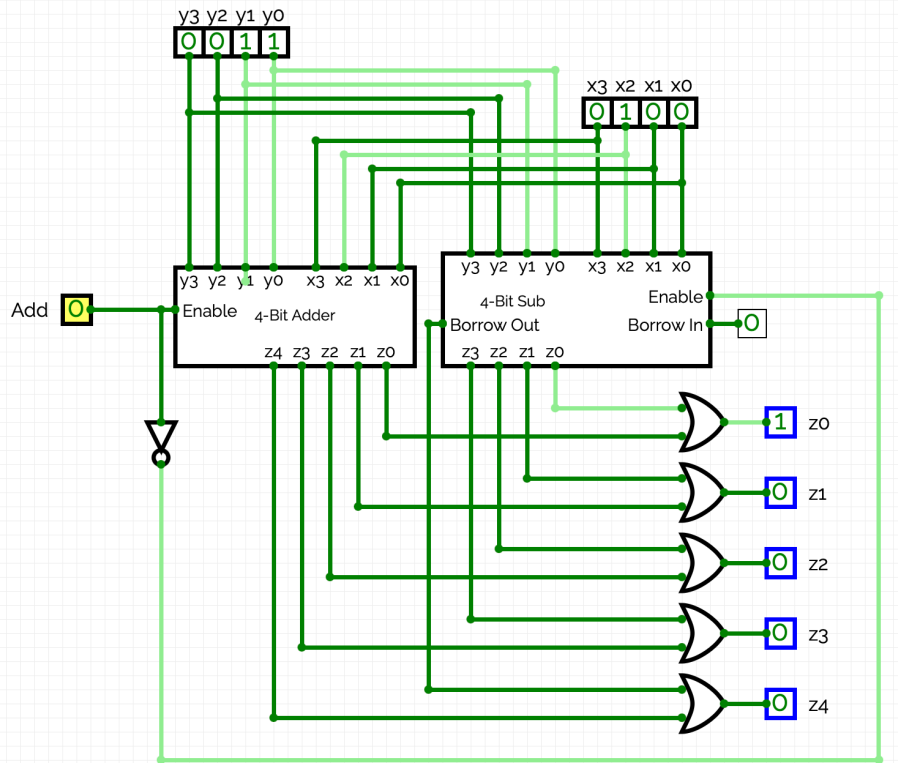
the power of composition



32x 8-bit

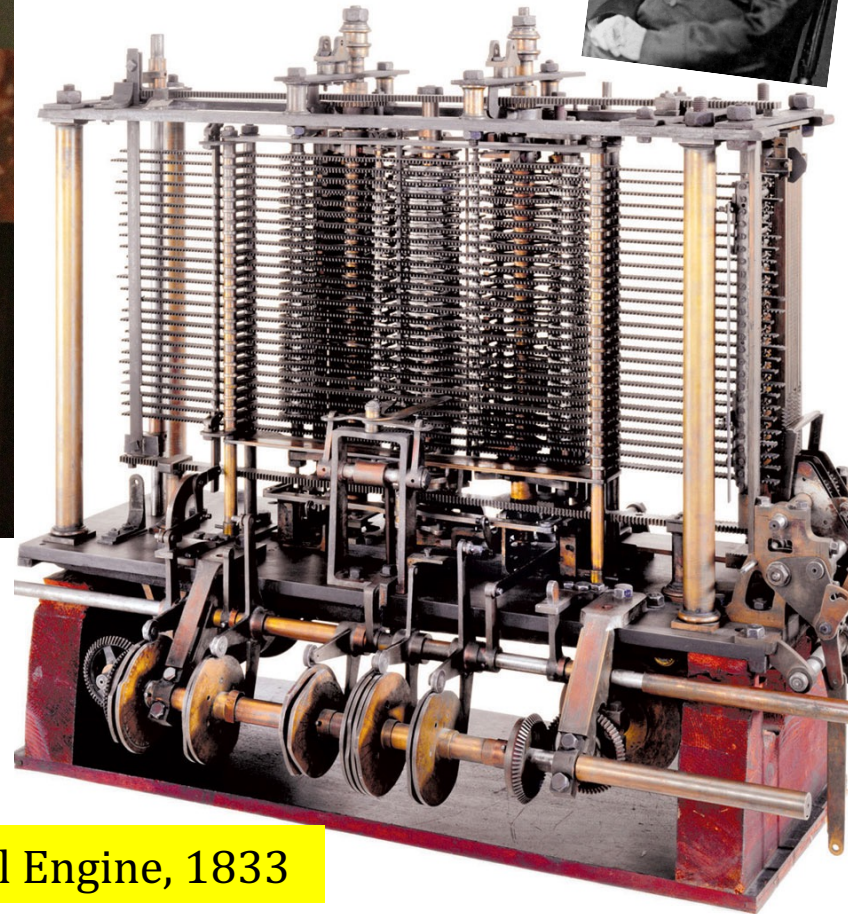
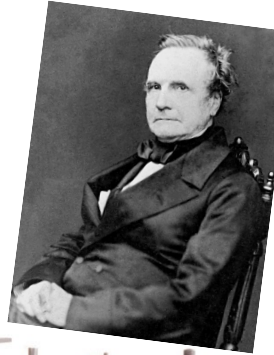


Fun with control?



Early Binary Control...

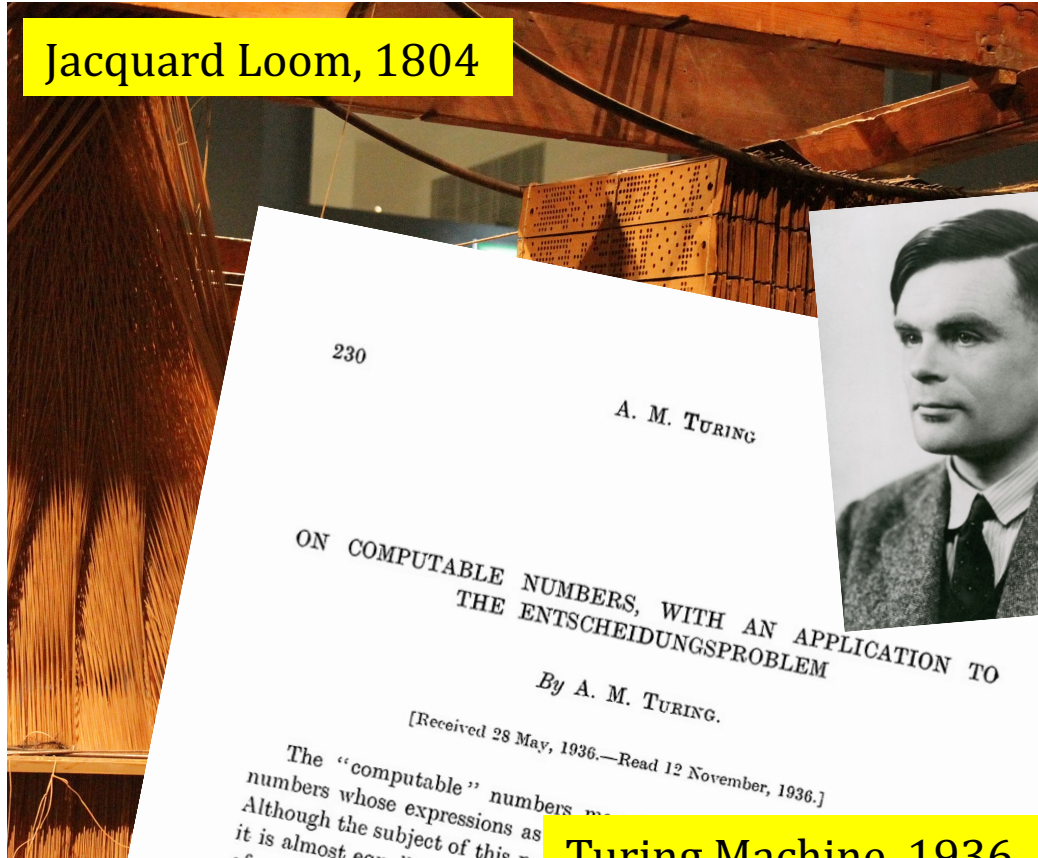
Jacquard Loom, 1804



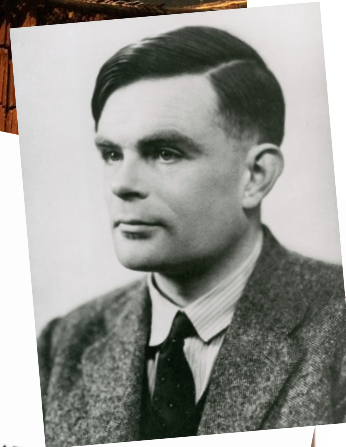
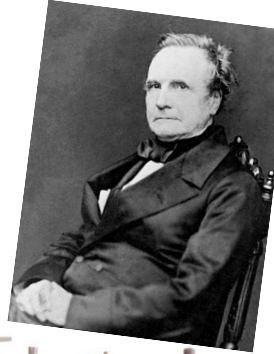
Babbage's Analytical Engine, 1833

Big idea: Control = Data

Jacquard Loom, 1804



A machine can use the same kind of storage for both code and data!



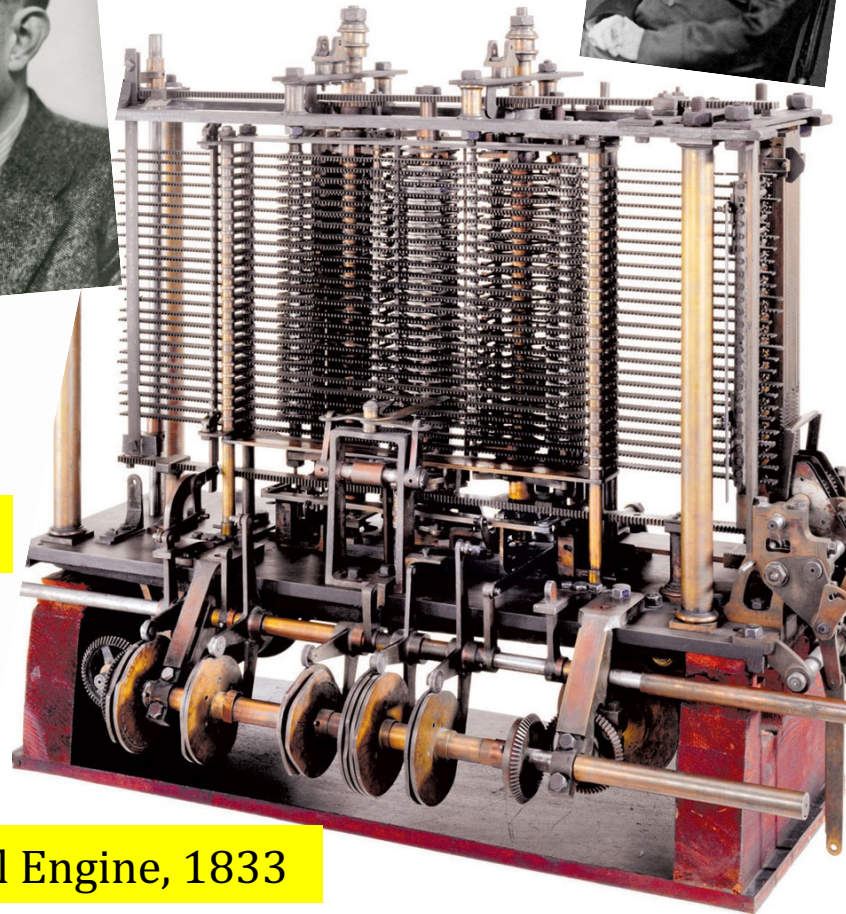
230
A. M. TURING
ON COMPUTABLE NUMBERS, WITH AN APPLICATION TO
THE ENTSCHEIDUNGSPROBLEM
By A. M. TURING.

[Received 28 May, 1936.—Read 12 November, 1936.]

Turing Machine, 1936

The "computable" numbers may be defined as the real numbers whose expressions as series of rational numbers can be obtained by a finite means. Although the subject of this paper is almost equally easy to define and investigate computable functions, it is almost equally easy to define and investigate computable numbers, of an integral variable or a real or computable variable, computable predicates, and so forth. The fundamental problems involved are, however, the same in each case, and I have chosen the computable numbers for explicit treatment as involving the least cumbersome technique. I hope shortly to give an account of the relations of the computable numbers, functions, and so forth to one another. This will include a development of the theory of functions of a real variable expressed in terms of computable numbers. According to my definition, a number is computable if its decimal can be written down by a machine.

In §§ 9, 10 I give some arguments with the intention of showing that computable numbers include all numbers which are regarded as computable. In particular, all algebraic numbers are computable.



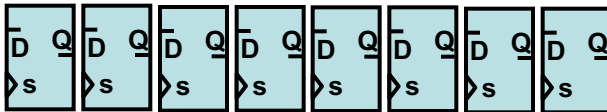
Cal Engine, 1833

Some memory is more equal than others...

Registers

on the Central Processing Unit

8 flip-flops are an 8-bit **register**



100 Registers of 64 bits each
~ 10,000 bits

memory from
logic gates

Main Memory (replaceable RAM)



10 GB memory
~ 100 billion bits

"Leaky Bucket"
capacitors

Disk Drive magnetic storage



4 TB drive
~ 42 trillion bits (or more)

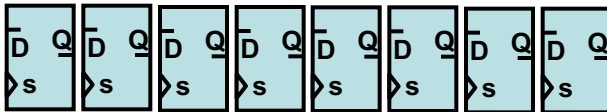
remagnetizing
surfaces

Some memory is more equal than others...

Registers

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8 flip-flops are an 8-bit **register**



100 Registers of 64 bits each

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Main Memory (replaceable RAM)



10 GB memory

~ 100 billion bits

Disk Drive magnetic storage



4 TB drive

~ 42 trillion bits (or more)

Price

~\$50

~\$50

~\$50

Time

1 clock cycle
 10^{-9} sec

100 cycles
 10^{-7} sec

10^7 cycles
 10^{-2} sec

If a clock cycle
== 1 minute

1 min

1.5 hours

19 YEARS

At least at my store!

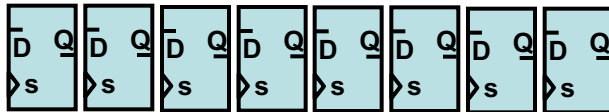


Some memory is more equal than others...

Registers

on the Central Processing Unit

8 flip-flops are an 8-bit **register**



100 Registers of 64 bits each

~ 10,000 bits

**programs are
fetched and
executed 1
instruction at a
time here...**

1 min

Main Memory (replaceable RAM)



10 GB memory

~ 100 billion bits

**running
programs
are stored
here...**

1.5 hours

Disk Drive magnetic storage



4 TB drive

~ 42 trillion bits (or more)

**"Off" data is
saved way
out here...**

10^{-2} sec

19 YEARS

If a c...
== 1 minute

How do we execute *sequences* of operations?

processor

CPU

stores all instructions and almost all data

RAM

live memory

the instruction's bits select which circuit to use...

sends next instruction to the CPU ...

divider

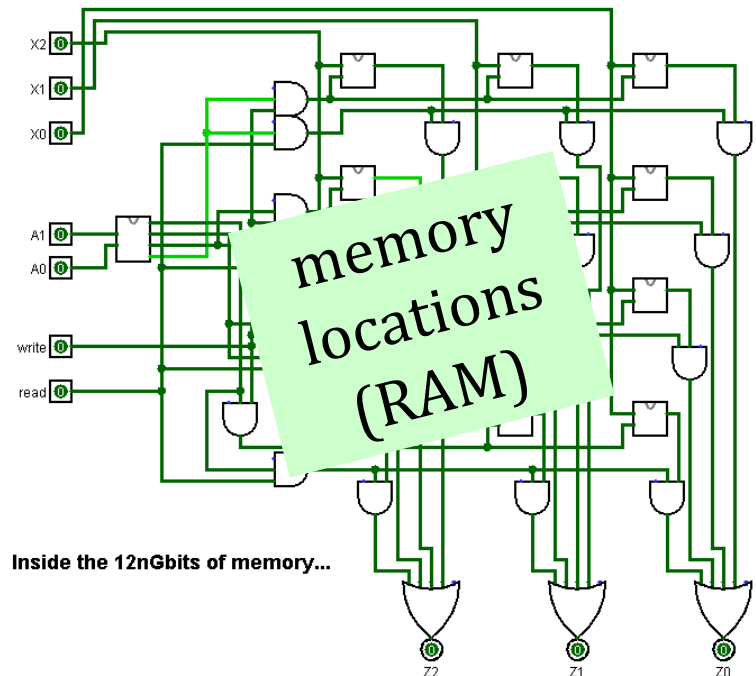
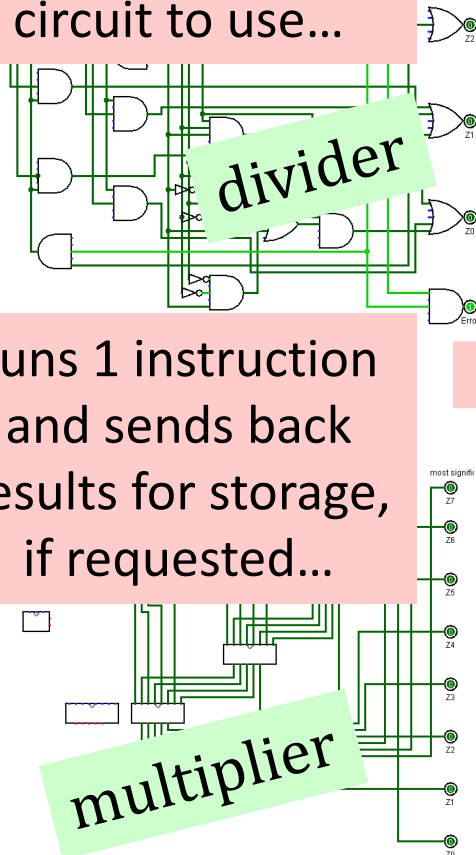
runs 1 instruction and sends back results for storage, if requested...

memory locations (RAM)

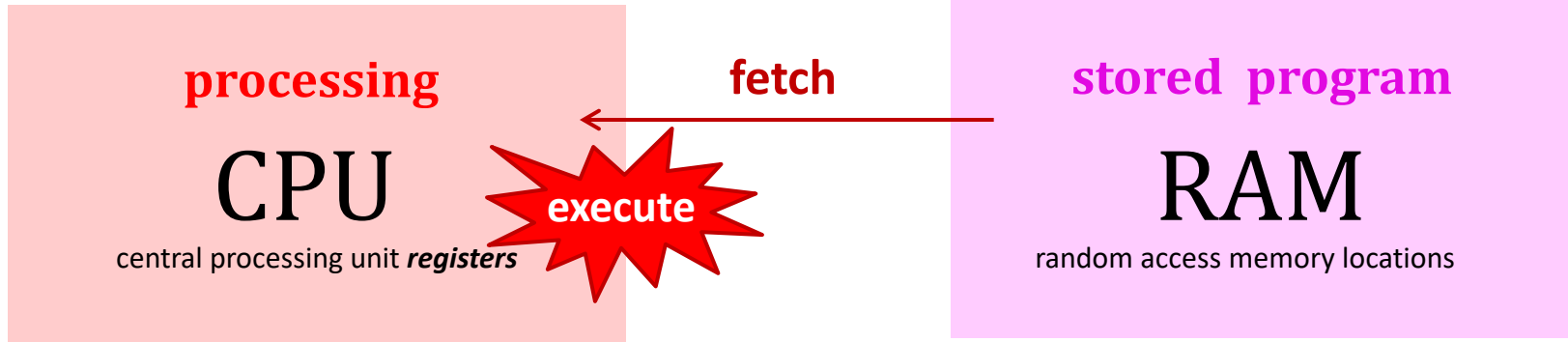
multiplier

Inside the 12nGbits of memory...

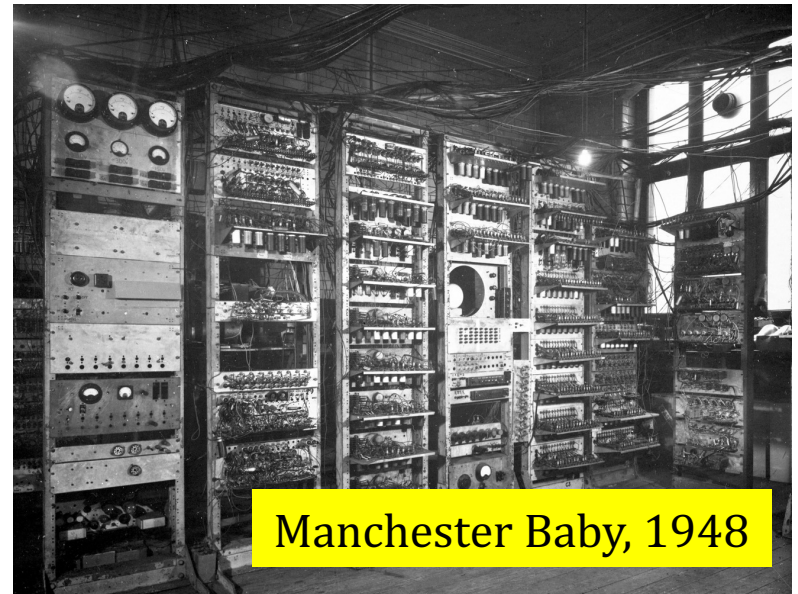
sends next instruction to the CPU ...



75 years ago...



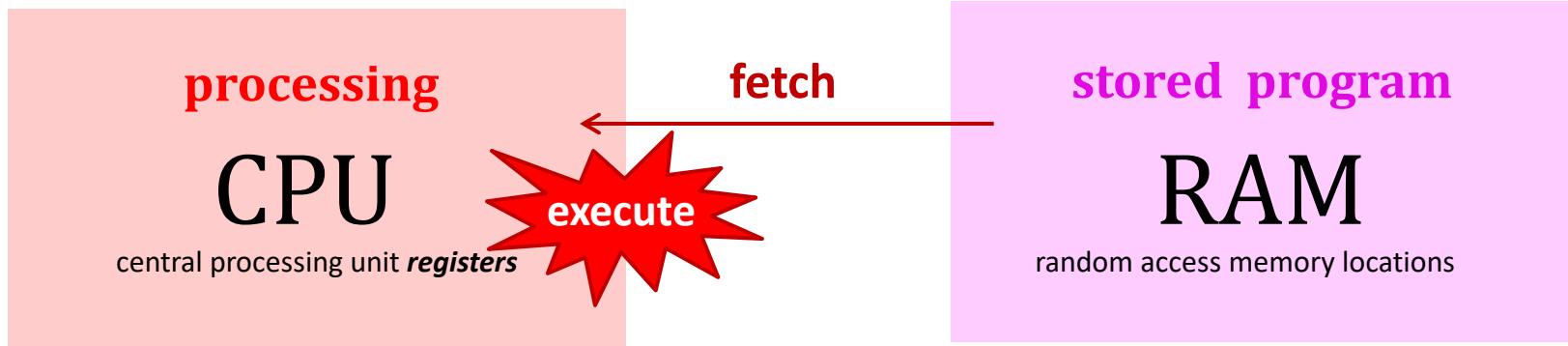
limited, fast **registers**
+ arithmetic



Manchester Baby, 1948

larger, slower **memory**
+ *no* computation

75 years later...

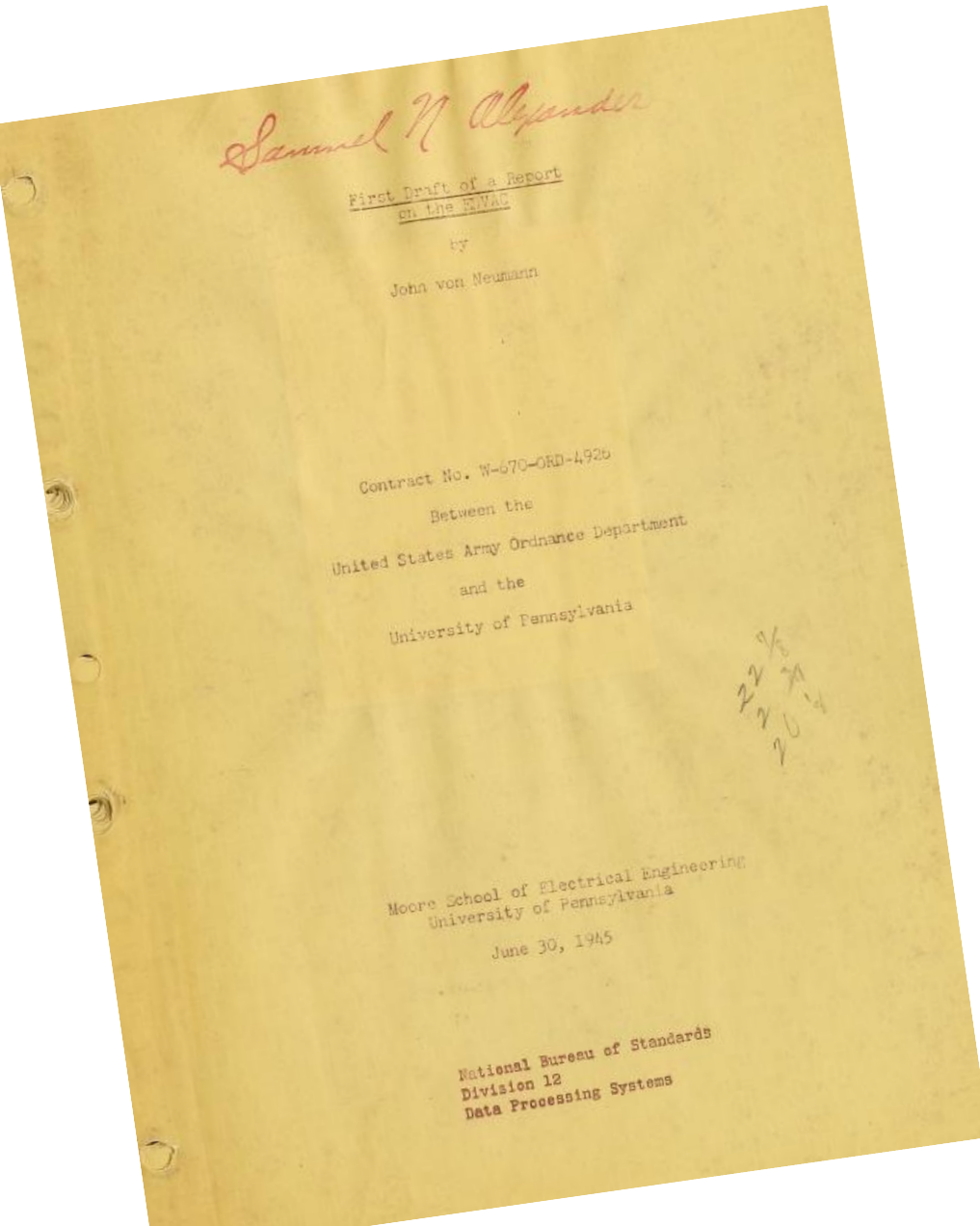


limited, fast **registers**
+ arithmetic



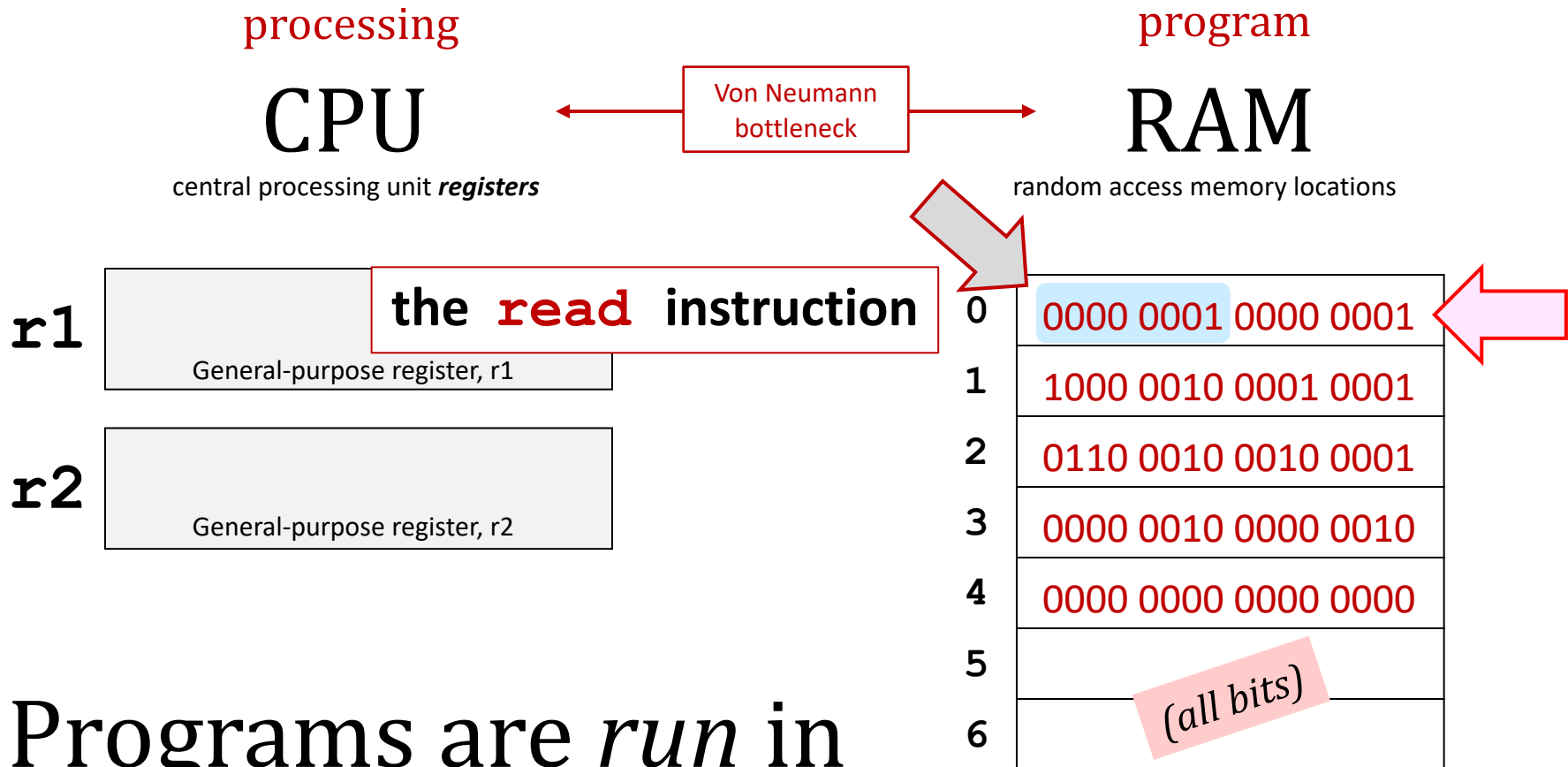
larger, slower **memory**
+ *no* computation

John von Neumann



- Polymath
- On EDVAC team...
 - Wasn't first stored-program computer!
- Based on the work of J. Presper Eckert and John Mauchly and other EDIAC/EDVAC designers.
 - Prevented their patent.

“Von Neumann” Architecture



Programs are run in
machine language

The Hmmm Instruction Set

There are 26 different instructions in Hmmm, each of which accepts between 0 and 3 arguments. Two of the instructions, `setn` and `addn`, accept a signed numerical argument between -128 and 127. The `load`, `store`, `call`, and `jump` instructions accept an unsigned numerical argument between 0 and 255. All other instruction arguments are registers. In the code below, register arguments will be represented by 'rX', 'rY', and 'rZ', while numerical arguments will be represented by '#'. In real code, any of the 16 registers could take the place of 'rX' 'rY' or 'rZ'. The available instructions are:

| Assembly | Binary | Description |
|-----------------------|---------------------|---|
| halt | 0000 0000 0000 0000 | Halt program |
| nop | 0110 0000 0000 0000 | Do nothing |
| read rX | 0000 xxxx 0000 0001 | Stop for user input, which will then be stored in register rX (input is an integer from -32768 to +32767). Prints "Enter number: " to prompt user for input |
| write rX | 0000 xxxx 0000 0010 | Print the contents of register rX on standard output |
| setn rX, # | 0001 xxxx #### #### | Load an 8-bit integer # (-128 to +127) into register rX |
| loadr rX, rY | 0100 xxxx yyyy 0000 | Load register rX from memory word addressed by rY: $rX = \text{memory}[rY]$ |
| storer rX, rY | 0100 xxxx yyyy 0001 | Store contents of register rX into memory word addressed by rY: $\text{memory}[rY] = rX$ |
| popr rX rY | 0100 xxxx yyyy 0010 | Load contents of register rX from stack pointed to by register rY: $rY -= 1$; $rX = \text{memory}[rY]$ |
| pushr rX rY | 0100 xxxx yyyy 0011 | Store contents of register rX onto stack pointed to by register rY: $\text{memory}[rY] = rX$; $rY += 1$ |
| loadn rX, # | 0010 xxxx #### #### | Load register rX with memory word at address # |
| storen rX, # | 0011 xxxx #### #### | Store contents of register rX into memory word at address # |
| addn rX, # | 0101 xxxx #### #### | Add the 8-bit integer # (-128 to 127) to register rX |
| copy rX, rY | 0110 xxxx yyyy 0000 | Set $rX = rY$ |
| neg rX, rY | 0111 xxxx 0000 yyyy | Set $rX = -rY$ |
| add rX, rY, rZ | 0110 xxxx yyyy zzzz | Set $rX = rY + rZ$ |
| sub rX, rY, rZ | 0111 xxxx yyyy zzzz | Set $rX = rY - rZ$ |
| mul rX, rY, rZ | 1000 xxxx yyyy zzzz | Set $rX = rY * rZ$ |
| div rX, rY, rZ | 1001 xxxx yyyy zzzz | Set $rX = rY // rZ$ |
| mod rX, rY, rZ | 1010 xxxx yyyy zzzz | Set $rX = rY \% rZ$ |
| jumpr rX | 0000 xxxx 0000 0011 | Set program counter to address in rX |
| jumpn n | 1011 0000 #### #### | Set program counter to address # |
| jeqzn rX, # | 1100 xxxx #### #### | If $rX = 0$ then set program counter to address # |
| jnezn rX, # | 1101 xxxx #### #### | If $rX \neq 0$ then set program counter to address # |
| jgtzn rX, # | 1110 xxxx #### #### | If $rX > 0$ then set program counter to address # |
| jltzn rX, # | 1111 xxxx #### #### | If $rX < 0$ then set program counter to address # |
| calln rX, # | 1011 xxxx #### #### | Set rX to (next) program counter, then set program counter to address # |

Machine Language

The Hmmm Instruction Set

There are 26 different instructions in Hmmm, each of which accepts between 0 and 3 arguments. Two of the instructions, `setn` and `addn`, accept a signed numerical argument between -128 and 127. The `load`, `store`, `call`, and `jump` instructions accept an unsigned numerical argument between 0 and 255. All other instruction arguments are registers. In the code below, register arguments will be represented by 'rX', 'rY', and 'rZ', while numerical arguments will be represented by '#'. In real code, any of the 16 registers could take the place of 'rX' 'rY' or 'rZ'. The available instructions are:

| Assembly | Binary | Description |
|-----------------------------|---------------------|---|
| <code>halt</code> | 0000 0000 0000 0000 | Halt program |
| <code>nop</code> | 0110 0000 0000 0000 | Do nothing |
| <code>read rX</code> | 0000 xxxx 0000 0001 | Stop for user input, which will then be stored in register rX (input is an integer from -32768 to +32767). Prints "Enter number: " to prompt user for input |
| <code>write rX</code> | 0000 xxxx 0010 0000 | Print the contents of register rX on standard output |
| <code>setn rX, #</code> | 0001 xxxx # 0000 | Load an 8-bit integer # (-128 to +127) into register rX |
| <code>loadr rX, rY</code> | 0100 xxxx YYY 0000 | Load memory word addressed by rY: $rX = \text{memory}[rY]$ |
| <code>storer rX, rY</code> | 0100 xxxx YYY 0001 | Store rX into memory word addressed by rY: $\text{memory}[rY] = rX$ |
| <code>popr rX rY</code> | 0100 xxxx YYY 0010 | Pop rX from stack pointed to by register rY: $rY -= 1$; $rX = \text{memory}[rY]$ |
| <code>pushr rX rY</code> | 0100 xxxx YYY 0011 | Push rX onto stack pointed to by register rY: $\text{memory}[rY] = rX$; $rY += 1$ |
| <code>loadn rX, #</code> | 0010 xxxx #### | Load memory word at address # |
| <code>storen rX, #</code> | 0011 xxxx #### | Store rX into memory word at address # |
| <code>addn rX, #</code> | 0101 xxxx #### | Add the 8-bit integer # (-128 to 127) to register rX |
| <code>copy rX, rY</code> | 0110 xxxx YYY 0000 | Set $rX = rY$ |
| <code>neg rX, rY</code> | 0111 xxxx 0000 YYY | Set $rX = -rY$ |
| <code>add rX, rY, rZ</code> | 0110 xxxx YYY ZZZ | Set $rX = rY + rZ$ |
| <code>sub rX, rY, rZ</code> | 0111 xxxx YYY ZZZ | Set $rX = rY - rZ$ |
| <code>mul rX, rY, rZ</code> | 1000 xxxx YYY ZZZ | Set $rX = rY * rZ$ |
| <code>div rX, rY, rZ</code> | 1001 xxxx YYY ZZZ | Set $rX = rY / rZ$ |
| <code>mod rX, rY, rZ</code> | 1010 xxxx YYY ZZZ | Set $rX = rY \% rZ$ |
| <code>jumpr rX</code> | 0000 xxxx 0000 0011 | Jump to address # |
| <code>jumpn n</code> | 1011 0000 #### | Jump to address # |
| <code>jeqzn rX, #</code> | 1100 xxxx #### | If $rX == \#$ then set program counter to address # |
| <code>jnezn rX, #</code> | 1101 xxxx #### | If $rX \neq \#$ then set program counter to address # |
| <code>jgtzn rX, #</code> | 1110 xxxx #### | If $rX > \#$ then set program counter to address # |
| <code>jltzn rX, #</code> | 1111 xxxx #### | If $rX < \#$ then set program counter to address # |
| <code>calln rX, #</code> | 1011 xxxx #### | Set rX to (next) program counter, then set program counter to address # |

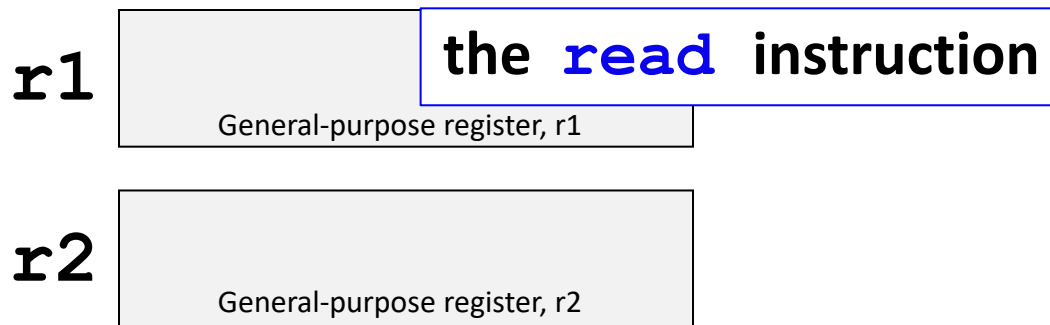
the **read** instruction

which register to **read** into?

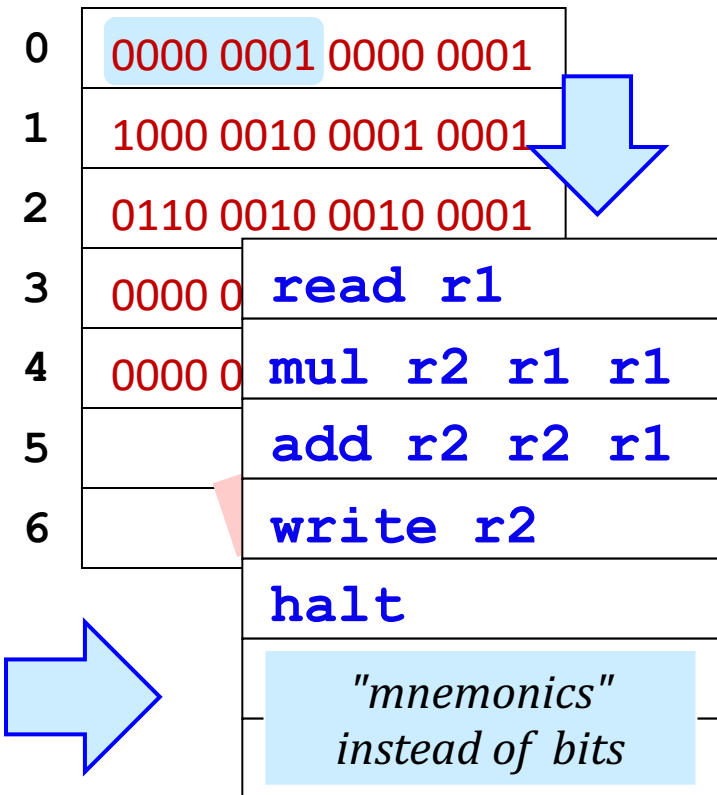
the "bitpatterns" do matter!

Machine Language

“Von Neumann” Architecture



Programs are shown
in *assembly language*



The Hmmm Instruction Set

There are 26 different instructions in Hmmm, each of which accepts between 0 and 3 arguments. Two of the instructions, `setn` and `addn`, accept a signed numerical argument between -128 and 127. The `load`, `store`, `call`, and `jump` instructions accept an unsigned numerical argument between 0 and 255. All other instruction arguments are registers. In the code below, register arguments will be represented by 'rX', 'rY', and 'rZ', while numerical arguments will be represented by '#'. In real code, any of the 16 registers could take the place of 'rX' 'rY' or 'rZ'. The available instructions are:

| Assembly | Binary | Description |
|-----------------------------|---------------------|--|
| <code>halt</code> | 0000 0000 0000 0000 | Stop the program |
| <code>nop</code> | 0111 0000 0000 0000 | No operation |
| <code>read rX</code> | 0000 0000 0000 0000 | Read an integer from standard input and store it in register rX (input is an integer from -32768 to +32767). |
| <code>write rX</code> | 0000 xxxx 0000 0000 | Write the integer in register rX to standard output |
| <code>setn rX, #</code> | 0101 xxxx ##### | Set register rX to 8-bit integer # (-128 to +127) into register rX |
| <code>loadr rX, rY</code> | 0110 xxxx ##### | Load the integer in register rY from memory word addressed by rY: $rX = \text{memory}[rY]$ |
| <code>storer rX, rY</code> | 0111 xxxx ##### | Store the integer in register rX into memory word addressed by rY: $\text{memory}[rY] = rX$ |
| <code>popr rX rY</code> | 0111 xxxx ##### | Pop the integer from stack pointed to by register rY: $rY -= 1$; $rX = \text{memory}[rY]$ |
| <code>pushr rX rY</code> | 0111 xxxx ##### | Push the integer in register rX onto stack pointed to by register rY: $\text{memory}[rY] = rX$; $rY += 1$ |
| <code>loadn rX, #</code> | 0000 0000 ##### | Load the integer # into register rX with memory word at address # |
| <code>storen rX, #</code> | 0000 0000 ##### | Store the integer in register rX into memory word at address # |
| <code>addn rX, #</code> | 0101 xxxx ##### | Add the 8-bit integer # (-128 to +127) to register rX |
| <code>copy rX, rY</code> | 0110 xxxx ##### | Set $rX = rY$ |
| <code>neg rX, rY</code> | 0111 xxxx ##### | Set $rX = -rY$ |
| <code>add rX, rY, rZ</code> | 0110 xxxx ##### | Set $rX = rY + rZ$ |
| <code>sub rX, rY, rZ</code> | 0111 xxxx ##### | Set $rX = rY - rZ$ |
| <code>mul rX, rY, rZ</code> | 1000 xxxx ##### | Set $rX = rY * rZ$ |
| <code>div rX, rY, rZ</code> | 1001 xxxx ##### | Set $rX = rY // rZ$ |
| <code>mod rX, rY, rZ</code> | 1010 xxxx ##### | Set $rX = rY \% rZ$ |
| <code>jumpr rX</code> | 0000 xxxx 0000 0011 | Jump to address # if register rX is not zero |
| <code>jumpn n</code> | 1011 0000 ##### | Jump to address # |
| <code>jeqzn rX, #</code> | 1100 xxxx ##### | Jump to address # if register rX is equal to address # |
| <code>jnezn rX, #</code> | 1101 xxxx ##### | Jump to address # if register rX is not equal to address # |
| <code>jgtzn rX, #</code> | 1110 xxxx ##### | Jump to address # if register rX is greater than address # |
| <code>jltzn rX, #</code> | 1111 xxxx ##### | Jump to address # if register rX is less than address # |
| <code>calln rX, #</code> | 1011 xxxx ##### | Set rX to (next) program counter, then set program counter to address # |

the read instruction

which register to read into?

the "bitpatterns" don't matter!

Assembly Language

The Hmmm I

There are 26 diff
signed numerical
and 255. All othe
numerical argum
instructions are:

Documentation for HMMM (Harvey Mudd Miniature Machine)

Last update: 2024

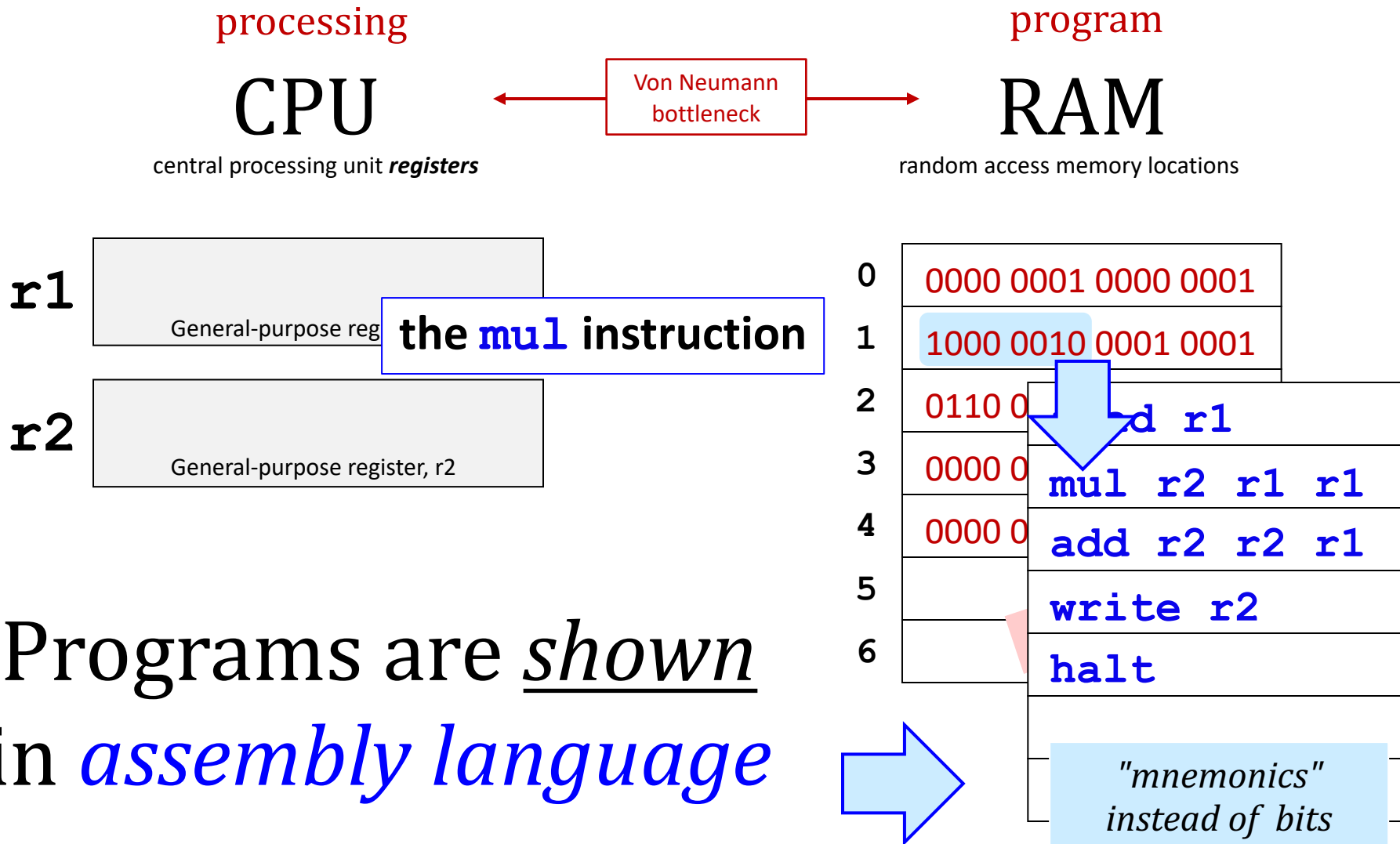
Quick reference: Table of Hmmm Instructions

| Instruction | Description | Aliases |
|-------------------------------|--|---------------|
| System instructions | | |
| halt | Stop! | |
| read rX | Place user input in register rX | |
| write rX | Print contents of register rX | |
| nop | Do nothing | |
| Setting register data | | |
| setn rX N | Set register rX equal to the integer N (-128 to +127) | |
| addn rX N | Add integer N (-128 to 127) to register rX | |
| copy rX rY | Set rX = rY | |
| Arithmetic | | |
| add rX rY rZ | Set rX = rY + rZ | |
| sub rX rY rZ | Set rX = rY - rZ | |
| neg rX rY | Set rX = -rY | |
| mul rX rY rZ | Set rX = rY * rZ | |
| div rX rY rZ | Set rX = rY // rZ (integer di | |
| mod rX rY rZ | Set rX = rY % rZ (returns the | |
| Jumps! | | |
| jumpn N | Set program counter to address N | |
| jumpr rX | Set program counter to address in rX | |
| jeqzn rX N | If rX == 0, then jump to line N | jump |
| jnezn rX N | If rX != 0, then jump to line N | jeqz |
| jgtzn rX N | If rX > 0, then jump to line N | jnez |
| jltzn rX N | If rX < 0, then jump to line N | jgtz |
| calln rX N | Copy addr. of next instr. into rX and then jump to mem. addr. N | jltz |
| Interacting with memory (RAM) | | |
| pushr rX rY | Store contents of register rX onto stack pointed to by reg. rY | |
| popr rX rY | Load contents of register rX from stack pointed to by reg. rY | |
| loadn rX N | Load register rX with the contents of memory address N | |
| storen rX N | Store contents of register rX into memory address N | |
| loadr rX rY | Load register rX with data from the address location held in reg. rY | loadi, load |
| storer rX rY | Store contents of register rX into memory address held in reg. rY | storei, store |

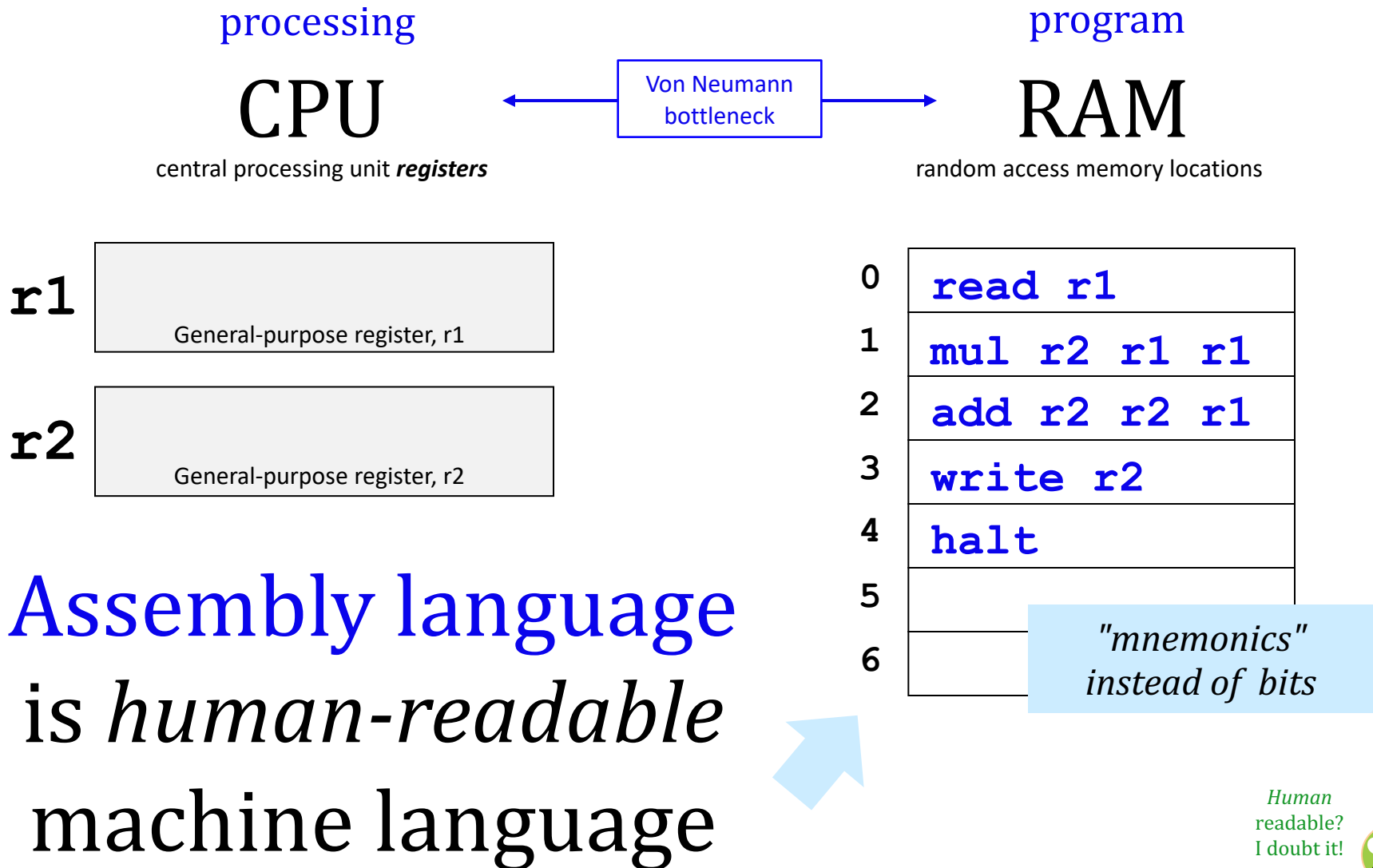
the realm of
"instructions"

Assembly Language

“Von Neumann” Architecture



“Von Neumann” Architecture



Human
readable?
I doubt it!



Example #1:

Screen

6 (input)

a five-line assembly-
language program



0 **read** r1

6

1 **mul** r2 r1 r1

2 **add** r2 r2 r1

3 **write** r2

4 **halt**

Demo

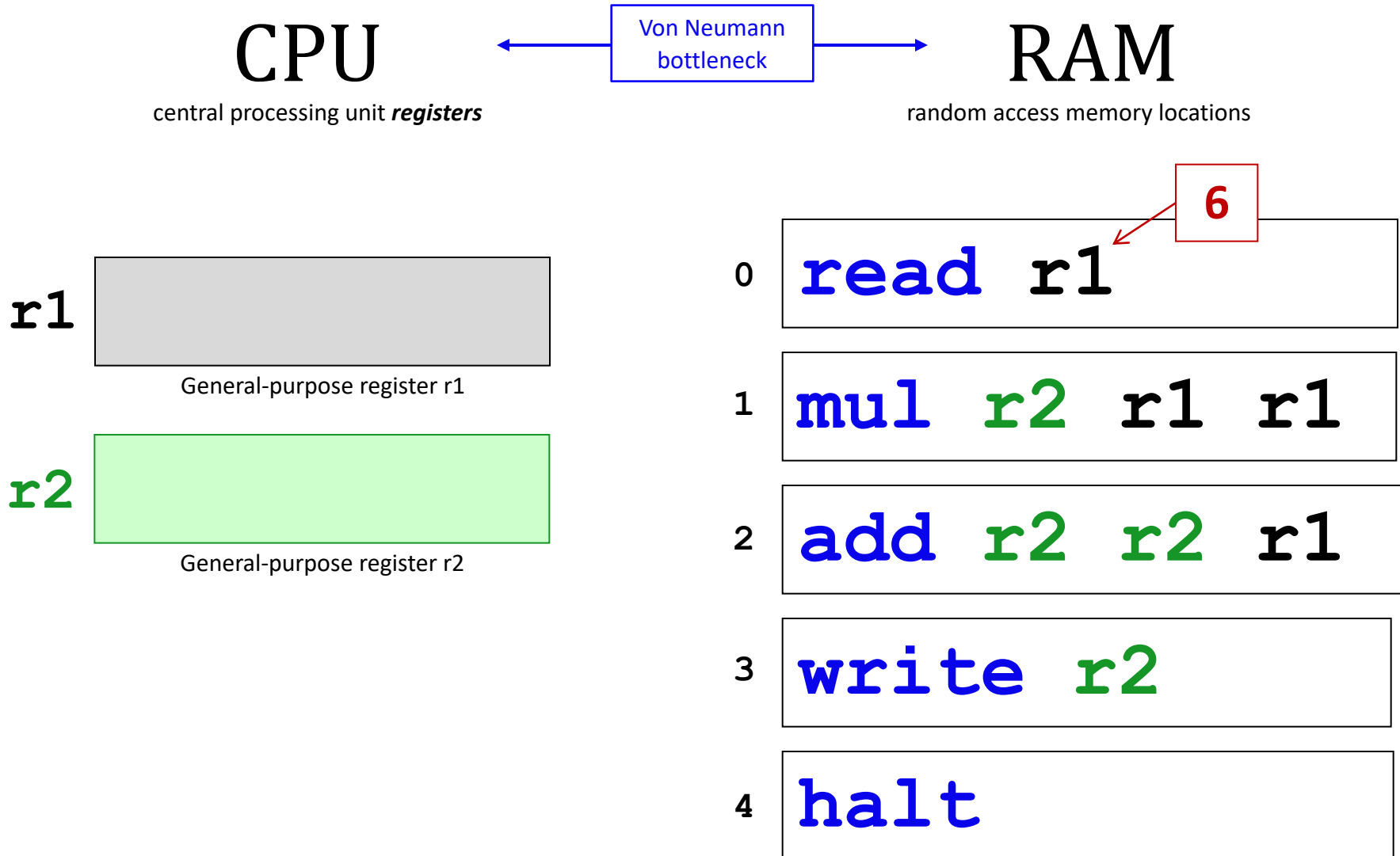
of assembly-language programming in **Hmmm**...

in hw6, CS
stands for
Chin-
Scratching?!

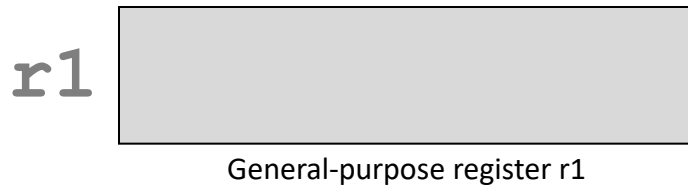
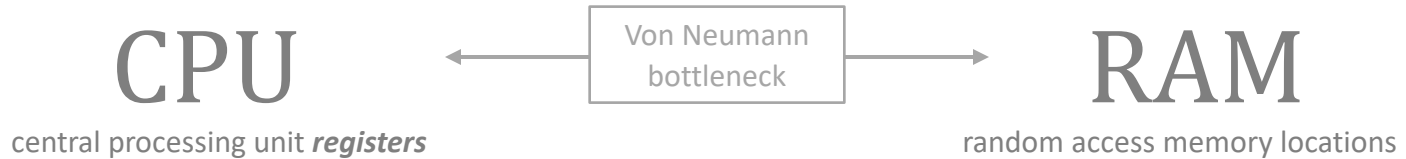
Example #1:

Screen

6 (input)



Hmmm: Harvey mudd miniature machine



16 registers



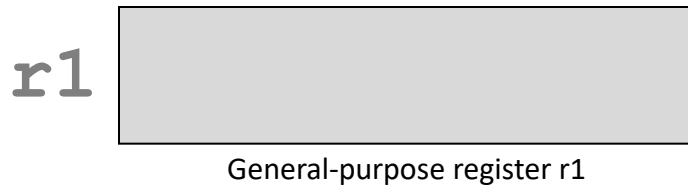
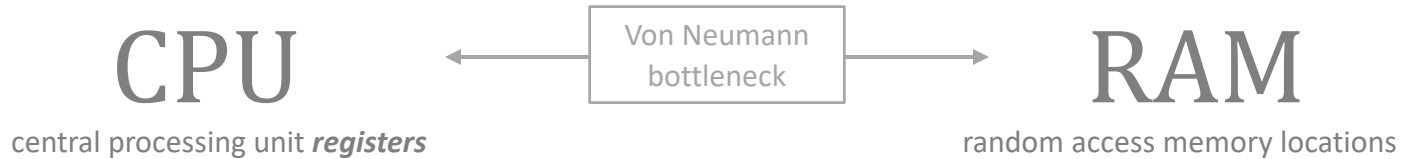
256 memory
locations

vs. 2024 ?

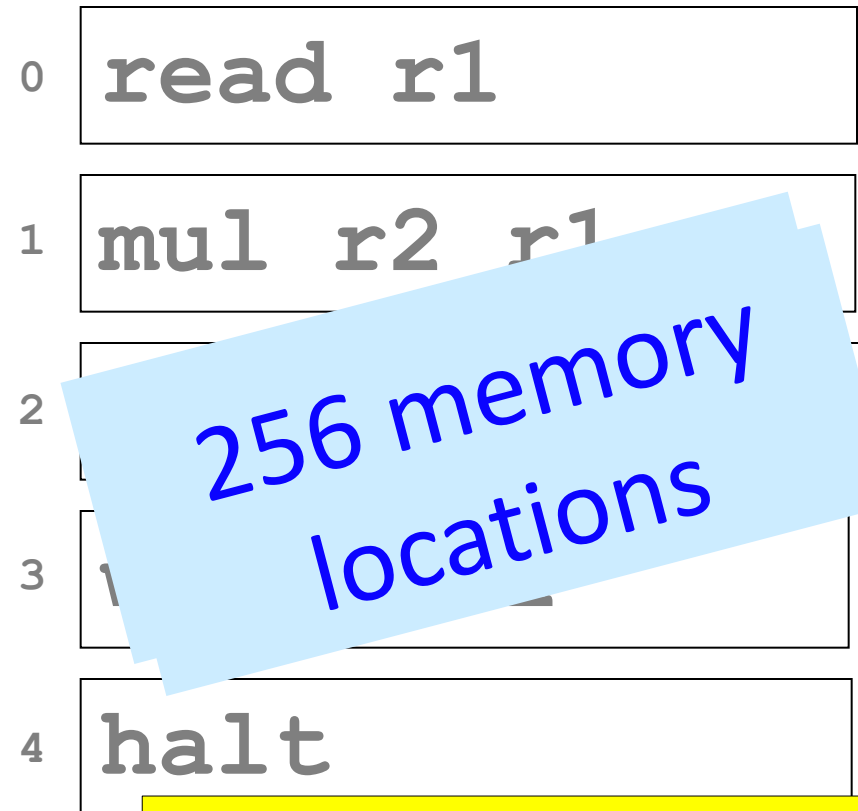
Really, it's only 15,
r0 is special



Hmmm vs 2024



16 registers



256 memory locations

2022 Arm M1: 37-40 registers per core

2024: ~16,000,000,000 mem loc's

Why Assembly?

| Oct 2018 | Oct 2017 | Change | Programming Language | Ratings | Change |
|----------|----------|--------|----------------------|---------|--------|
| 1 | 1 | | Java | 17.801% | +5.37% |
| 2 | 2 | | C | 15.376% | +7.00% |
| 3 | 3 | | C++ | 7.593% | +2.59% |
| 4 | 5 | ⬆ | Python | 7.156% | +3.35% |
| 5 | 8 | ⬆ | Visual Basic .NET | 5.884% | +3.15% |
| 6 | 4 | ⬇ | C# | 3.485% | -0.37% |
| 7 | 7 | | PHP | 2.794% | +0.00% |
| 8 | 6 | ⬇ | JavaScript | 2.280% | -0.73% |
| 9 | - | ⬆ | SQL | 2.038% | +2.04% |
| 10 | 16 | ⬆ | Swift | 1.500% | -0.17% |
| 11 | 13 | ⬆ | MATLAB | 1.317% | -0.56% |
| 12 | 20 | ⬆ | Go | 1.253% | -0.10% |
| 13 | 9 | ⬇ | Assembly language | 1.245% | -1.13% |
| 14 | 15 | ⬆ | R | 1.214% | -0.47% |
| 15 | 17 | ⬆ | Objective-C | 1.202% | -0.31% |

15

14



R



Unsafe vehicles, hills, and philosophy go hand in hand.

2018

Why Assembly?

| Oct 2019 | Oct 2018 | Change | Programming Language | Ratings | Change |
|----------|----------|--------|----------------------|---------|--------|
| 1 | 1 | | Java | 16.884% | -0.92% |
| 2 | 2 | | C | 16.180% | +0.80% |
| 3 | 4 | ⬆ | Python | 9.089% | +1.93% |
| 4 | 3 | ⬇ | C++ | 6.229% | -1.36% |
| 5 | 6 | ⬆ | C# | 3.860% | +0.37% |
| 6 | 5 | ⬇ | Visual Basic .NET | 3.745% | -2.14% |
| 7 | 8 | ⬆ | JavaScript | 2.076% | -0.20% |
| 8 | 9 | ⬆ | SQL | 1.935% | -0.10% |
| 9 | 7 | ⬇ | PHP | 1.909% | -0.89% |
| 10 | 15 | ⬆ | Objective-C | 1.501% | +0.30% |
| 11 | 28 | ⬆ | Groovy | 1.394% | +0.96% |
| 12 | 10 | ⬇ | Swift | 1.362% | -0.14% |
| 13 | 18 | ⬆ | Ruby | 1.318% | +0.21% |
| 14 | 13 | ⬇ | Assembly language | 1.307% | +0.06% |
| 15 | 14 | ⬇ | R | | |

2019

Why Assembly?









| Oct 2018 | Oct 2017 | Change | Programming Language | Ratings | Change |
|----------|----------|--------|----------------------|---------|--------|
| 1 | 1 | | Java | 16.884% | 0.02% |

| May 2021 | May 2020 | Change | Programming Language | Ratings | Change |
|----------|----------|--------|----------------------|---------|--------|
| 1 | 1 | | C | 13.38% | -3.68% |
| 2 | 3 | ⬆ | Python | 11.87% | +2.75% |
| 3 | 2 | ⬇ | Java | 11.74% | -4.54% |
| 4 | 4 | | C++ | 7.81% | +1.69% |
| 5 | 5 | | C# | 4.41% | +0.12% |
| 6 | 6 | | Visual Basic | 4.02% | -0.16% |
| 7 | 7 | | JavaScript | 2.45% | -0.23% |
| 8 | 14 | ⬆ | Assembly language | 2.43% | +1.31% |

| | | | | | |
|----|----|---|-------------------|--|--------|
| 13 | 18 | ⬆ | Ruby | | |
| 14 | 13 | ⬇ | Assembly language | | |
| 15 | 14 | ⬇ | R | | +0.05% |

2021...

Why Assembly?

| May 2022 | May 2021 | Change | Programming Language | | Ratings | Change |
|----------|----------|--------|---|-------------------|---------|--------|
| 1 | 2 | ▲ |  | Python | 12.74% | +0.86% |
| 2 | 1 | ▼ |  | C | 11.59% | -1.80% |
| 3 | 3 | |  | Java | 10.99% | -0.74% |
| 4 | 4 | |  | C++ | 8.83% | +1.01% |
| 5 | 5 | |  | C# | 6.39% | +1.98% |
| 6 | 6 | |  | Visual Basic | 5.86% | +1.85% |
| 7 | 7 | |  | JavaScript | 2.12% | -0.33% |
| 8 | 8 | |  | Assembly language | 1.92% | -0.51% |

| Change |
|--------|
| -3.68% |
| +2.75% |
| -4.54% |
| +1.69% |
| +0.12% |
| -0.16% |
| -0.23% |
| +1.31% |















May 2022 ...

Unsafe vehicles, hills, and philosophy go hand in hand.

Why Assembly?

TIOBE Index for October 2022

October Headline: The big 4 languages keep increasing their dominance

| Oct 2022 | Oct 2021 | Change | Programming Language | | Ratings | Change |
|----------|----------|--------|---|-------------------|---------|--------|
| 1 | 1 | |  | Python | 17.08% | +5.81% |
| 2 | 2 | |  | C | 15.21% | +4.05% |
| 3 | 3 | |  | Java | 12.84% | +2.38% |
| 4 | 4 | |  | C++ | 9.92% | +2.42% |
| 5 | 5 | |  | C# | 4.42% | -0.84% |
| 6 | 6 | |  | Visual Basic | 3.95% | -1.29% |
| 7 | 7 | |  | JavaScript | 2.74% | +0.55% |
| 8 | 10 | ▲ |  | Assembly language | 2.39% | +0.33% |
| 9 | 9 | |  | PHP | 2.04% | -0.06% |
| 10 | 8 | ▼ |  | SQL | 1.78% | -0.39% |
| 11 | 12 | ▲ |  | Go | 1.27% | -0.01% |
| 12 | 14 | ▲ |  | R | 1.22% | +0.03% |
| 13 | 29 | ▲▲ |  | Objective-C | 1.21% | +0.55% |
| 14 | 13 | ▼ |  | MATLAB | | |

Change

-3.68%

+2.75%

-4.54%

+1.69%

+0.12%

-0.16%

-0.23%

+1.31%

October 2022

Unsafe vehicles, hills, and philosophy go hand in hand.

Feb 2024





















Feb 2023

Change


Programming Language

Ratings

Change

| | | | | | |
|----|----|----|--|--------|--------|
| 1 | 1 | |  Python | 15.16% | -0.32% |
| 2 | 2 | |  C | 10.97% | -4.41% |
| 3 | 3 | |  C++ | 10.53% | -3.40% |
| 4 | 4 | |  Java | 8.88% | -4.33% |
| 5 | 5 | |  C# | 7.53% | +1.15% |
| 6 | 7 | ^ |  JavaScript | 3.17% | +0.64% |
| 7 | 8 | ^ |  SQL | 1.82% | -0.30% |
| 8 | 11 | ^ |  Go | 1.73% | +0.61% |
| 9 | 6 | v |  Visual Basic | 1.52% | -2.62% |
| 10 | 10 | |  PHP | 1.51% | +0.21% |
| 11 | 24 | ^^ |  Fortran | 1.40% | +0.82% |
| 12 | 14 | ^ |  Delphi/Object Pascal | 1.40% | +0.45% |
| 13 | 13 | |  MATLAB | 1.26% | +0.27% |
| 14 | 9 | vv |  Assembly language | 1.19% | -0.19% |
| 15 | 18 | ^ |  Scratch | 1.18% | +0.42% |
| 16 | 15 | v |  Swift | 1.16% | +0.23% |
| 17 | 33 | ^^ |  Kotlin | | |
| 18 | 20 | ^ |  Rust | | |
| 19 | 30 | ^^ |  COBOL | | |
| 20 | 16 | vv |  Ruby | | |

I FEEL LIKE



Change

-3.68%

+2.75%

-4.54%

+1.69%

+0.12%

-0.16%

-0.23%

+1.31%

2022

February 2024


















| Feb 2024 | Feb 2023 | Change | Programming Language | | Ratings | Change |
|----------|----------|--------|---|----------------------|---------|--------|
| 1 | 1 | |  | Python | 15.16% | -0.32% |
| 2 | 2 | |  | C | 10.97% | -4.41% |
| 3 | 3 | |  | C++ | 10.53% | -3.40% |
| 4 | 4 | |  | Java | 8.88% | -4.33% |
| 5 | 5 | |  | C# | 7.53% | +1.15% |
| 6 | 7 | ^ |  | JavaScript | 3.17% | +0.64% |
| 7 | 8 | ^ |  | SQL | 1.82% | -0.30% |
| 8 | 11 | ^ |  | Go | 1.73% | +0.61% |
| 9 | 6 | v |  | Visual Basic | 1.52% | -2.62% |
| 10 | 10 | |  | PHP | 1.51% | +0.21% |
| 11 | 24 | ^^ |  | Fortran | 1.40% | +0.82% |
| 12 | 14 | ^ |  | Delphi/Object Pascal | 1.40% | +0.45% |
| 13 | 13 | |  | MATLAB | 1.26% | +0.27% |
| 14 | 9 | v |  | Assembly language | 1.19% | -0.19% |
| 15 | 18 | |  | Scratch | 1.18% | +0.42% |
| | | |  | Swift | 1.16% | +0.23% |
| | | |  | Kotlin | | |
| | | |  | Rust | | |
| | | |  | COBOL | 1.01% | |
| | | |  | Ruby | | |

| Change |
|--------|
| -3.68% |
| +2.75% |
| -4.54% |
| +1.69% |
| +0.12% |
| -0.16% |
| -0.23% |
| +1.31% |

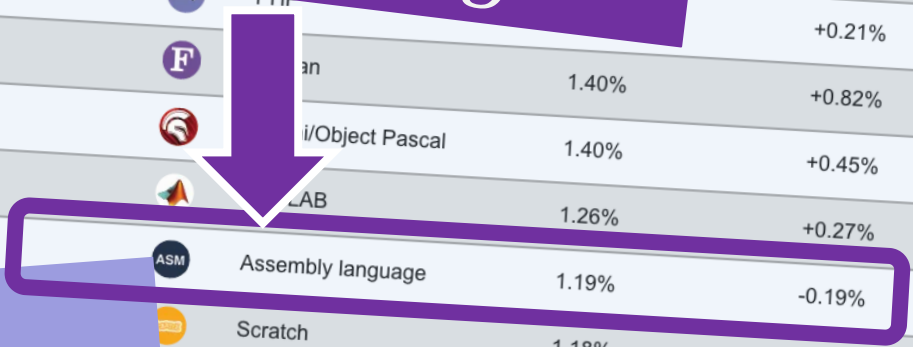
Software is written in many languages

February 2024

2022

| Feb 2024 | Feb 2023 | Change | Programming Language | Ratings | Change |
|----------|----------|--------|--|---------|--------|
| 1 | 1 | |  Python | 15.16% | -0.32% |
| 2 | 2 | |  C | 10.97% | -4.41% |
| 3 | 3 | |  C++ | 10.53% | -3.40% |
| 4 | 4 | |  Java | 8.88% | -4.33% |
| 5 | 5 | |  C# | 7.53% | +1.15% |
| 6 | | |  JavaScript | 3.17% | +0.64% |
| 7 | | | | | -0.30% |
| 8 | | | | | +0.61% |
| 9 | | | | | -2.62% |
| 10 | 10 | |  PHP | | +0.21% |
| 11 | 24 | ⬆️ |  Fortran | 1.40% | +0.82% |
| 12 | 14 | ⬆️ |  Pascal | 1.40% | +0.45% |
| 13 | 13 | |  LabVIEW | 1.26% | +0.27% |
| 14 | 9 | ⬇️ |  Assembly language | 1.19% | -0.19% |
| 15 | | |  Scratch | 1.18% | +0.42% |
| | | |  Swift | 1.16% | +0.23% |
| | | |  Kotlin | | |
| | | |  Rust | | |
| | | |  COBOL | 1.01% | |
| | | |  Ruby | | |

... but the CPU only RUNS in only one language!



Software is written in many languages

February 2024

2022

| Change |
|--------|
| -3.68% |
| +2.75% |
| -4.54% |
| +1.69% |
| +0.12% |
| -0.16% |
| -0.23% |
| +1.31% |

L

The Economist explains

Explaining the world, daily

[Previous](#) | [Next](#) | [Latest The Economist explains](#)

The Economist explains

What is code?

Sep 8th 2015, 23:50 BY T.S.

```
for i in people.data.users:
    response = client.api.statuse
    print 'Got', len(response.data)
    if len(response.data) != 0:
        ldate = response.data[0]
        ldate2 = datetime.strptime(ldate, '%a %b %d %H:%M:%S +0000 %Y')
        today = datetime.now()
        howlong = (today-ldate2).days
        if howlong < daywindow:
            print i.screen_name, 'has tweeted in the past', daywindow,
            totaltweets += len(response.data)
            for j in response.data:
                if j.entities.urls:
                    for k in j.entities.urls:
                        newurl = k['expanded_url']
                        urlset.add((newurl, j.user.screen_name))
        else:
            print i.screen_name, 'has not tweeted in the past', daywindow
```

FROM lifts to cars to airliners to smartphones, modern civilisation is powered by software, the digital instructions that allow computers, and the devices they control, to perform calculations and respond to their surroundings. How did that software get there? Someone had to write it. But code, the sequences of symbols painstakingly created by programmers, is not quite the same as software, the sequences of instructions that computers execute. So what exactly is it?

syntax

Coding, or programming, is a way of writing instructions for computers that bridges the gap between how humans like to express themselves and how computers actually work.

Programming languages, of which there are hundreds, cannot generally be executed by computers directly. Instead, programs written in a particular "high level" language such as C++, Python or Java are translated by a special piece of software (a compiler or an interpreter) into low-level instructions which a computer can actually run. In some cases programmers write software in low-level instructions directly, but this is fiddly. It is usually much easier to use a high-level programming language, because such languages make it

to write
Python!

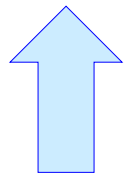
to run!

Assembly!!

| Instruction | Description | Aliases |
|--------------------------------------|--|---------------|
| System instructions | | |
| halt | Stop! | |
| read rX | Place user input in register rX | |
| write rX | Print contents of register rX | |
| nop | Do nothing | |
| Setting register data | | |
| setn rX N | Set register rX equal to the integer N (-128 to +127) | |
| addn rX N | Add integer N (-128 to 127) to register rX | |
| copy rX rY | Set rX = rY | mov |
| Arithmetic | | |
| add rX rY rZ | Set rX = rY + rZ | |
| sub rX rY rZ | Set rX = rY - rZ | |
| neg rX rY | Set rX = -rY | |
| mul rX rY rZ | Set rX = rY * rZ | |
| div rX rY rZ | Set rX = rY / rZ (integer division; no remainder) | |
| mod rX rY rZ | Set rX = rY % rZ (returns the remainder of integer division) | |
| Jumps! | | |
| jumpn N | Set program counter to address N | |
| jumpr rX | Set program counter to address in rX | jump |
| jeqzn rX N | If rX == 0, then jump to line N | jeqz |
| jnezn rX N | If rX != 0, then jump to line N | jnez |
| jgtzn rX N | If rX > 0, then jump to line N | jgtz |
| jltzn rX N | If rX < 0, then jump to line N | jltz |
| calln rX N | Copy the next address into rX and then jump to mem. addr. N | call |
| Interacting with memory (RAM) | | |
| pushr rX rY | Store contents of register rX onto stack pointed to by reg. rY | |
| popr rX rY | Load contents of register rX from stack pointed to by reg. rY | |
| loadn rX N | Load register rX with the contents of memory address N | |
| storen rX N | Store contents of register rX into memory address N | |
| loadr rX rY | Load register rX with data from the address location held in reg. rY | loadi, load |
| storer rX rY | Store contents of register rX into memory address held in reg. rY | storei, store |

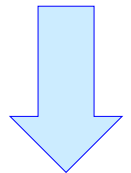
Hmmm
the complete reference

At www.cs.hmc.edu/~cs5grad/cs5/hmmm/documentation/documentation.html



Today

Thursday



Assembly Language

ought to be called *register* language



`read r1`

reads from keyboard into **reg r1**

`write r2`

outputs **reg r2** onto the screen

`setn r1 42`

`reg1 = 42`

you can replace 42 with anything from -128 to 127

`addn r1 -1`

`reg1 = reg1 - 1`

a shortcut

←
This is why assignment is written R to L in Python!

`add r3 r1 r2`

`reg3 = reg1 + reg2`

`sub r3 r1 r2`

`reg3 = reg1 - reg2`

`mul r2 r1 r1`

`reg2 = reg1 * reg1`

`div r1 r1 r2`

`reg1 = reg1 / reg2`

ints only!

Names(s): _____

Try it!

CPU

central processing unit

r1

100

General-purpose register r1

r2

7

General-purpose register r2

r3

42

~~43~~

~~50~~

General-purpose register r3

r4

2

General-purpose register r4

Hmmm...!?

Extra! Change only the instruction on line 4 to create the overall output of 56 or 349 or 0 or 6 ... ?

"Quiz"

screen

100 (input)

(output)

RAM

random access memory

Python

0

read r1

100

r1 = 100

1

setn r2 7

r2 = 7

2

mod r4 r1 r2

r4 = r1 % r2

3

div r3 r1 r4

r3 = r1 // r4

4

sub r3 r3 r2

r3 = r3 - r2

5

addn r3 -1

r3 = r3 + -1

6

write r3

print r3

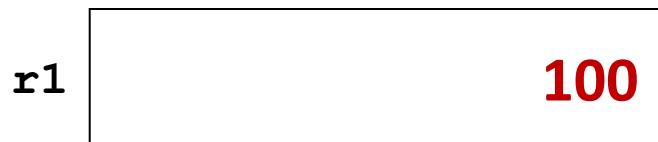
7

halt

Try this on the back page first!

CPU

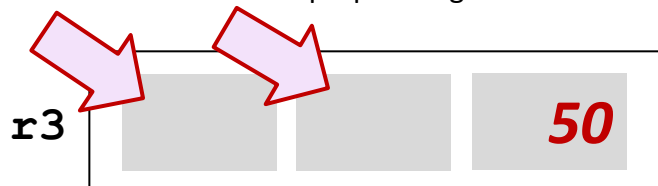
central processing unit



General-purpose register r1



General-purpose register r2



General-purpose register r3



General-purpose register r4

Hmmm...!?

Extra! Change the instruction on line 4 to create the overall output of 56 or 349 or 0 or 6 ... ?

mod

0

div

6

mul

349

screen

100 (input)

42 (output)

Quiz

RAM

random access memory

Python

| | | | |
|---|----------------------------------|---|----------------------------|
| 0 | <code>read r1</code> | 100 | <code>r1 = 100</code> |
| 1 | <code>setn r2 7</code> | | <code>r2 = 7</code> |
| 2 | <code>mod <u>r4</u> r1 r2</code> | | <code>r4 = r1 % r2</code> |
| 3 | <code>div r3 r1 r4</code> | | <code>r3 = r1 // r4</code> |
| 4 | <code>sub r3 r3 r2</code> | | <code>r3 = r3 - r2</code> |
| 5 | <code>addn r3 -1</code> | | <code>r3 = r3 + -1</code> |
| 6 | <code>write r3</code> | | <code>print r3</code> |
| 7 | <code>halt</code> | | |

Is this all we need?

What's
missing
here?

0 **read r1**

1 **mul r2 r1 r1**

2 **add r2 r2 r1**

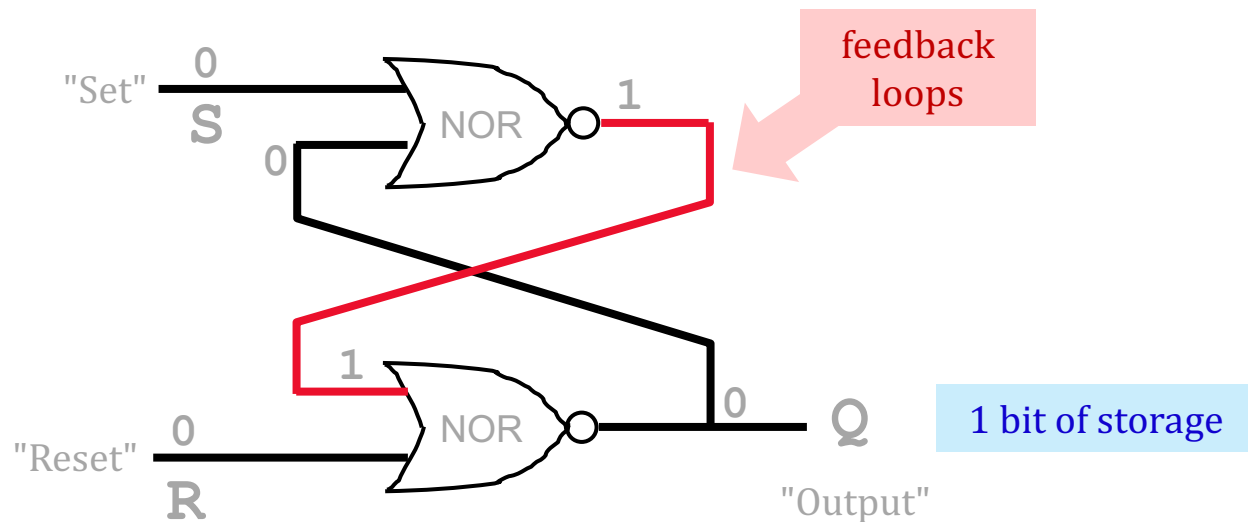
3 **write r2**

4 **halt**



Why *couldn't* we implement Python using only our Hmmm assembly language up to this point?

For systems, innovation is adding an edge to *create a cycle*, not just an additional node.



Loops and ifs

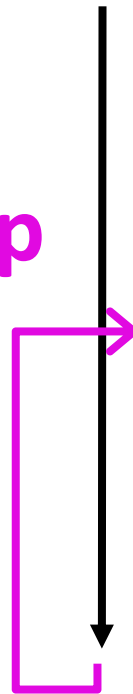
We *couldn't* implement Python using Hmmm so far...

It's too linear!



jumpn!

loop



"straight-line code"

0 **setn r1 42**

1 **write r1**

2 **addn r1 1**

3 **jumpn 1**

4 **halt**

CPU

central processing unit

r1



General-purpose register r1

r2



General-purpose register r2

Screen



jumpn!

RAM

random access memory

0

setn r1 42

1

write r1

2

addn r1 1

3

jumpn 1

4

halt

CPU

central processing unit

r1

... ~~45~~ ~~44~~ ~~43~~ ~~42~~

General-purpose register r1

r2

not used in this program...

General-purpose register r2

42

Screen

43

44

45

46

47

...

crash!

RAM

random access memory

0

setn r1 42

1

write r1

2

addn r1 1

3

jumpn 1

4

halt

← if we jumpn 1

What would happen *IF...*

- we replace line 3 with jumpn 0
- we replace line 3 with jumpn 2
- we replace line 3 with jumpn 3
- we replace line 3 with jumpn 4

CPU

central processing unit

r1



General-purpose register r1

r2



General-purpose register r2

RAM

random access memory

0

setn r1 42

1

write r1

2

addn r1 1

3

jumpn 1 ➡

4

halt

Screen

42

43

44

45

46

...

crash!

Screen

42

no
crash

Screen

42



no
crash

Screen

42



crash!

Screen

42

42

42

42

42

...

no
crash



*What would happen **IF**...*

- we replace line 3 with **jumpn 0**
- we replace line 3 with **jumpn 2**
- we replace line 3 with **jumpn 3**
- we replace line 3 with **jumpn 4**

CPU

central processing unit



General-purpose register r1



jumpn answers

jumpn 4

Screen

42

no
crash

jumpn 3

Screen

42



no
crash

jumpn 2

Screen

42



crash!

jumpn 0

Screen

42

42

42

42

42

... no
crash

RAM

random access memory

0 setn r1 42

1 write r1

2 addn r1 1

3 jumpn 1 →

4 halt

Screen

42

43

44

45

46

...

crash!

← What would happen **IF**...

- we replace line 3 with jumpn 0
- we replace line 3 with jumpn 2
- we replace line 3 with jumpn 3
- we replace line 3 with jumpn 4

Jumps in Hmmm

Conditional jumps

| | | | |
|--------------|-----------|-----------|--|
| jeqzn | r1 | 42 | IF r1 == 0 THEN jump to line number 42 |
| jgtzn | r1 | 42 | IF r1 > 0 THEN jump to line number 42 |
| jltzn | r1 | 42 | IF r1 < 0 THEN jump to line number 42 |
| jnezn | r1 | 42 | IF r1 != 0 THEN jump to line number 42 |

This is making me
jumpy!



Unconditional jump

jumpn **42** Jump to program line # **42**

Jumps in Hmmm

Conditional jumps

| | | |
|--------------|--|------------------------------------|
| jeqzr | if <u>e</u> qual to <u>z</u> ero... | THEN jump to line number 42 |
| jgtzr | if <u>g</u> reater <u>t</u> han <u>z</u> ero ... | EN jump to line number 42 |
| jltzr | if <u>l</u> ess <u>t</u> han <u>z</u> ero... | THEN jump to line number 42 |
| jnezr | if <u>n</u> ot <u>e</u> qual to <u>z</u> ero ... | HEN jump to line number 42 |

Mnemonics!

This is making me
jumpy!



Unconditional jump

jumpn 42

Jump to program line # **42**

| Instruction | Description | Aliases |
|--------------------------------------|--|---------------|
| System instructions | | |
| halt | Stop! | |
| read rX | Place user input in register rX | |
| write rX | Print contents of register rX | |
| nop | Do nothing | |
| Setting register data | | |
| setn rX N | Set register rX equal to the integer N (-128 to +127) | |
| addn rX N | Add integer N (-128 to 127) to register rX | |
| copy rX rY | Set rX = rY | mov |
| Arithmetic | | |
| add rX rY rZ | Set rX = rY + rZ | |
| sub rX rY rZ | Set rX = rY - rZ | |
| neg rX rY | Set rX = -rY | |
| mul rX rY rZ | Set rX = rY * rZ | |
| div rX rY rZ | Set rX = rY / rZ (integer division; no remainder) | |
| mod rX rY rZ | Set rX = rY % rZ (returns the remainder of integer division) | |
| Jumps! | | |
| jumpn N | Set program counter to address N | |
| jumpr rX | Set program counter to address in rX | jump |
| jeqzn rX N | If rX == 0, then jump to line N | jeqz |
| jnezn rX N | If rX != 0, then jump to line N | jnez |
| jgtzn rX N | If rX > 0, then jump to line N | jgtz |
| jltzn rX N | If rX < 0, then jump to line N | jltz |
| calln rX N | Copy the next address into rX and then jump to mem. addr. N | call |
| Interacting with memory (RAM) | | |
| pushr rX rY | Store contents of register rX onto stack pointed to by reg. rY | |
| popr rX rY | Load contents of register rX from stack pointed to by reg. rY | |
| loadn rX N | Load register rX with the contents of memory address N | |
| storen rX N | Store contents of register rX into memory address N | |
| loadr rX rY | Load register rX with data from the address location held in reg. rY | loadi, load |
| storer rX rY | Store contents of register rX into memory address held in reg. rY | storei, store |

Hmmm
the complete reference

At www.cs.hmc.edu/~cs5grad/cs5/hmmm/documentation/documentation.html

Jumps!

off-processor
 access...
 [Thursday]



What Python f'n is this?

CPU

central processing unit

r1

| | |
|--|--|
| | |
|--|--|

General-purpose register r1

r2

| | |
|--|--|
| | |
|--|--|

General-purpose register r2

Screen

-6 (input)

| | |
|--|--|
| | |
|--|--|

RAM

random access memory

| | |
|---|---------------------------|
| 0 | <code>read r1</code> |
| 1 | <code>jgtzn r1 7</code> |
| 2 | <code>setn r2 -1</code> |
| 3 | <code>mul r1 r1 r2</code> |
| 4 | <code>nop</code> |
| 5 | <code>nop</code> |
| 6 | <code>nop</code> |
| 7 | <code>write r1</code> |
| 8 | <code>halt</code> |

space for
future
expansion!

With an input of **-6**, what does this code write out?

Try it!

I think this language has injured my *craniuhmmm!*



1

Follow this Hmmm program.

First run: use **r1 = 42** and **r2 = 5**.

Next run: use **r1 = 5** and **r2 = 42**.

Registers - CPU

| | Run 1 | Run 2 |
|----|----------|----------|
| r1 | 42 | 5 |
| r2 | 5 | 42 |
| r3 | | |
| | ↓ | ↓ |
| | Output 1 | Output 2 |

Memory - RAM

| | |
|---|--------------|
| 0 | read r1 |
| 1 | read r2 |
| 2 | sub r3 r1 r2 |
| 3 | nop |
| 4 | jgtzn r3 7 |
| 5 | write r1 |
| 6 | jumpn 8 |
| 7 | write r2 |
| 8 | halt |

(1) What **common function** does this compute?

Hint: try the inputs in both orders...

(2) **Extra!** How could you change only line 3 so that, if inputs **r1** and **r2** are **equal**, the program will ask for new inputs?

2

Write an assembly-language program that reads a positive integer into **r1**. The program should compute the **factorial** of the input in **r2**. Once it's computed, it should write out that factorial. Two lines are provided:

Registers - CPU

| | | |
|----|-----------------------|---|
| r1 | input | 5 |
| r2 | result – so far | 1 |
| r3 | | |
| | not needed; OK to use | |

Memory - RAM

| | |
|---|-----------|
| 0 | read r1 |
| 1 | setn r2 1 |
| 2 | |
| 3 | |
| 4 | |
| 5 | |
| 6 | |
| 7 | |
| 8 | write r2 |
| 9 | halt |

Hint: On line 2, could you write a test that checks if the factorial is finished; if it's not, compute one piece and then jump back!

Extra! How few lines can you use here? (Fill the rest with **nops**...)