

CS140: Algorithms

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Lecture 0: Introduction
9/5/01

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1

Today

- Who Am I?
- Who Are You?
- Course Overview
- Some Basics

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2

Course Overview

1. How to analyze algorithms
2. How to design algorithms
3. NP-completeness

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3

Course Requirements

- Homework
- Exams
- Class participation

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4

Homework

- Assignments will be posted to the web page on Monday and Wednesday.
- Assignments are due at the start of the next class period.
- Solutions will be posted at due time so no late homework will be accepted
- Your three lowest homework scores will be dropped in computing your final homework grade.
- Solutions should be prepared in LaTeX

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5

Exam Due Dates (Tentative Schedule)

- Exam I: Oct. 8
 - Exam II: Nov. 12
 - Exam III: Dec. 3
 - Final: Dec. 18
-
- Exams will be take-home, timed, closed book

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6

Course Requirements – class participation

- Show up to class
- Speak up in class
- Hand in daily “worksheets”

Advice

- Stay on top of things
- Seek help when necessary

The Problem

Computational Problem: Specified by input/output pair

Sorting

- Sorting Integers in Ascending Order (SIAO):
Input: A list of integers
Output: The input integers sorted in ascending order
- Example:
Input: 5,3,8,1,2
Output:

The Algorithm

- Computational Problem: Specified by input/output pair.
- Algorithm: Well-defined sequence of computational steps that produce a correct output for every valid input.

What should an algorithm for SIAO do?

- Example 1:
 - Input: 5,3,8,1,2
 - Output:
- Example 2:
 - Input: 3,a,5.27,mudder
 - Output:

Software Development

The problem:
Huh?



The idea:
A-ha!



The program:
Ta-da!



The algorithm exists somewhere between
a-ha and ta-da.

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13

An Algorithm for SIAO?

Sort1(S)

```
While there are integers x and y in S such that x precedes y
in S and  $x > y$ 
    Swap x and y in S
Return S
```

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14

Sort1 Example

Input: 5, 3, 8, 1, 2

Swap 3 and 2: 5, 2, 8, 1, 3

Swap 5 and 3: 3, 2, 8, 1, 5

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15

Is Sort1 an algorithm?

Is Sort1(S) well-defined?

No! We need to specify a selection rule.

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16

An Algorithm for SIAO?

Sort2(S)

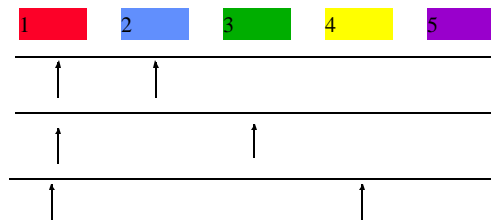
```
Assume a fixed order on pairs of elements in S
While there are integers x and y in S such that x precedes y
in S and  $x > y$ 
    Choose first pair x,y that is out of order
    Swap x and y in S
Return S
```

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17

Order on pairs of indices



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18

Sort2 Example

Input: 5, 3, 8, 1, 2

Swap 5 and 3: 3, 5, 8, 1, 2

Swap 3 and 1: 1, 5, 8, 3, 2

Is Sort2 an algorithm?

Is Sort2(S) well-defined? Yes.

Does it produce the correct output for any valid input?

Proof of correctness

- When the algorithm halts S is sorted.
- The algorithm halts on all input.
 - How can we measure the progress the algorithm makes from a swap?

Claim

- The number of “out-of-order” pairs decreases with each swap.
- Progress: Eventually we’ll have 0 “out-of-order” pairs and the algorithm will halt.

Illustration of claim

Input: 5, 3, 8, 1, 2 #out-of-order pair: 7

Swap 5 and 3: 3, 5, 8, 1, 2 6

Swap 3 and 1: 1, 5, 8, 3, 2 5

Proof of claim:

- Suppose the algorithm chooses to swap x and y:

$u_0, \dots, u_i, \mathbf{x}, v_0, \dots, v_j, \mathbf{y}, w_0, \dots, w_k$

- Change in status of a pair:

- Out-of-Order \rightarrow Ordered
- Ordered \rightarrow Out-of-order

- Pairs to consider

- $u-u, u-x, u-v, u-y, u-w,$
- $\boxed{x-v} \quad \boxed{x-y}, x-w,$
- $v-v, \boxed{v-y}, v-w$
- $y-w, w-w$

Ordered→unordered

- Suppose an x -v pair goes from ordered to out-of-order.
- Then _____

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Ordered→unordered

- Suppose an x -y pair goes from ordered to out-of-order.
- Then _____

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26

Ordered→unordered

- Suppose an v -y pair goes from ordered to out-of-order.
- Then _____

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27

Proof

- The x-y pair goes from out-of-order to ordered.
- For every pair that goes from ordered to out-of-order _____.

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28

Is Sort2 a good algorithm?

- Is it easy to understand?
- Is it easy to implement?
- Is it fast?
- Is it space-efficient?

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29

How fast is Sort2?

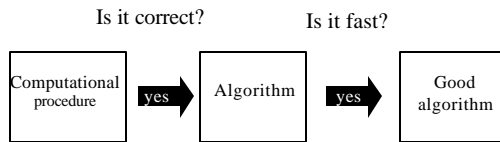
- How many swaps can the algorithm make?
- How many comparisons does the algorithm need to make to find a pair to swap?
- How many comparisons does the algorithm make?
- The running time is _____

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30

CS140: Two questions



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31

Another algorithm for SAIO?

```
Sort3(S)
  If ||S|| ≤ 1
    Return: S
  Else
    Return: Sort3(S[max-element(S)], max-element(S))
```

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32

Proof of correctness

(Prove Def. 1 \approx Def. 2)

- Def. 1 (non-recursive)

s_0, s_1, \dots, s_n is sorted if for every i, j such that $i < j$ it holds that $s_i \leq s_j$

- Def. 2 (recursive)

The list s_0, s_1, \dots, s_n is sorted if

1. $n=0$, or
2. $n>0$ and s_0, s_1, \dots, s_{n-1} is sorted and for all $i < n$ it holds that $s_n \geq s_i$

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33

Example: Sort3(3,1,5,2,4)

```
Sort3(3,1,5,2,4) = Sort3(3,1,2,4),5
                  = Sort3(3,1,2),4,5
                  = Sort3(1,2),3,4,5
                  = Sort3(1),2,3,4,5
                  = 1,2,3,4,5
```

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34

Recursive Algorithms

What about Sort3?

```
Sort3(S)
  If ||S|| ≤ 1
    Return: S
  Else
    Return: Sort3(S[max-element(S)], max-element(S))
```

Let $T(n)$ be the running time of Sort3:

$$T(1) = c_2$$

$$T(n) = c_1 n + T(n-1), n > 1$$

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35

Unwinding

$$\begin{aligned} T(n) &= c_1 n + T(n-1) \\ &= c_1 n + c_1(n-1) + T(n-2) \\ &= c_1 n + c_1(n-1) + c_1(n-2) + T(n-3) \end{aligned}$$

\vdots

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36

Basic skills

- LaTeX – HW0
- Ch 2.2 of CLR