

Overview



Patch Peer Review

Drivers

Low-Level Mid-Level Upper-Level

Unusual Devices



Multiple execvs run out of memory

splhigh prevents interrupts in file operations

Negative process exit codes don't work



ents on a Few Commen

free_kpages is a no-op in dunbym.c

#define foo(x) do {...} while (0)
is weird, obscure, and norstandard

Multiple execvs run out of memory

free_kpages is a no-op in dumbvm.c

splhigh prevents interrupts in file operations

Negative process exit codes don't work



mments on a Few Comments

Multiple execvs run out of memory free_kpages is a no-op in dunbym.c

splhigh prevents interrupts in file operations
mi_switch requires high spl—and drops it

Negative process exit codes don't work

#define foo(x) do {...} while (0)
is weird, obscure, and nonstandard

Multiple execvs run out of memory

free_kpages **is a no-op in** dumbvm.c

splhigh prevents interrupts in file operations

mi_switch requires high spl—and drops it

Negative process exit codes don't work



mments on a Few Comments

Multiple execvs run out of memory free_kpages is a no-op in dumbrm.c

splhigh prevents interrupts in file operations mi_switch requires high spl—and drops it

Negative process exit codes don't work Exit codes are 8 bits, unsigned

#define foo(x) do {...} while (0)
is weird, obscure, and nonstandard

Multiple execvs run out of memory

free_kpages is a no-op in dumbvm.c

splhigh prevents interrupts in file operations

mi_switch requires high spl—and drops it

Negative process exit codes don't work

Exit codes are 8 bits, unsigned



mments on a Few Comments

Multiple execvs run out of memory free_kpages is a no-op in durbor.c

splhigh prevents interrupts in file operations mi_switch requires high spl—and drops it

Negative process exit codes don't work Exit codes are 8 bits, unsigned

Multiple execvs run out of memory

free_kpages is a no-op in dumbvm.c

${\tt splhigh} \ {\tt prevents} \ {\tt interrupts} \ {\tt in} \ {\tt file} \ {\tt operations}$

mi_switch requires high spl—and drops it

Negative process exit codes don't work Exit codes are 8 bits. unsigned

#define foo(x) do {...} while (0)
is weird, obscure, and nonstandard

It's a standard C idiom because...

Patch Peer Review

Numeric Evaluations

Group	Clarity	Concise	Fit	Correct	Docs	Total
hertz	4.78	4.67	4.67	4.00	4.78	22.90
fewer	4.67	4.33	4.67	2.67	4.33	20.67
dwarf	4.33	4.33	5.00	2.67	4.00	20.33
yank	5.00	4.33	4.33	2.33	4.00	19.99
radon	3.67	4.33	3.67	4.00	2.67	18.34
race	4.00	4.33	4.67	2.00	3.33	18.33
wilt	4.33	4.33	4.00	1.33	3.33	17.32
bides	3.67	3.00	3.00	2.67	4.00	16.34
cue	3.67	4.00	2.67	1.67	4.00	16.01
diner	4.33	4.67	4.00	1.33	1.33	15.66



Group	Clarity	Concise	Fit	Correct	Docs	Tota
hertz	4.78	4.67	4.87	4.00	4.78	22.9
fewer	4.67	4.33	4.67	2.67	4.33	20.6
dwarf	4.33	4.33	5.00	2.67	4.00	20.3
vank	5.00	4.33	4.33	2.33	4.00	19.95
nadon	3.67	4.33	3.67	4.00	2.67	18.3
race	4.00	4.33	4.67	2.00	3.33	18.3
wit	4.33	4.33	4.00	1.33	3.33	17.3
bides	3.67	3.00	3.00	2.67	4.00	16.3
cue	3.67	4.00	2.67	1.67	4.00	16.0
diner	4.33	4.67	4.00	1.33	1.33	15.68

Ranking



Rank	Group		
1.00	hertz		
2.00	radon		
2.00	yank		
2.33	fewer		
3.00	bides		
3.00	race		
3.33	dwarf		
3.67	cue		
3.67	diner		
4.00	wilt		

Device Drivers



Device-access code can run in kernel or user mode (but usually kernel).

Driver must abstract control registers to OS's read/write model:

- Validate request
- Wait for idle
- Issue commands through control registers
- Possibly block waiting for interrupt
- Possibly invoke scheduler



Many devices have common characteristics; e.g., different brands of disk or printer

Makes sense to abstract common parts

Resulting structure is uniform driver sitting above specific one



Desirable to collect input before delivering it, accept output before device swallows it

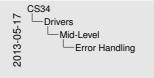
Kernel buffers allow both features

Wise to have extra buffers to allow overlapped I/O

Many devices need buffers, so common kernel mechanism makes sense



Drivers Mid-Level



Error Handling

Best option on errors: retry and hide from upper levels Atternative: return error code to application & let it handle Worst option: ask user what to do (user usually has insufficien information to make wise decision)

Best option on errors: retry and hide from upper levels

Alternative: return error code to application & let it handle

Worst option: ask user what to do (user usually has insufficient information to make wise decision)

Abstractions



Some devices need more than just read and write:

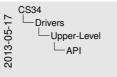
- Disks need filesystems
- Network cards need routing and connection management
- Graphics displays need windowing
- Keyboard needs editing
- Mouse needs pointing to particular windows

▶ ...

OS must provide sensible interposition/interface

Drivers Upper-Leve

API



User-space applications need standardized interface • Open, close, read, write, issek • What to do shoct unusual closes like reject CD'? Sometimes need even higher-level abstractions • Mountumnoute • Printer spooling

User-space applications need standardized interface

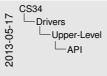
- Open, close, read, write, lseek
- What to do about unusual cases like "eject CD"?

Sometimes need even higher-level abstractions

- Mount/unmount
- Printer spooling

Drivers Upper-Leve

API



User-space applications need standardized interface • Open, close, read, write, leaek • What to do about unsult classes like "ight CD"? Social Somstimes need even higher-level abstractions • Mountlummount • Printer spooling

User-space applications need standardized interface

- Open, close, read, write, lseek
- What to do about unusual cases like "eject CD"? ioctl

Sometimes need even higher-level abstractions

- Mount/unmount
- Printer spooling

Drivers Upper-Level

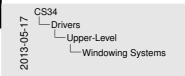
Windowing Systems

Modern GUIs need window management:

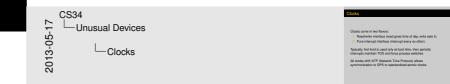
- Overlapping windows
- High-performance drawing
- "Events" (keystrokes and mouse clicks) delivered to selected windows
- Window manager to decide which window is on top and which is active

In Unix, all of this is implemented as a network-connected server that runs the display, mouse, and keyboard: the **X Window System**

Applications are *clients* that connect to the server and ask for windows to be drawn, keystrokes delivered, etc.



٧	Mindowing Systems
	Modern GUIs need window management:
	 Overlapping windows
	 High-performance drawing
	 "Events" (keystrokes and mouse clicks) delivered to selected windows
	 Window manager to decide which window is on top and which is active
	In Unix, all of this is implemented as a network-connected server that runs the display, mouse, and keyboard: the X Window System
	 Applications are clients that connect to the server and ask for windows to be drawn, levstrokes delivered, etc.



Clocks come in two flavors:

Clocks

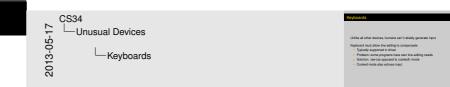
- 1. Read/write interface (read gives time of day, write sets it)
- 2. Pure interrupt interface (interrupt every so often)

Typically, first kind is used only at boot time, then periodic interrupts maintain TOD and force process switches

All clocks *drift*; NTP (Network Time Protocol) allows synchronization to GPS or standardized atomic clocks

Unusual Devices

Keyboards



Unlike all other devices, humans can't reliably generate input

Keyboard must allow *line editing* to compensate

- Typically supported in driver
- Problem: some programs have own line-editing needs
- Solution: raw (as opposed to cooked!) mode
- Cooked mode also echoes input



Mouse generates periodic updates: Δx , Δy , buttons

Problem: to whom do mouse events go?

 Unusual Devices
 CS34
 More

 Unusual Devices
 Unusual Devices
 More greaters profile optications A dy bothoms

 Profile
 Mice
 More greaters profile optications A dy bothoms

Mouse generates periodic updates: Δx , Δy , buttons

Problem: to whom do mouse events go?

Solution: Send to windowing system, let it decide which window is interested