CS 181AG
Lecture 12

## Packet Classification Cont.

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## Feedback

## - Common themes:

- Office hours > lecture + activities > assignments > slides
- Assignments reinforce material without taking $10+$ hours a week
- Having lecture slides beforehand (with room to annotate) is important
- Slides are not a great study material on their own
- Lecture recordings are available
- Strongly encourage you to come to office hours
- Starter code could be explained more
- Will try!


## Midterm

- Take-home midterm, released Oct 25 8am, due Oct 30 10pm (on Gradescope)
- Must be done in one sitting (no limit on time), though it will be written to take about 3 hrs
- Open: lecture slides, videos, textbook, assignments, NOT open: anything else on internet, conversations with classmates
- Will cover everything through 10/12
- For fairness, I will not respond to any questions about midterm content. If anything seems ambiguous, please clearly write down your assumptions and I will grade accordingly (within reason :) )


## Upcoming Schedule

- 10/10, 10/12: Finish packet classification
- 10/17-10/18: Fall break!
- No assignment due week of fall break
-10/19: Start new topic (switching)
- 10/24: Midterm Review in class
- No assignment due week of midterm
-10/25 - 10/30: Midterm
- No assignment due right after midterm
- Assignment 7 due Nov. 7 (Monday)


## 2D Schemes: Grid of Tries

| Rule | Destination | Source |
| :---: | :---: | :---: |
| R1 | D1 = 0* | S1 = 10* |
| R2 | D2 $=0$ * | S2 $=01^{*}$ |
| R3 | D3 $=00 *$ | S3 = 11* |
| R4 | D4 $=00$ * | S4 $=1$ * |
| R5 | D5 = 0* | S5 = 1* |
| R6 | D6 = 10* | S6 = 1* |
| R7 | D7 = * | S7 = 00* |
| R8 | D8 $=*$ | S8 = * |

## 2D Schemes: Grid of Tries

- Grid of tries using backtracking (each source appears exactly once in entire grid of tries). Find longest destination D, check corresponding source trie, then backtrack to all prefixes of $D$, keeping track of lowest-cost rule so far



## 2D Schemes: Grid of Tries

- Grid of tries without backtracking: once longest matching destination is found, all possible sources are contained in corresponding source trie



## Trie of Tries

- Construct a trie of destination prefixes
- Each valid destination prefix (D) points to a trie of source prefixes
- The source trie contains source prefixes for all rules with a destination field exactly equal to $D$
- Instead of backtracking, let's add switch pointers, pointing us to the next location we would backtrack to if it exists


## Switch Pointers



## One More Example

- Draw all switch pointers for correctness, but remember: it is very important that when looking up the source trie, we keep track of the best rule we've seen so far
- i.e., it is possible for switch pointers to point us to a worse rule, but we need to remember the best we've seen so far

| Rule | Destination | Source |
| :---: | :---: | :---: |
| R1 | D1 = 101* | S1 = 0* |
| R2 | D2 $=10 *$ | S2 = $1^{*}$ |
| R3 | D3 $=0$ * | S3 = 10* |
| R4 | D4 $=0$ * | S4 = 00* |
| R5 | D5 = * | S5 = 00* |
| R6 | D6 $=$ * | S6 = 01* |
| R7 | D7 = * | S7 = * |

## Geometrical View of Classification

- Each combination of fields represents a single point in space
- Can assess complexity of classifier by number of disjoint regions
- Observations:
- Number of disjoint regions is small compared to worst case ( N rules, K fields $>$ worst case of $\mathrm{N}^{\mathrm{k}}$. Reality is closer to $\mathrm{N}^{*} \mathrm{k}$
- Once source and destination are matched, less than 20 remaining matching rules for all other fields combined


## Divide and Conquer Algorithms

- Solve best match by field, then efficiently put results together


## Using Bitmaps

- Key observation: operations on bitmaps can be sped up by hardware

| Rule | Destination | Source | Dest. Port | Src Port | Flags |
| :--- | :--- | :--- | :--- | :--- | :--- |
| R1 | D1 $=0^{*}$ | S1 $=10^{*}$ | 25 | $*$ | $*$ |
| R2 | D2 $=0^{*}$ | S2 $=01^{*}$ | 25 | 123 | $*$ |
| R3 | D3 $=0^{*}$ | S3 $=1^{*}$ | 53 | $*$ | UDP |
| R4 | D4 $=0^{*}$ | S4 $=1^{*}$ | $*$ | $*$ | $*$ |
| R5 | D5 $=0^{*}$ | S5 $=1^{*}$ | 20 | $*$ | $*$ |
| R6 | D6 $=$ 10 $^{*}$ | S6 $=1^{*}$ | $*$ | $*$ | TCP_ack |
| R7 | D7 $=*$ | S7 $=00^{*}$ | $*$ | $*$ | UDP |
| R8 | D8 $=*$ | S8 $=*$ | $*$ | $*$ | $*$ |

## Using Bitmaps

- Find least cost rule using bitmaps for the following (for simplicity, IPs will be 4 bits):
- (0001, 0011, 53, 35, UDP)
- (0111, 1011, 25, 30, -)
- (0011, 1000, 20, 30, UDP)

| Rule | Destination | Source | Dest. Port | Src Port | Flags |
| :--- | :--- | :--- | :--- | :--- | :--- |
| R1 | D1 $=0^{*}$ | S1 $=10^{*}$ | 25 | $*$ | $*$ |
| R2 | D2 $=0^{*}$ | S2 $=01^{*}$ | 25 | 123 | $*$ |
| R3 | D3 $=00^{*}$ | S3 $=11^{*}$ | 53 | $*$ | UDP |
| R4 | D4 $=00^{*}$ | S4 $=1^{*}$ | $*$ | $*$ | $*$ |
| R5 | D5 $=0^{*}$ | S5 $=1^{*}$ | 20 | $*$ | $*$ |
| R6 | D6 $=10^{*}$ | S6 $=1^{*}$ | $*$ | $*$ | TCP_ack |
| R7 | D7 $=*$ | S7 $=00^{*}$ | $*$ | $*$ | UDP |
| R8 | D8 $=*$ | S8 $=*$ | $*$ | $*$ | $*$ |

## Using Bitmaps

- Does require linear search
- Summary bits can help, but still requires linear search of summary bits

| Rule | Destination | Source | Dest. Port | Src Port | Flags |
| :--- | :--- | :--- | :--- | :--- | :--- |
| R1 | D1 $=0^{*}$ | S1 $=10^{*}$ | 25 | $*$ | $*$ |
| R2 | D2 $=0^{*}$ | S2 $=01^{*}$ | 25 | 123 | $*$ |
| R3 | D3 $=0^{*}$ | S3 $=1^{*}$ | 53 | $*$ | UDP |
| R4 | D4 $=0^{*}$ | S4 $=1^{*}$ | $*$ | $*$ | $*$ |
| R5 | D5 $=0^{*}$ | S5 $=1^{*}$ | 20 | $*$ | $*$ |
| R6 | D6 $=$ 10 $^{*}$ | S6 $=1^{*}$ | $*$ | $*$ | TCP_ack |
| R7 | D7 $=*$ | S7 $=00^{*}$ | $*$ | $*$ | UDP |
| R8 | D8 $=*$ | S8 $=*$ | $*$ | $*$ | $*$ |

## Cross Producting

- Divide and Conquer algorithm (solve best match by field, then put results together)
- Precompute result for each possible best match
- Problem?

