

## Higher-Order Functions

CS 131: Programming Languages  
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## Cons vs. Append

- There are two binary operators used with lists
  - The `::` operator, called cons
  - The `@` operator, called append

- Note:

```
[1,2,3] :: [] = [[1,2,3]]  
[1,2,3] @ [] = [1,2,3]
```

- Also, only `::` can appear in patterns.

## Cons vs. Append

- The type of `@` is:
  
- The type of `::` is:

## Applying a Function to a List

- Problem: Apply some function  $f$  to every element of a list, return the list of results
  - That is, given the input

```
[x1, ..., xn]
```

```
return
```

```
[f(x1), ..., f(xn)].
```

```
fun loop [] =  
  | loop (x::xs) =
```

## Applying a Function to a List

- New Problem: Write a function `map` that *given* `f`, returns a function that applies `f` to every element of a list.

```
fun map f =  
  let  
    fun loop []      = []  
      | loop (x::xs) = (f x)::(loop xs)  
  in  
    loop  
  end
```

## What is the type of `map` ?

- The argument can be any function.
- If we assume that `f : 'a -> 'b`, what is the type of the locally defined function `loop`?
- Then what is the type of `map`?

## Doubling Lists

```
val doubler : int list -> int list  
  = map (fn x => x*2)  
val l      = doubler [1,2,3]
```

```
val l = map (fn x => x*2) [1,2,3]
```

```
fun double x = x*2  
val l = map double [1,2,3]
```

## Higher-Order Functions

- A function that takes a function as its argument or returns a function as its result is said to be a *higher-order function*.
  - e.g., `map` is higher-order
- Let's look at some more examples

## Building Functions that Add

- Consider the following functions:

```
fun addone(x) = x + 1
fun addtwo(x) = x + 2
fun addsix(x) = x + 6
```

- Can we generalize this construction?
- Goal: a function that, given *n*, returns the function which adds *n* to its argument

## Building Functions that Add

- It may help to consider the fully-expanded code for the functions on the previous slide:

```
val addone = (fn x => x+1)
val addtwo = (fn x => x+2)
val addsix = (fn x => x+6)
```

- Exercise: Define  
add : int -> (int->int)

## Using add

```
val addone = add 1
val addtwo = add 2
val addsix = add 6

fun increment_list lst =
  map (add 1) lst

val increment_list' =
  map (add 1)
```

## Types for Curried Functions

- The type of add is  
int -> (int -> int)
- Function types are right associative, so can write  
int -> int -> int
- There are two ways to think about this type.
  - The function add takes an int and returns a function on ints  
add 4 : int->int
  - The function add takes two integer arguments  
*in succession*  
add 4 7 : int

## Syntax For Curried Functions

- Functions like `add` that do nothing but return another function are said to be *curried*.
- SML has special syntax for defining curried functions
  - Function argument patterns are separated by spaces

```
fun add n m = n+m

fun map f [] = []
  | map f (x::xs) = (f x)::(map f xs)
```

## Staged Computation

- Consider the following variation of `add`, which adds the square of its argument:

```
fun addsquare 0 m = m
  | addsquare n m = (n*n) + m
```

Equivalently:

```
fun addsquare n m =
  (case n of
   0 => m
  | _ => (n*n)+m)
```

## Staged Computation

- If we now execute the code

```
val f : int->int = addsquare 6
val l = map f [1,2,3,4,5]
```

then how many times do we compare 6 against 0?

- How many times do we compute  $6*6$ ?

## Staged Computation

- More efficient version:

```
fun addsquare n =
  (case n of
   0 => (fn m => m)
  | _ => (fn m => (n*n)+m))
```

## Staged Computation

- Even more efficient:

```
fun addsquare n =  
  (case n of  
    0 => (fn m => m)  
  | _ => let  
        val n2 = n*n  
        in  
          fn m => n2+m  
        end)
```

## Function Composition

- Goal: a function `compose` that, given functions `f` and `g`, returns their composite.
  - Recall: composite of `f` and `g` is the function which given `x`, returns `f(g(x))`.
  - Also, what is the type of this function?

## Functions and Re-binding

- Consider the following definitions:

```
val x = 3  
fun addx (y:int) = y+x
```

- Now, what is the value of `addx 2` ?

## Functions and Re-binding

- Consider the following definitions:

```
val x = 3  
fun addx (y:int) = y+x  
val x = 5
```

- Now, what is the value of `addx 2` ?

## Functions and Re-binding

- Consider the following input:

```
fun add1 x = x+1
fun add2 x = add1(add1(x))
val x = add2 4

fun add1 x = x+3
val y = add2 4
```

- Now, what are the values of x and y ?

## The exists Function

```
(* exists : ('a -> bool) -> 'a list -> bool *)

fun exists p [] = false
  | exists p (x::xs) = (p x) orelse (exists p xs)

fun exists p =
  let fun loop [] = false
      | loop (x::xs) = (p x) orelse (loop xs)
  in
    loop
  end
```

## The all Function

```
(* all : ('a -> bool) -> 'a list -> bool *)

fun all p [] = true
  | all p (x::xs) = (p x) andalso (all p xs)

fun all p =
  let fun loop [] = true
      | loop (x::xs) = (p x) andalso (loop xs)
  in
    loop
  end
```

## The find Function

```
(* find : ('a -> bool) -> 'a list -> 'a option *)

fun find p [] = NONE
  | find p (x::xs) =
    if (p x) then (SOME x) else (find p xs)

fun find p =
  let fun loop [] = NONE
      | loop (x::xs) =
        if (p x) then (SOME x) else (loop xs)
  in
    loop
  end
```

## The partition Function

```
(* partition : ('a -> bool) ->
   'a list -> 'a list * 'a list *)

fun partition p []      = ([], [])
  | partition p (x::xs) =
    let
      val (yes,no) = partition p xs
    in
      if p x then
        (x::yes, no)
      else
        (yes, x::no)
    end
```

## A Question of Style

```
fun map f []      = []
  | map f (x::xs) = (f x) :: (map f xs)

fun map2 (f,[])   = []
  | map2 (f,x::xs) = (f x) :: (map2 (f,xs))

fun map3 []      f = []
  | map3 (x::xs) f = (f x) :: (map3 xs f)

fun map4 ([],f)   = []
  | map4 (x::xs,f) = (f x) :: (map4 (f,xs))
```