

# Grammars and Recursive Descent Parsing

March 28–29, 2012

CS 60: Principles of Computer Science

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Assignment 8 (Full Unicalc) due Monday, April 2.

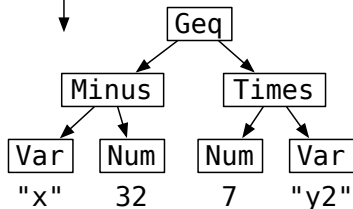
## TRADITIONAL LEXING AND PARSING

( x - 3 2 ) > = 7 \* y 2

LEXING

LPAREN ID DASH INT RPAREN GEQ NUM STAR ID  
 "x" 32 7 "y2"

PARSING



# SPECIFYING SYNTAX VIA CFGs

A *context-free grammar* is a set of rules for producing a set of strings (a *language*).

$$\begin{aligned} S &\rightarrow V + S \mid V \\ V &\rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9 \end{aligned}$$

Ingredients:

- ✓ Nonterminals:  $S, V$
- ✓ Terminals:  $+, 0, 1, 2, \dots, 9$
- ✓ Production rules: (see above)
- ✓ Where to start:  $S$

# PARSE TREES

The *parse tree* of a string makes explicit how a string was produced:

- ✓ Root is the start symbol
- ✓ When we apply a rule, items on the right-hand-side become children

Parse trees for 4 and 4+5?

$$S \rightarrow V + S \mid V$$

$$V \rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$$

## EXERCISE

$$P \rightarrow S * P \mid S / P \mid V$$
$$S \rightarrow V + S \mid V - S \mid V$$
$$V \rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$$

1. Draw the parse tree for  
 $7 * 3 + 9 / 2$
2. How could we change the grammar to get a parse tree with the *usual* mathematical meaning of this expression?
3. What could we add to permit parentheses ( and ) to group sub-expressions?

NAME:

“QUIZ”

Here is the grammar that you will use for the Homework 8.

$$S \rightarrow \text{def } V L \mid L$$

$$L \rightarrow \# E \mid E$$

$$E \rightarrow P + E \mid P - E \mid P$$

$$P \rightarrow K * P \mid K / P \mid K$$

$$K \rightarrow - K \mid Q$$

$$Q \rightarrow R \mid R Q$$

$$R \rightarrow V \mid V \wedge J$$

$$V \rightarrow D \mid D \sim D \mid W \mid ( L )$$

$$J \rightarrow I \mid - I$$

$I \rightarrow$  (any nonnegative integer)

$D \rightarrow$  (any nonnegative number)

$W \rightarrow$  (any word that is a unit name)

Draw the *parse tree* for the input

def accel 2 meter / second^2

# AMBIGUOUS GRAMMARS

A grammar is *ambiguous* if there are strings that have more than one possible parse tree.

$$\begin{aligned} E &\rightarrow E + E \mid E - E \\ &\mid E * E \mid E / E \\ &\mid (E) \mid 0 \mid \dots \mid 9 \end{aligned}$$

$$\begin{aligned} E &\rightarrow E + T \mid E - T \mid T \\ T &\rightarrow T * F \mid T / F \mid F \\ F &\rightarrow (E) \mid 0 \mid \dots \mid 9 \end{aligned}$$

# ABSTRACT SYNTAX TREES

- ✓ Ambiguous grammar have very nice parse trees
- ✓ Unambiguous grammars work better

Compromise:

- ✓ Parse input with an unambiguous grammar
- ✓ Produce “simplified” parse trees

*Abstract Syntax Trees (ASTs) preserve the essential structure of parse trees, without extra crud.*

# RECURSIVE DESCENT PARSING

The simplest algorithm for parsing (which works on some, but not all grammars) is called “recursive descent.”

- ✓ We write one function for each variable/nonterminal in the grammar.
- ✓ Each function  $A()$  will “consume” tokens that match the right-hand-side of some rule  $A \rightarrow \dots$

## TRIVIAL EXAMPLE OF RECURSIVE DESCENT

```
S -> read V
    | write V
V -> x | y | z
```

```
S():
# consume tokens matching
# S in the grammar
peek at the next token
if (it's "read"):
    consume it
    V()
else if (it's "write"):
    consume it
    V()
else:
    report a parsing error
```

```
V():
# consume tokens matching
# V in the grammar
peek at the next token
if (it's "x"):
    consume it
else if (it's "y"):
    consume it
else if (it's "z"):
    consume it
else:
    report a parsing error
```

## SIMPLE EXAMPLE OF RECURSIVE DESCENT

$$L \rightarrow V$$

$$| V, L$$

$$V \rightarrow x | y | z$$

```
L():
  # consume tokens matching
  # L in the grammar
  V()

  peek at the next token
  if (it's ","):
    # there's more to this L
    consume the comma
    L()
  else
    # there was only one V.
    return
```

```
V():
  # consume tokens matching
  # V in the grammar
  peek at the next token
  if (it's "x"):
    consume it
  else if (it's "y"):
    consume it
  else if (it's "z"):
    consume it
  else:
    report a parsing error
```

## EXERCISE

$$J \rightarrow I \mid - I$$
$$V \rightarrow D \mid D \sim D \mid W \mid ( L )$$
$$I \rightarrow (\text{any nonnegative integer})$$
$$D \rightarrow (\text{any nonnegative number})$$
$$W \rightarrow (\text{any word that is a unit name})$$

Sketch pseudocode for  $J()$  and  $V()$ .

(You may assume that  $I()$ ,  $D()$ , and  $W()$  already exist.)

## MINILOG.JAVA

```
// Minilog
// Author: Chris Stone
// Simple logical calculator, for inputs like
//   true or false
//   ( true and false) or false
//   ...etc...
```