

# Dynamic Programming Concluded

April 9–10, 2012

CS 60: Principles of Computer Science

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Assignment 10 (Dynamic Programming) due Monday, April 16.

## RECALL: FIBONACCI NUMBERS

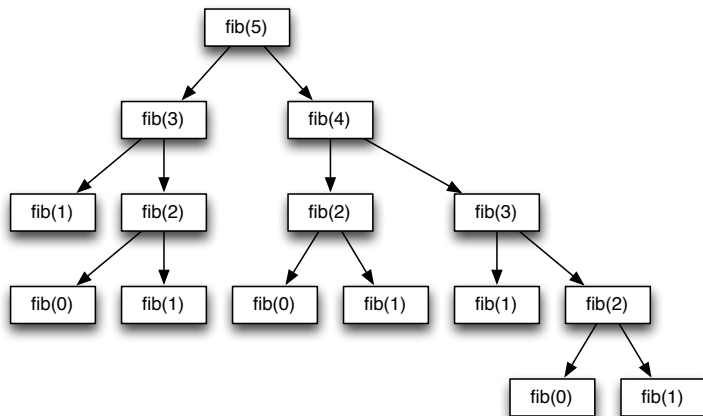
The Fibonacci numbers are easy to compute!

$$\text{fib}(n) = \begin{cases} n & \text{if } n = 0 \text{ or } n = 1 \\ \text{fib}(n - 2) + \text{fib}(n - 1) & \text{if } n \geq 2 \end{cases}$$

The corresponding recursive code is exponentially slow for large  $n$ .

## THE PROBLEM

$$\text{fib}(n) = \begin{cases} n & \text{if } n = 0 \text{ or } n = 1 \\ \text{fib}(n-2) + \text{fib}(n-1) & \text{if } n \geq 2 \end{cases}$$

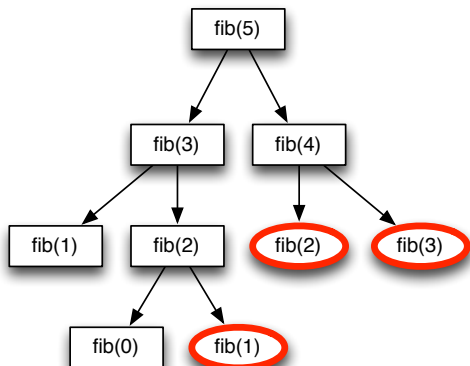


# IDEA 1: "MEMOIZING"

Remember all the inputs and output so far

Return precomputed answers for repeated questions.

$$\text{fib}(n) = \begin{cases} n & \text{if } n = 0 \text{ or } n = 1 \\ \text{fib}(n-2) + \text{fib}(n-1) & \text{if } n \geq 2 \end{cases}$$

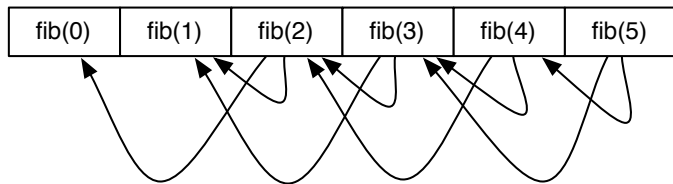


## IDEA 2: “DYNAMIC PROGRAMMING”

Compute each value exactly once, in a “clever” order.

Ensure that problems are solved after their subproblems.

$$\text{fib}(n) = \begin{cases} n & \text{if } n = 0 \text{ or } n = 1 \\ \text{fib}(n - 2) + \text{fib}(n - 1) & \text{if } n \geq 2 \end{cases}$$



## RECALL: THE KNAPSACK PROBLEM

Suppose we have  $n$  kinds of items.

- ✓ They have value (or utility)  $v_1, \dots, v_n$
- ✓ Each has weight (or size or cost)  $w_1, \dots, w_n$ .

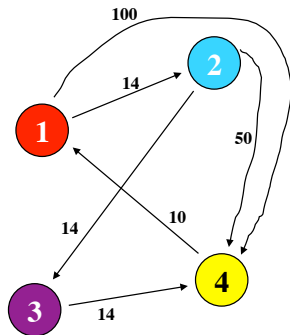
What should we choose, if the maximum total weight is  $W$ ?

$$\text{opt}_v(0) = 0$$

$$\text{opt}_v(W) = \max \begin{cases} \text{opt}_v(W - 1) \\ v_1 + \text{opt}_v(W - w_1) \\ v_2 + \text{opt}_v(W - w_2) \\ \vdots \\ v_n + \text{opt}_v(W - w_n) \end{cases}$$

# FLOYD-WARSHALL

An algorithm for finding *all-pairs* shortest paths!



For  $k = 0, 1, 2, \dots$ , compute:

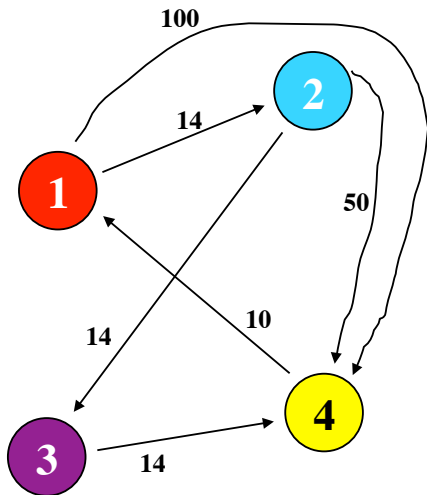
*What is the shortest path from  $s$  to  $d$  via nodes  $1..k$  only?*

# FLOYD-WARSHALL ALGORITHM

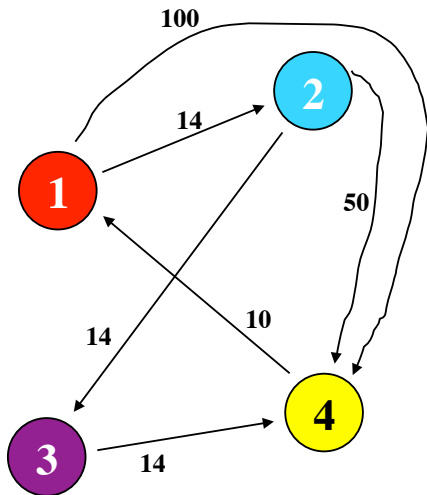
Minimum distance from **src** to  
**dst** using intermediate nodes 1..k

**T[k][src][dst]**

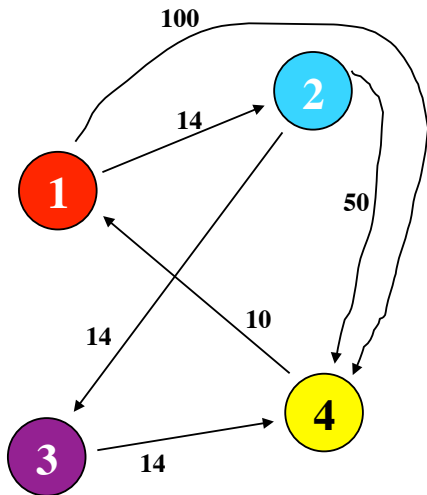
$$= \min \begin{cases} \text{lose k} \\ T[k-1][src][dst] \\ T[k-1][src][k] + \\ T[k-1][k][dst] \\ \text{use k} \end{cases}$$

EXAMPLE:  $K = 0$ 

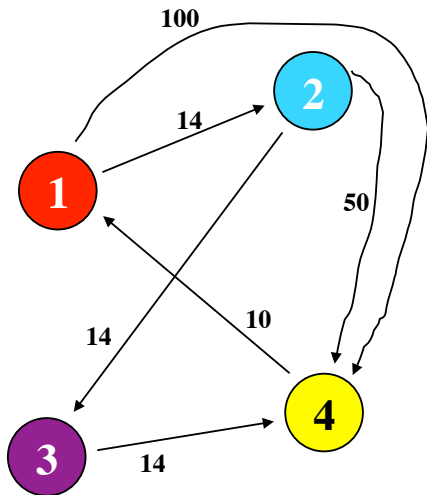
		dst "to"			
		1	2	3	4
src "from"	1	0	14	inf	100
	2	inf	0	14	50
	3	inf	inf	0	14
	4	10	inf	inf	0

EXAMPLE:  $K = 1$ 

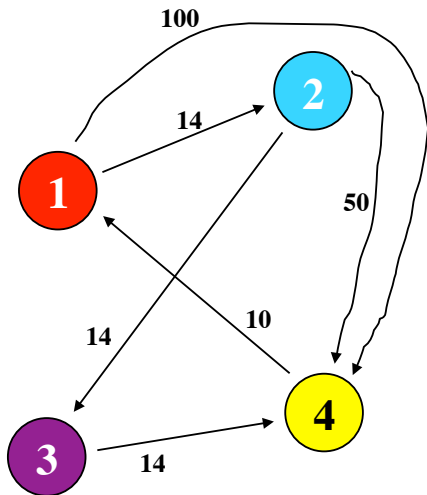
		dst "to"			
		1	2	3	4
src "from"	1	0	14	inf	100
	2	inf	0	14	50
	3	inf	inf	0	14
	4	10	24	inf	0

EXAMPLE:  $K = 2$ 

		dst "to"			
		1	2	3	4
SRC "from"	1	0	14	28	64
	2	inf	0	14	50
	3	inf	inf	0	14
	4	10	24	38	0

EXAMPLE:  $K = 2$ 

		dst "to"			
		1	2	3	4
src "from"	1	0	14	28	42
	2	inf	0	14	28
	3	inf	inf	0	14
	4	10	24	38	0

EXAMPLE:  $K = 4$  (DONE)

		dst "to"			
		1	2	3	4
SRC "from"	1	0	14	28	42
	2	38	0	14	28
	3	24	38	0	14
	4	10	24	38	0