Systolic Arrays and Algorithms

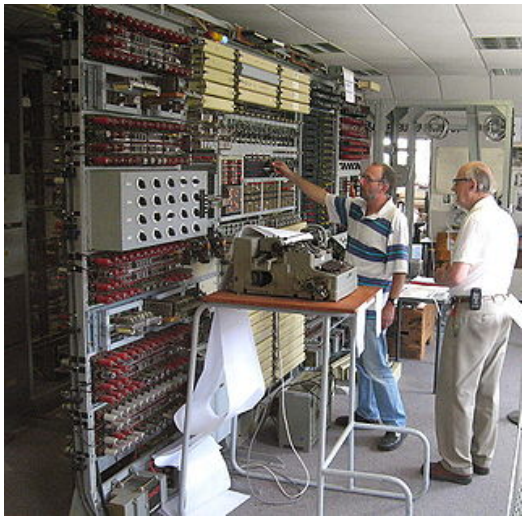
Systolic Arrays

- This is a form of **pipelining**, sometimes in more than one dimension.
- The term “systolic” was first used in this context by H.T. Kung, then at CMU; it refers to the “pumping” action of a heart.
- Machines have been constructed based on this principle, notably the iWARP, fabricated by Intel in the 80’s.

Early Systolic Operation

- Colossus Mark 2 at Bletchley Park, 1944
  (http://en.wikipedia.org/wiki/Colossus_computer)

- “Colossus included the first ever use of shift registers and systolic arrays, enabling five simultaneous tests, each involving up to 100 Boolean calculations, on each of the five channels on the punched tape (although in normal operation only one or two channels were examined in any run).”

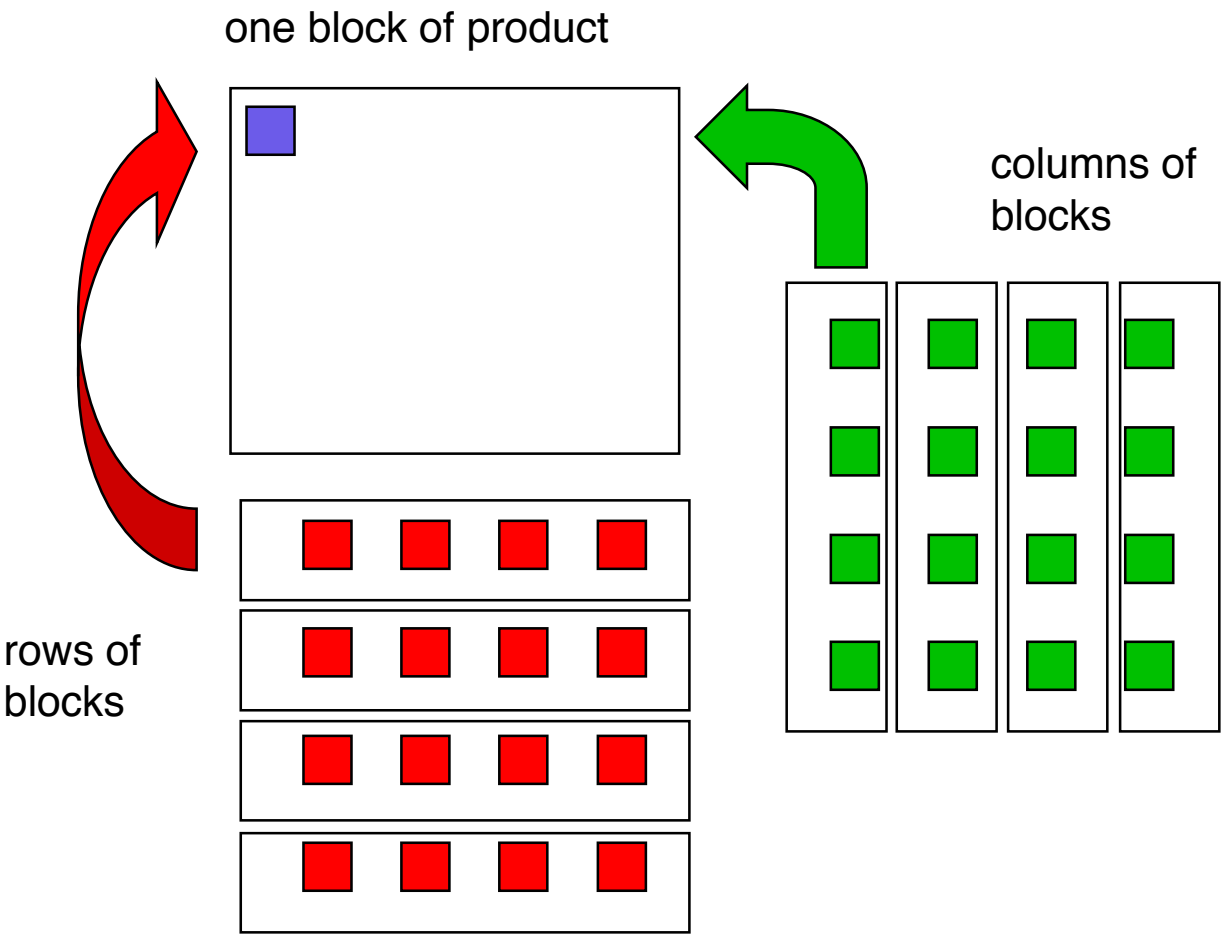


In 1994, a team led by Tony Sale (right) began a reconstruction of a Colossus at Bletchley Park. Here, in 2006, Sale supervises the breaking of an enciphered message with the completed machine.

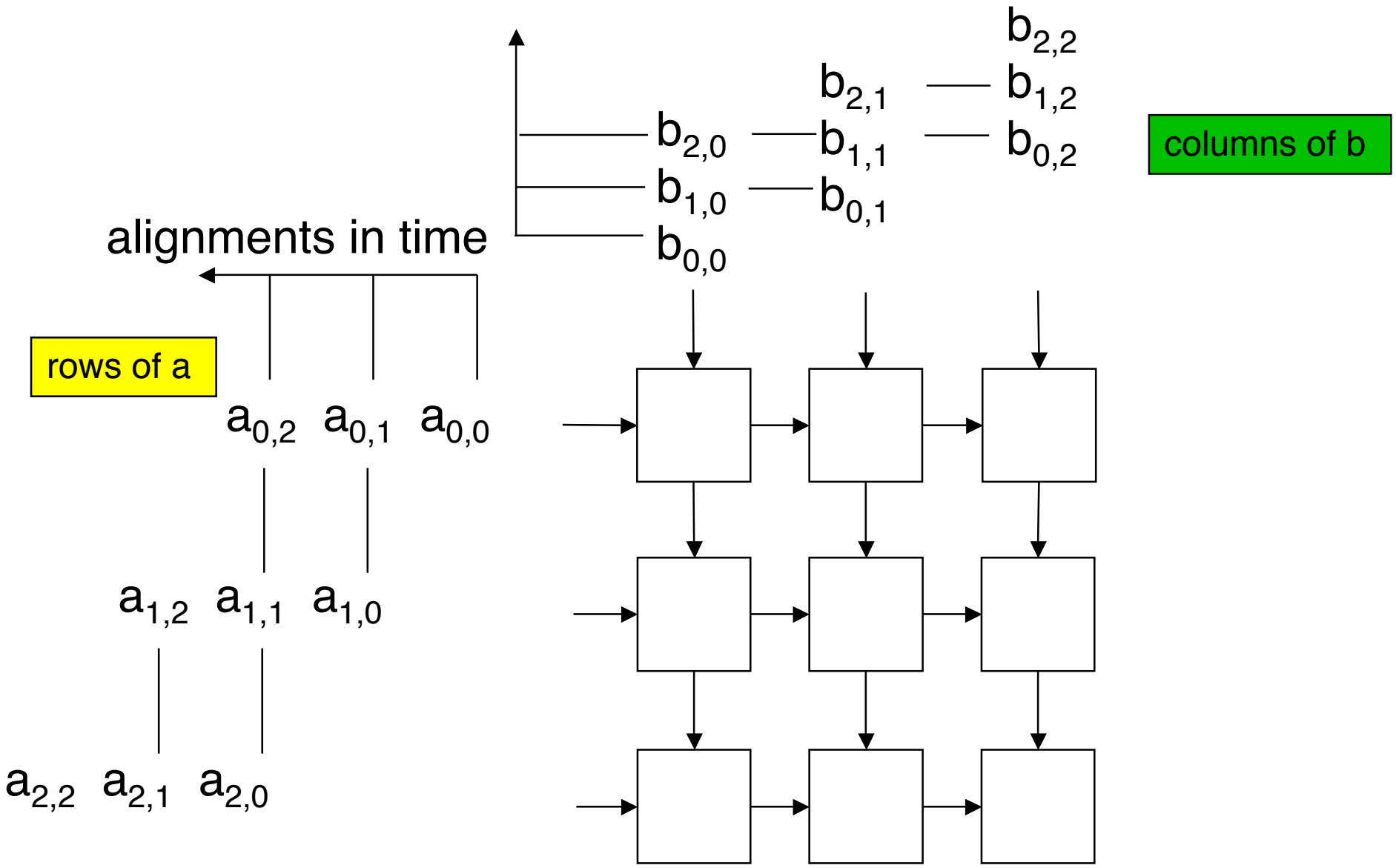
Systolic Matrix Multiplication

- Processors are arranged in a 2-D grid.
- Each processor accumulates one element of the product.
- The elements of the matrices to be multiplied are “pumped through” the array.
- We illustrate for individual elements, but can be applied to **block decompositions**.

Block Multiplication



Systolic Matrix Multiplication
Illustrated with two 3x3 matrices



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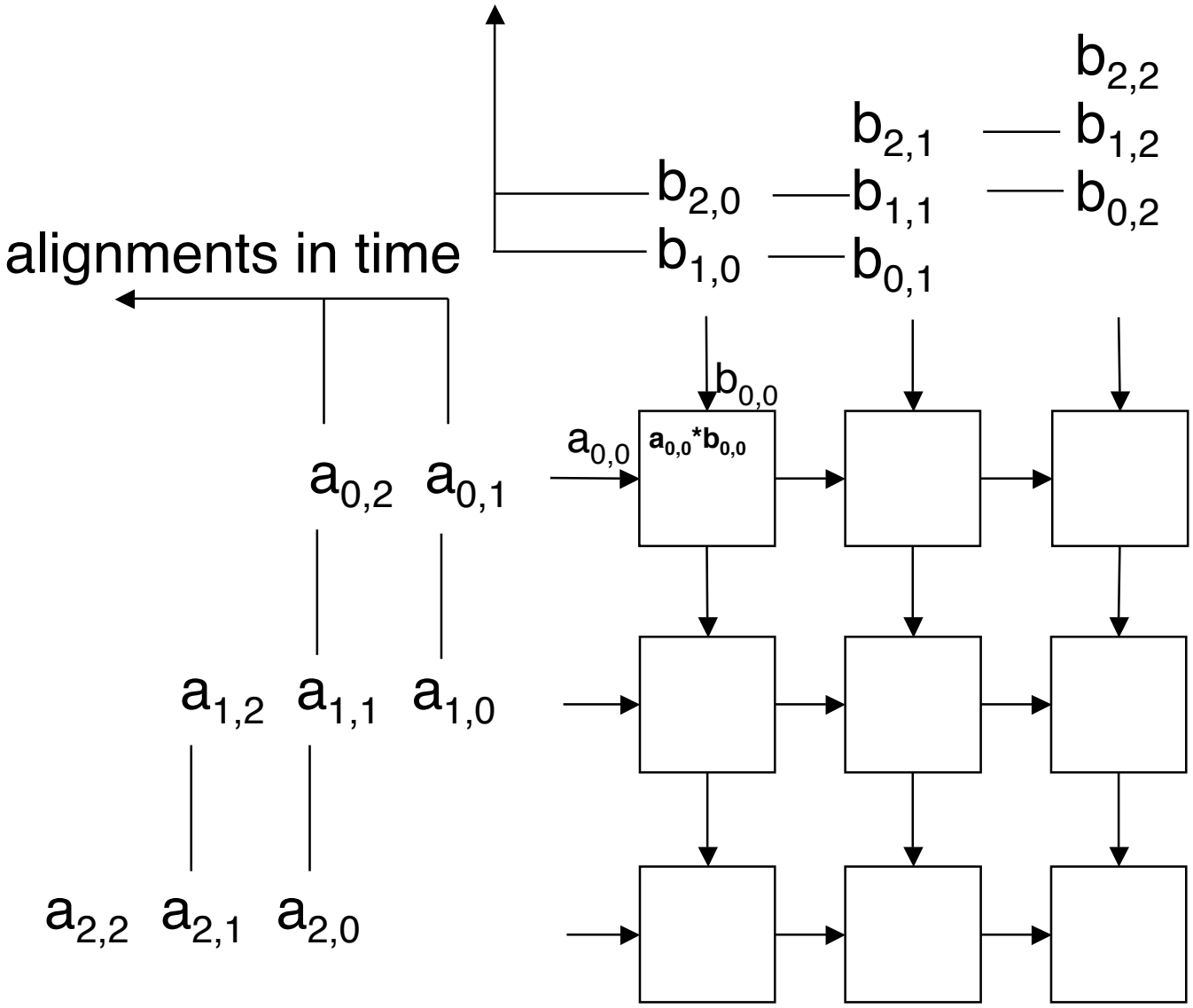

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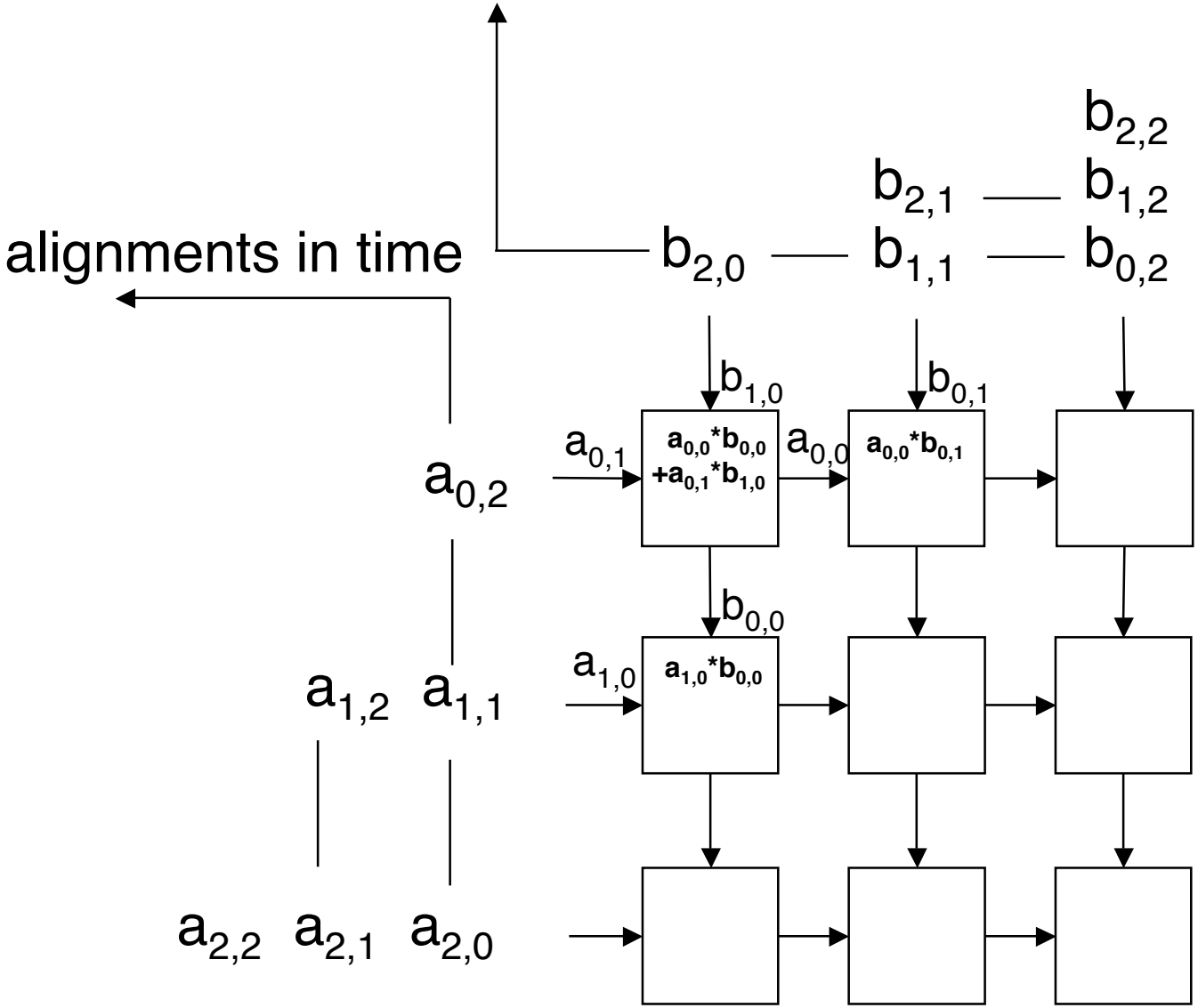

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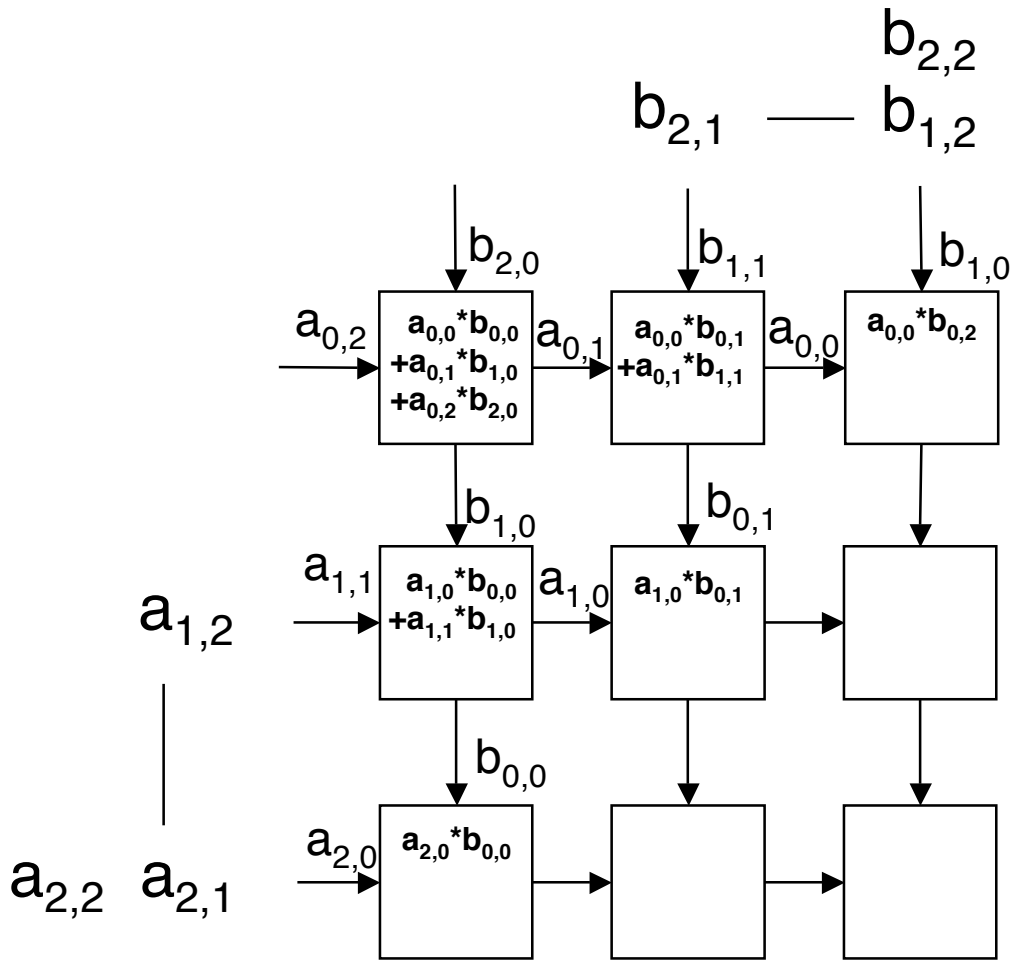

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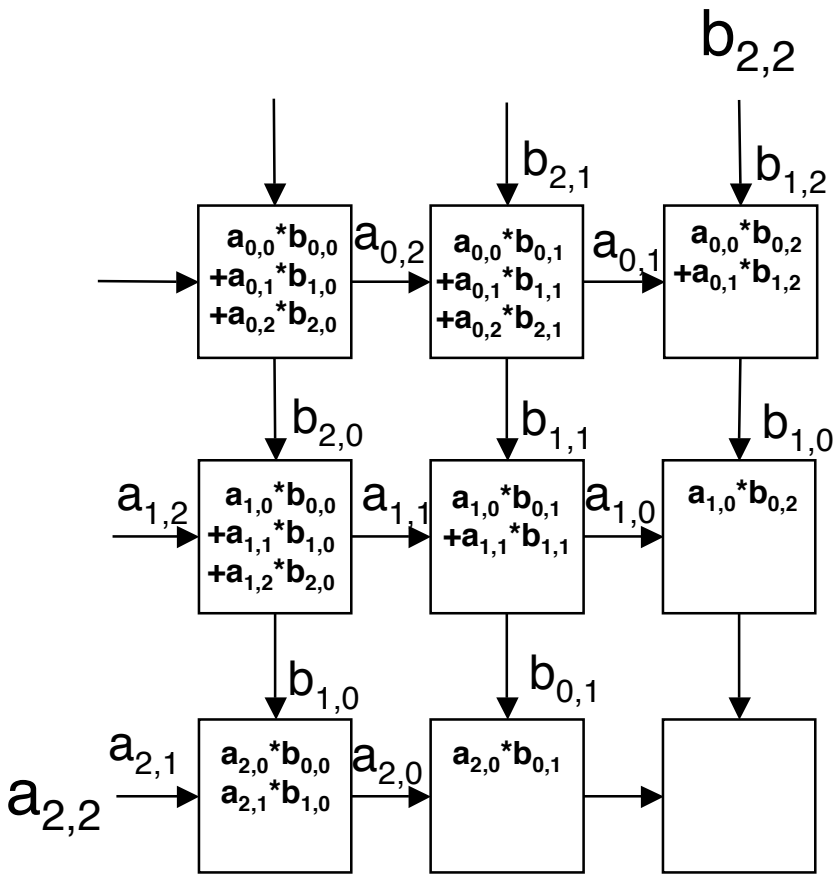

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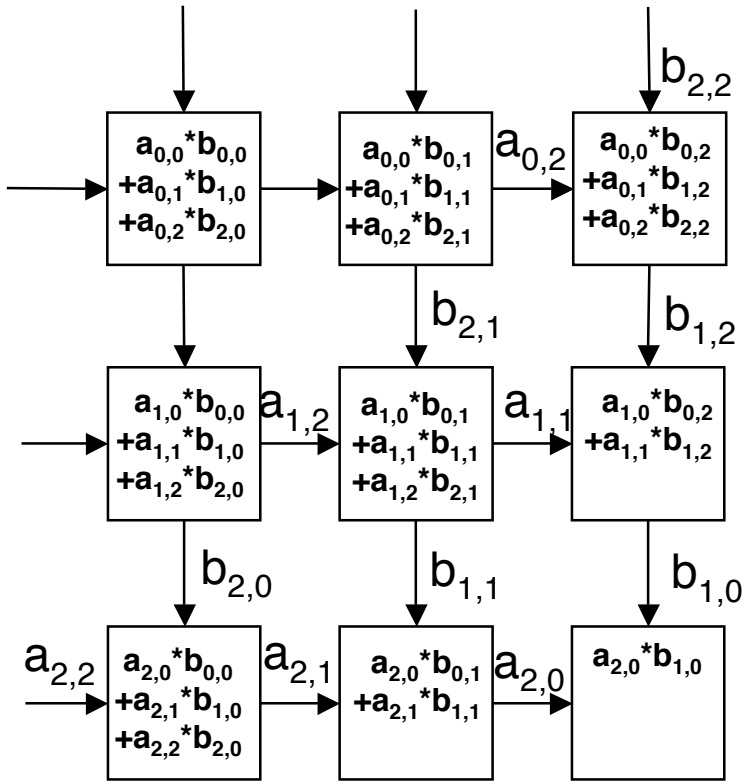

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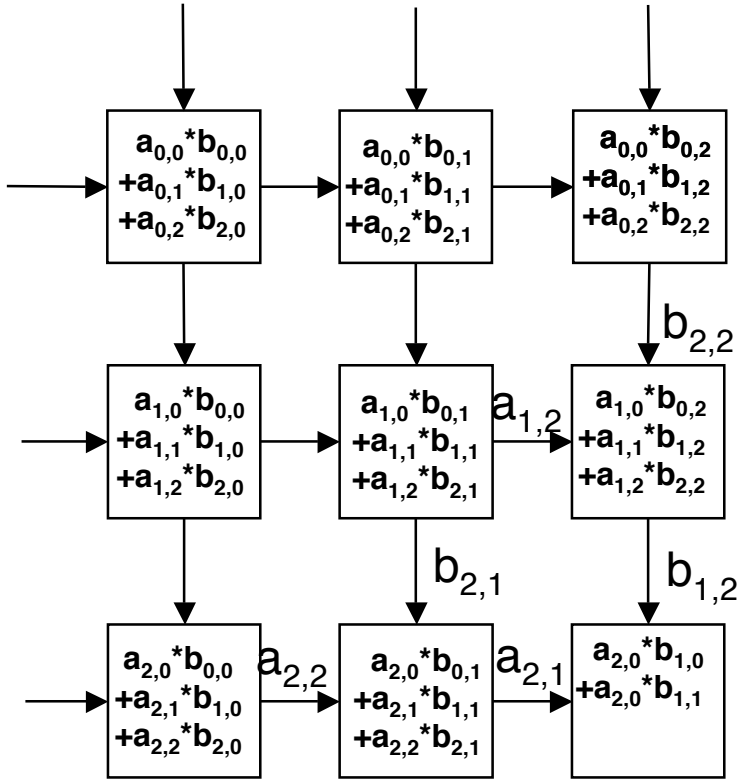

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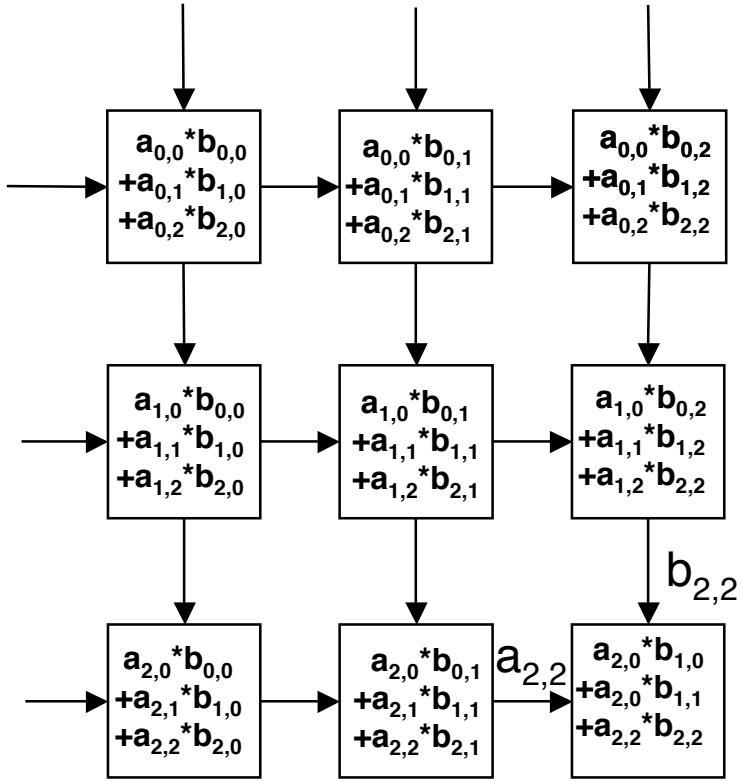


Analysis

- Assuming n x n matrices (n^2 elements)
- $O(n^3)$ operations
- O(n) time steps on n^2 processors

A Related Algorithm: Cannon's Method

- Let's take another view of systolic multiplication: Consider the rows and columns of the matrices to be multiplied as **strips** that slide past each other.

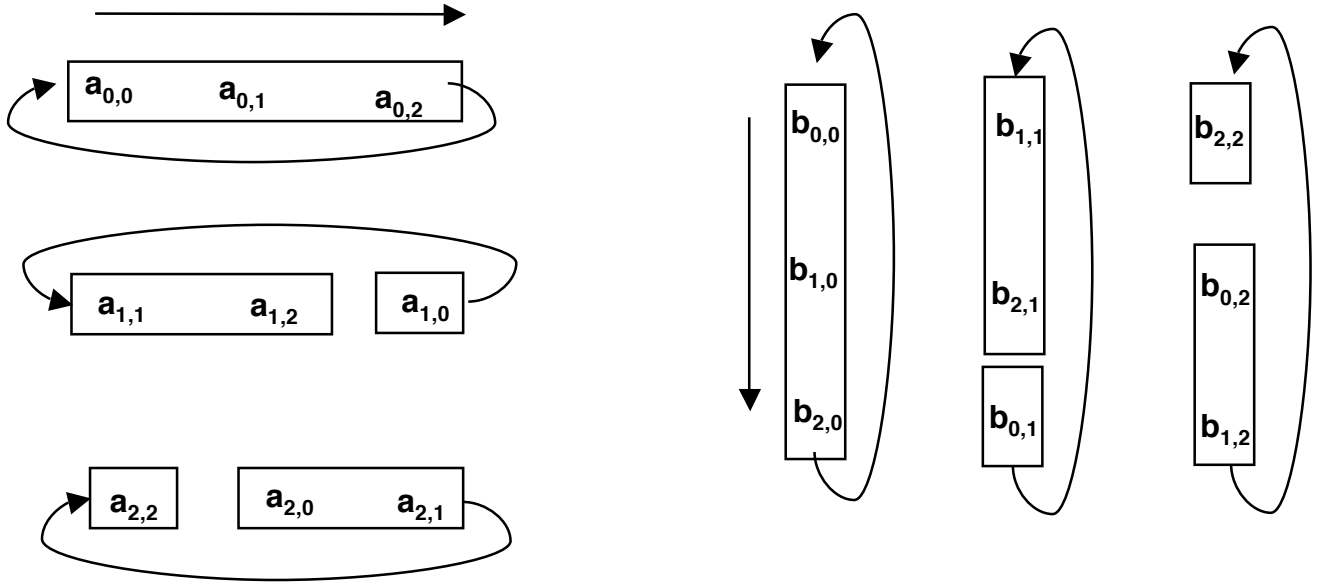
- The strips are first **pre-skewed** so that the correct elements are multiplied at each time step.

Cannon's Method

- Rather than have some processors idle,
- **wrap** the array rows and columns so that **every processor is doing something on each step**.
- In other words, rather than feeding in the elements, they are **rotated** around,
- starting in an initially staggered position as in the systolic model.
- We also change the order of products slightly, to make it correspond to more natural storage by rows and columns.

Cannon's Matrix Multiply ("Pipe-Roll" method)

Note that the **a diagonal is in the left column** and the **b diagonal is in the top row**.



Example sum: $c_{0,2} = a_{0,2}*b_{2,2} + a_{0,1}*b_{1,2} + a_{0,0}*b_{0,2}$

Products computed at each step:

Step 1

a00*b00	a01*b11	a02*b22
a11*b10	a12*b21	a10*b02
a22*b20	a20*b01	a21*b12

Step 2

a02*b20	a00*b01	a01*b12
a10*b00	a11*b11	a12*b22
a21*b10	a22*b21	a20*b02

Step 3

a01*b10	a02*b21	a00*b02
a12*b20	a10*b01	a11*b12
a20*b00	a21*b11	a22*b22

Cannon's Method for **Block** Multiplication

- At each step, entire **blocks** are transmitted down and to the left of neighboring PE's.
- Memory space is conserved.

Fox's Algorithm

- Also for block matrix multiplication, it has a resemblance to Cannon's algorithm.

- The difference is that **on each cycle**:
  - A row block is **broadcast** to **each other** processor in the row.
  - The column blocks are rolled cyclically.

Fox's Algorithm

Step 1

a00*b00	a00*b01	a00*b02
a11*b10	a11*b11	a11*b12
a22*b20	a22*b21	a22*b22

A different **row element** of **a** is **broadcast** in each step, starting with the diagonal

Step 2

a01*b10	a01*b11	a01*b12
a12*b20	a12*b21	a12*b22
a20*b00	a20*b01	a20*b02

b column elements are **rolled** each step, the next elements of a are the first first super-diagonal with wrap-around

Step 3

a02*b20	a02*b21	a02*b22
a10*b00	a10*b01	a10*b02
a21*b10	a21*b11	a21*b12

The sum is accumulated (not shown). For example, in upper-left corner we have a00*b00 + a01*b10 + a02*b20

Fox's Algorithm Pseudo-Code

```
q = sqrt(p) // number of rows, cols in processor grid

  // C = A*B matrix product
  // i,j = process row, column

  // src, dest rows for rotating 'up'
  src  = i+1 mod q;
  dest = i-1 mod q;

  for (stage = 0; stage < q; stage++) {
     k_bar = (i+stage) mod q;
     broadcast(A[i,k_bar]) to row i;
          C[i,j] = C[i,j] + A[i,k_bar]*B[k_bar,j]
     sendrecv(B[k_bar,j],src,dest);
  }
```

Fox's Algorithm in MPI (Pacheco)

```
n_bar = n/grid->q;
Set_to_zero(local_C);

/* Calculate addresses for circular shift of B */
source = (grid->my_row + 1) % grid->q;
dest = (grid->my_row + grid->q - 1) % grid->q;

/* Set aside storage for the broadcast block of A */
temp_A = Local_matrix_allocate(n_bar);

for (stage = 0; stage < grid->q; stage++) {
   bcast_root = (grid->my_row + stage) % grid->q;
   if (bcast_root == grid->my_col) {
      MPI_Bcast(local_A, 1, local_matrix_mpi_t, bcast_root, grid->row_comm);
      Local_matrix_multiply(local_A, local_B, local_C);
   } else {
      MPI_Bcast(temp_A, 1, local_matrix_mpi_t, bcast_root, grid->row_comm);
      Local_matrix_multiply(temp_A, local_B, local_C);
   }
   MPI_Sendrecv_replace(local_B, 1, local_matrix_mpi_t,  dest, 0, source, 0, grid->col_comm, &status);
} /* for */
```

Fox vs. Cannon

- Fox does not require **pre-skewing**.
- Broadcast *may* be cheaper than multiple block transfers.

Other Uses of Systolic Arrays

- Cholesky Factorization of Matrices
- Digital filters (FIR)
- Genome Matching (Smith-Waterman Algorithm, Paracel Inc.)
- http://www.timelogic.com/technology.html (FPGA-based computers)

References

- *Ziad A.A. Alqadi, Musbah Aqel, and Ibrahiem M. M. El Emary,* Performance Analysis and Evaluation of Parallel Matrix Multiplication Algorithms, World Applied Sciences Journal 5 (2): 211-214, 2008

- Analysis of a Class of Parallel Matrix Multiplication Algorithms http://www.cs.utexas.edu/users/plapack/papers/ipps98/ipps98.html

- Two-Dimensional **Block Cyclic** Data Distribution as a Key to Load Balancing and Software Reuse (part of ScaLAPACK User's Guide) http://www.netlib.org/scalapack/slug/node110.html