

CS140: Algorithms

Z Sweedyk
Lecture 17
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1

NP-completeness

Problem A is NP-Complete if

- A is in NP
- A is NP-hard

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2

Last Time

- Reduction by Restriction

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3

New Problems

- Minimum Cover
- Hitting Set
- Subgraph Isomorphism
- Bounded Degree Spanning Tree
- Knapsack
- Multiprocessor Scheduling

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4

Some Basics

- 3SAT
- 3DM
- X3C
- Partition
- HC
- HP
- Longest Path
- VC
- Clique
- Independent Set
- 3 Colorability
- 4 Colorability
- K Colorability

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5

Minimum Cover

- Input: Collection C of subsets of the set S and an integer K
- Question: Does C contain K or fewer sets whose union is S ?
- Reduction from Vertex Cover $G=(V,E), K'$
- $S=E, C=\text{Each adjacency list is a subset, } K=K'$

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6

Hitting Set

- Input: Collection C of subsets of a set S and an integer K
- Question: Is there a subset $W \subseteq S$ containing K or fewer elements such that every set in C has an element in W ?
- Reduction from Vertex Cover $G=(V,E),K'$
- $S=V, C=E, K=K'$

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7

Subgraph Isomorphism

- Input: Graphs G and H
- Question: Does G contain a subgraph isomorphic to H ?
- Reduction from Clique: G', K
- $G'=G, H$ is a fully connected graph on K vertices

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8

Bounded Degree Spanning Tree

- Input: Graph G and integer K
- Question: Does G have a spanning tree in which each vertex is incident to K or fewer edges.
- Reduction from Hamiltonian Path: G'
- $G=G; K=2$

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9

Knapsack

- Input: A set of "items" each having an integer "weight" w and an integer "value" v , plus integers B and K
- Question: Is there a subset of items that weighs at most B and has value at least K ?
- Reduction from Partition $S=\{s_1, \dots, s_n\}$
- Item i has $\text{weight}=\text{value}=s_i, B=K=(\sum s_i)/2$

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10

Multiprocessor Scheduling

- Input: A set of computational tasks where task i requires t_i seconds to process, an integer K representing the number of processors available, and an integer D representing the amount of seconds available for processing.
- Question: Can the tasks be scheduled on the processors so that all finish in D seconds?
(A task must be completed on a single processor without interruption. You can assume t_i is an integer.)

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11

Multiprocessor Scheduling

- Input: A set of computational tasks where task i requires t_i seconds to process, K (number processors), D (processing time)
- Question: Can the tasks be scheduled on the processors so that all finish in D seconds?
- Reduction from Partition: $S=(s_1, \dots, s_n)$
- Task i takes time $s_i, K=2, D=(\sum s_i)/2$

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12

Set Packing

- Input: Collection C of sets and an integer K
- Question: Does C contain K or more disjoint sets?
- Reduction from X3C: S, C' where $\|S\|=3q$
- $C'=C, K=q$

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13

Some New (and Harder) Reductions

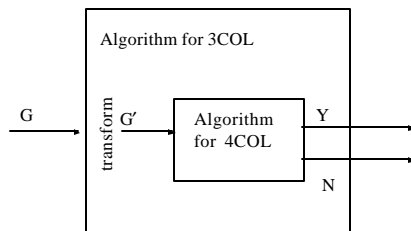
- 4 COL
- Partition into Triangles
- Minimum Test Collection
- Feedback Vertex Set
- Exact Cover by 4-Sets

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14

$3COL \approx_p 4COL$

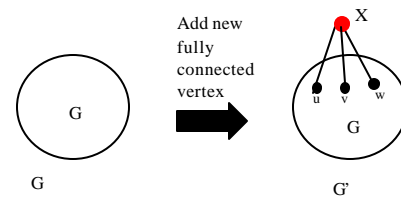


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Transform



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16

Claim

- G is 3 Colorable iff G' is 4 Colorable

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17

Some New (and Harder) Reductions

- | NEW | OLD |
|----------------------------|----------------|
| • Exact Cover by 4-Sets | • Vertex Cover |
| • Partition into Triangles | • 3DM |
| • Minimum Test Collection | • 3XC |
| • Feedback Vertex Set | |

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Exact Cover by 4-Sets

- Input: Collection C of subsets of a set S . Each set in C has cardinality 4.
- Question: Are there $\lceil |S|/4 \rceil$ sets in C that partition S ?

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19

Partition into Triangles

- Input: Graph G
- Question: Can the vertices of G be partitioned into sets of size three so that each set is fully connected in G ?

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20

Feedback Arc Set

- Input: Digraph G and integer K
- Question: Is there a subset W of K or fewer vertices of G such that every (directed) cycle of G passes through at least one vertex of W .

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21