

CS 181AG
Lecture 20

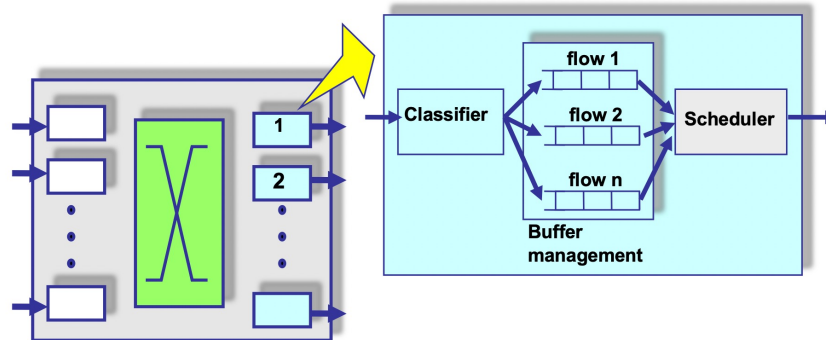
Output Scheduling (cont.)

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Nov 21, 2022

Recap

- Within each output queue, there are several flows (set of packets that take the same path through whole network)
- RED discards packets *before* they are placed in the output queue if the running average of queue length $>$ threshold. Why?
- Sometimes we need a way of controlling long-term rate and maximum burst size for a flow -> token bucket algorithm

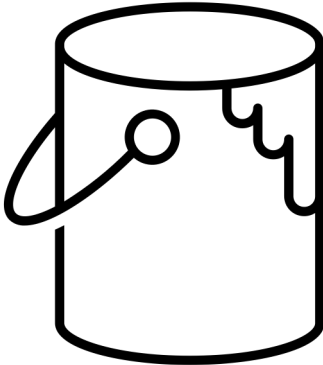


Token Bucket Shaping

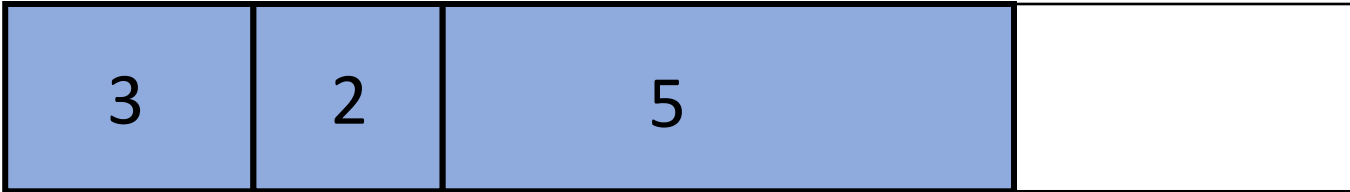
- Goal: for a particular flow, limit 1) the long-term rate of sending and 2) the maximum burst size
- Why: we might need to limit the amount of news traffic, or UDP traffic

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb



Token Bucket



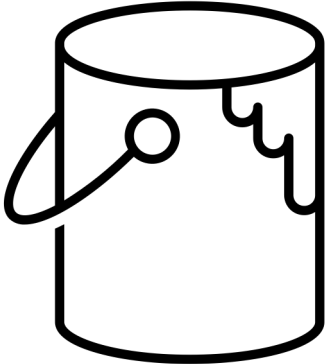
Buffer

Token Bucket Example

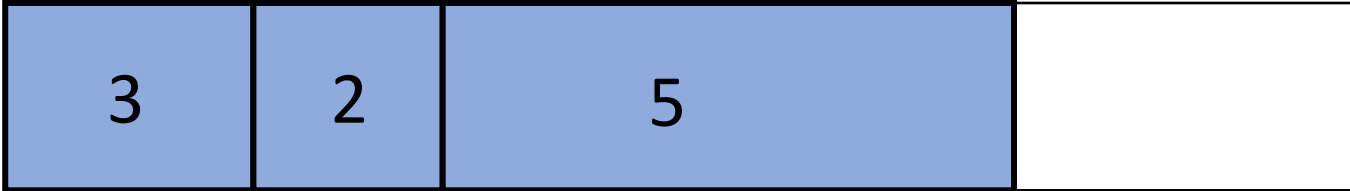
$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb

r tokens enter bucket
every time slot

Token can hold
max 5 tokens



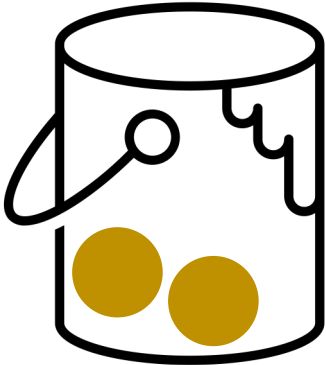
Token Bucket



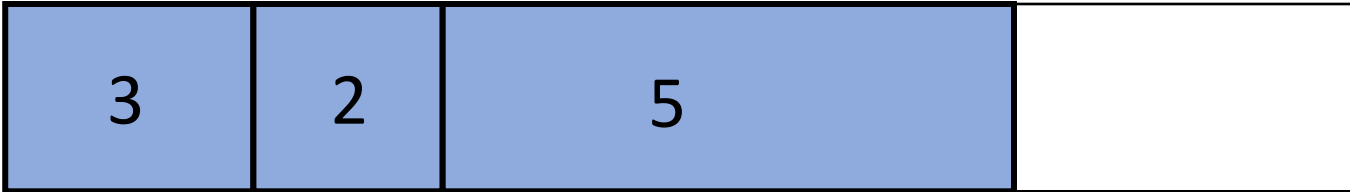
Buffer

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
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Token Bucket

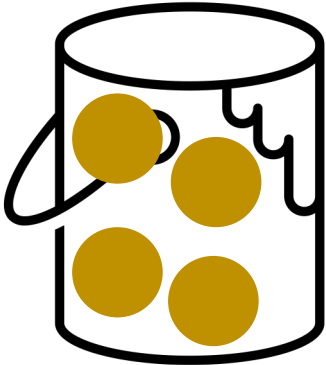


Buffer

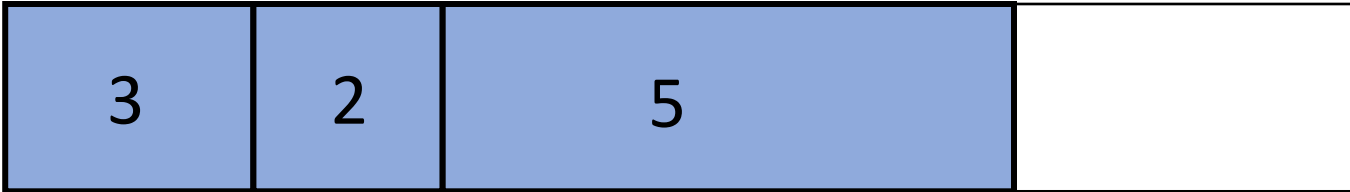
$t = 0$

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb



Token Bucket

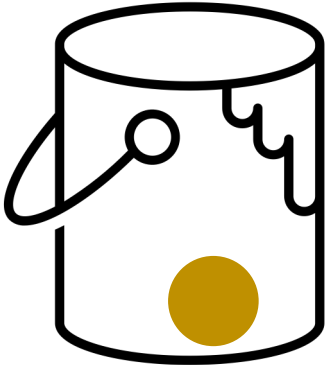


Buffer

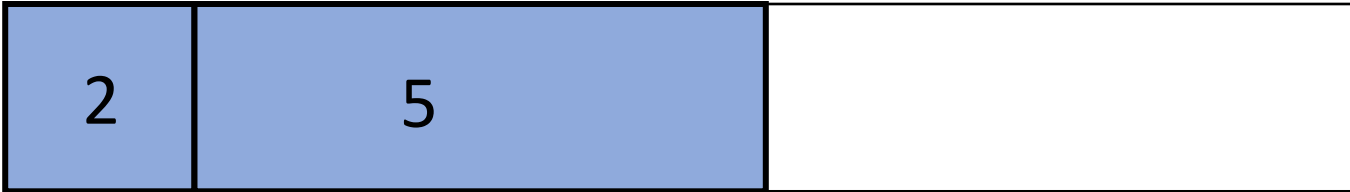
$t = 1$

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb



Token Bucket

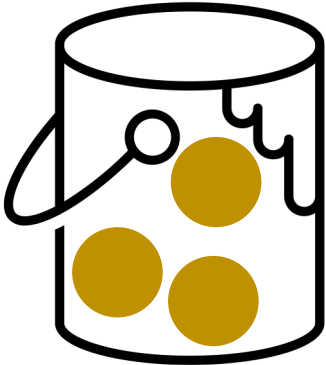


Buffer

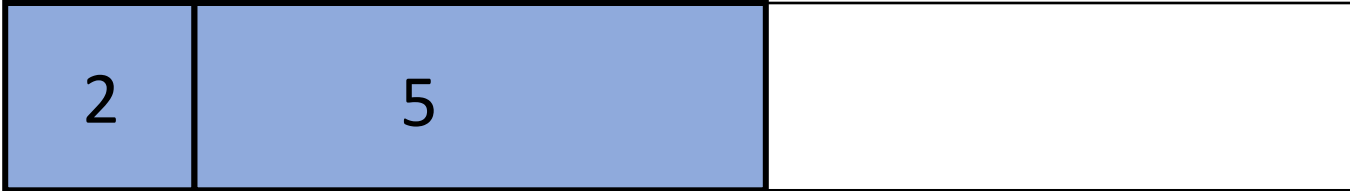
$t = 1$

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb



Token Bucket

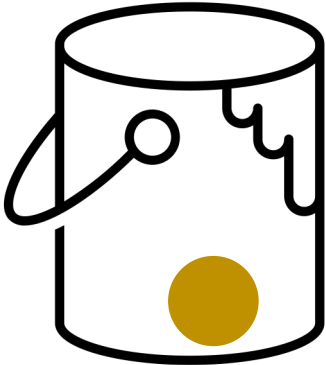


Buffer

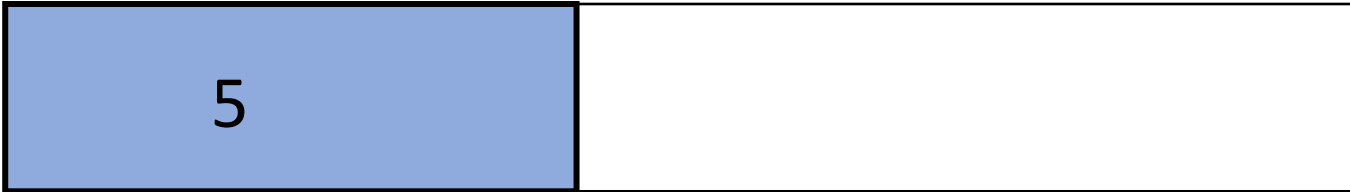
$t = 2$

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb



Token Bucket

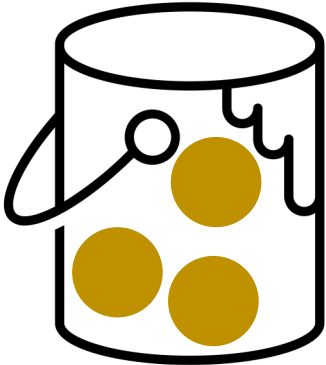


Buffer

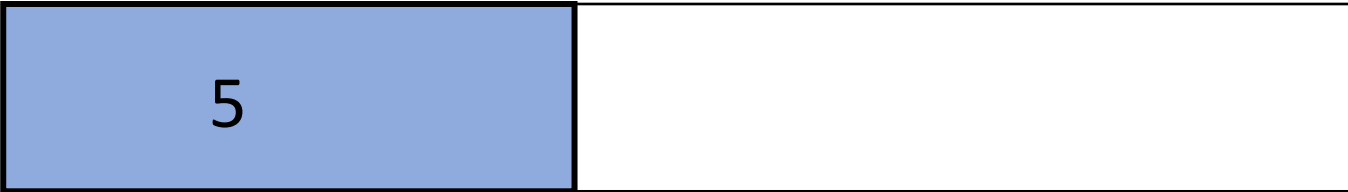
$t = 2$

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb



Token Bucket

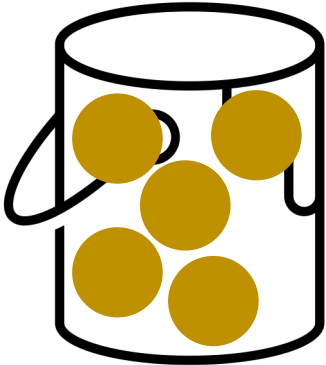


Buffer

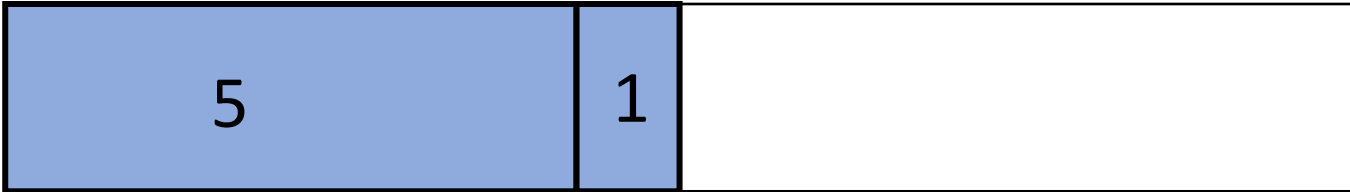
$t = 3$

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb



Token Bucket

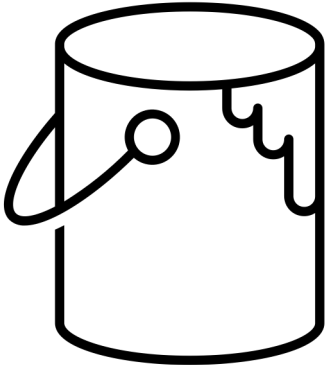


Buffer

$t = 4$

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb



Token Bucket

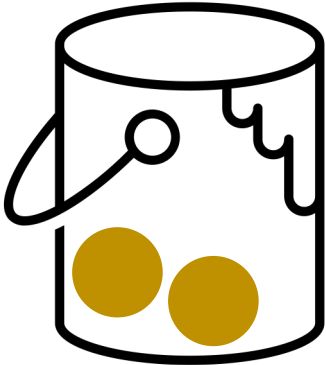


Buffer

$t = 4$

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb



Token Bucket

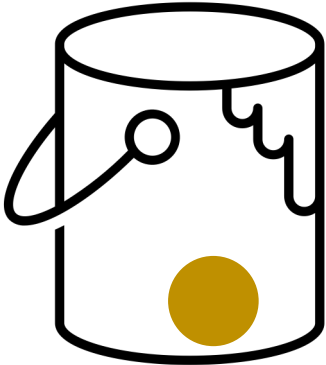


Buffer

$t = 5$

Token Bucket Example

$r = 2 \text{ Mbps}$
 $b = 5 \text{ Mb}$
1 token = 1 Mb



Token Bucket

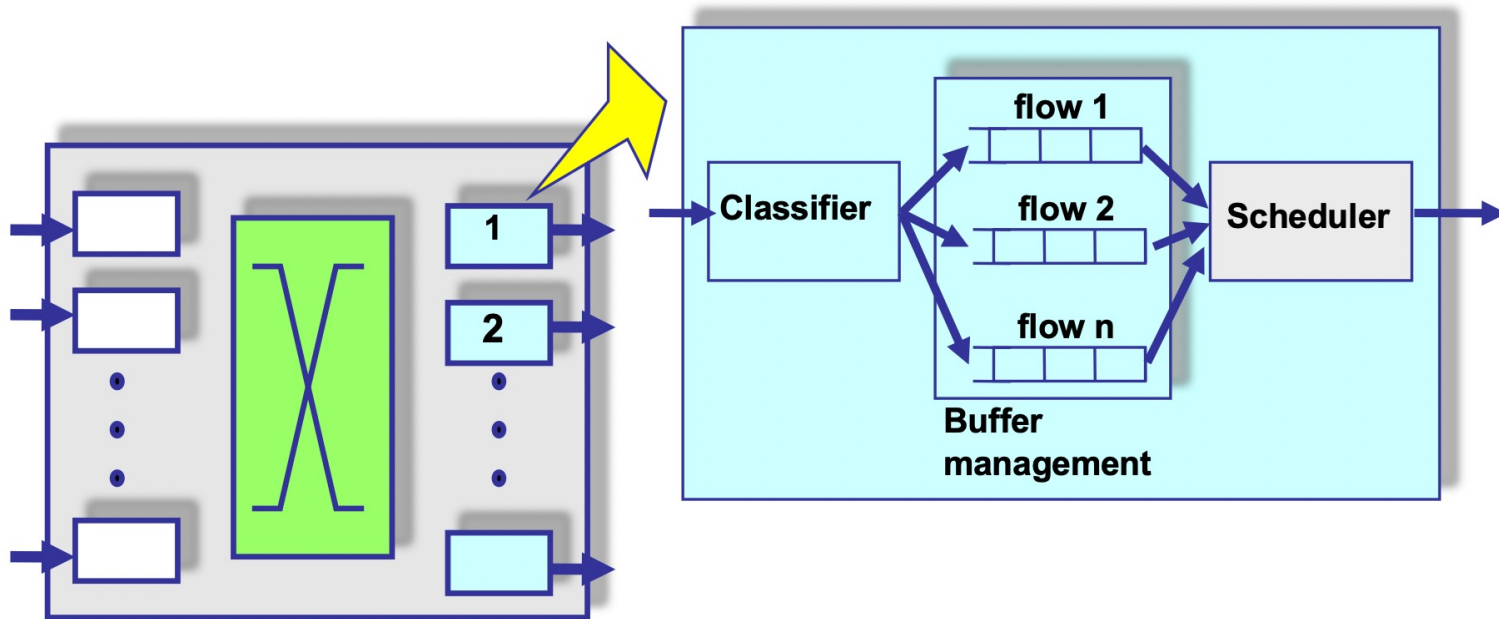


Buffer

$t = 5$

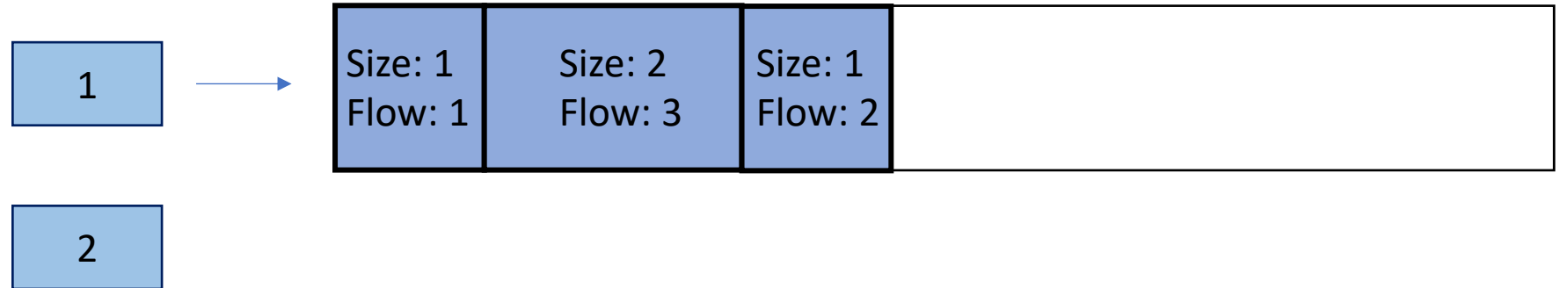
Token Bucket Per Flow

- So far, we have considered token bucket algorithm *per flow*



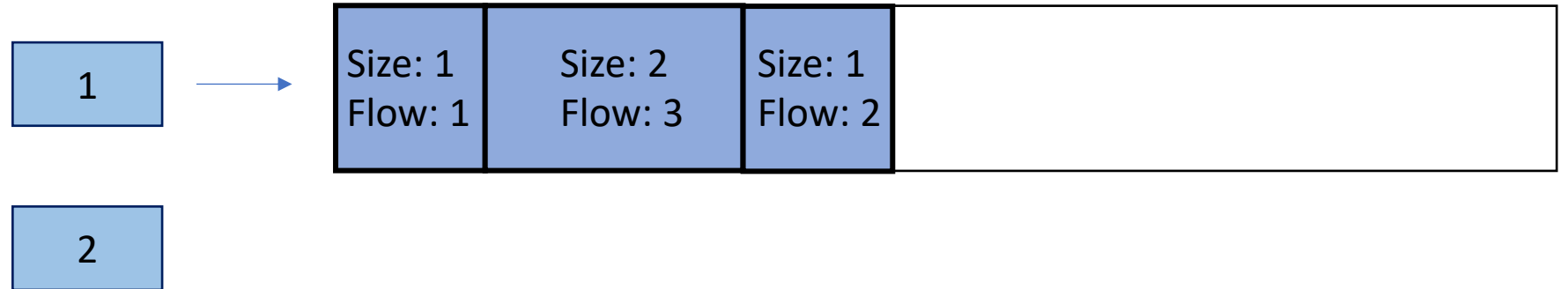
Token Bucket: One Buffer?

- What would be the issue if we kept all flows in the same buffer?



Token Bucket: One Buffer?

- What would be the issue if we kept all flows in the same buffer?



- Waiting for one packet to get enough tokens could block other flows that are able to send
- Solution: drop any packet that has insufficient tokens to send

Token Bucket Shaping vs Policing

- Shaping: one buffer per flow in the output queue; if there are insufficient tokens to send packet at head of buffer, wait
- Policing: one buffer for all packets in the output queue; if there are insufficient tokens to send packet at head of buffer, drop (still keep 1 token bucket per flow)

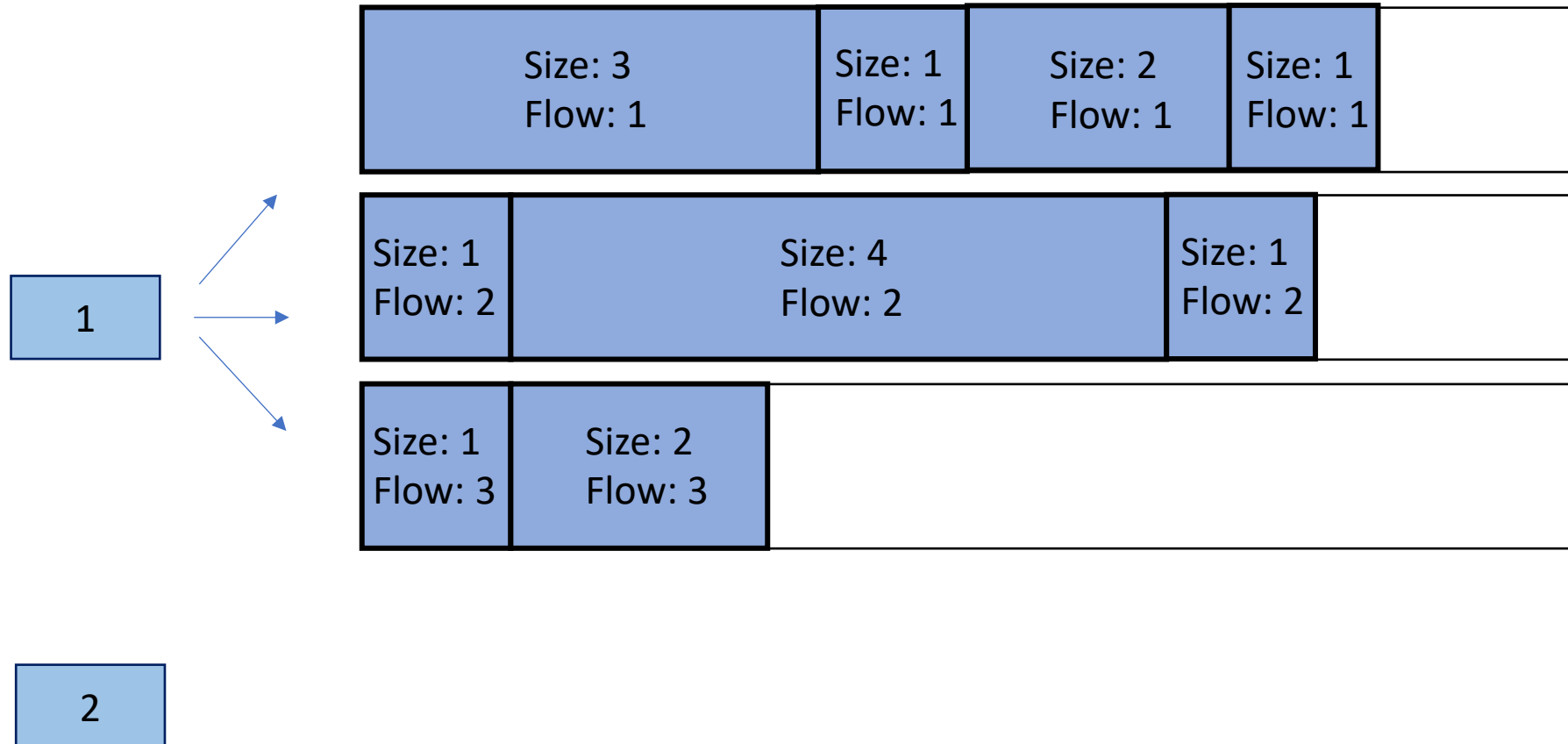
Token Bucket Policing Example

- It takes one unit of time to look at the head of the buffer and decide what to do with it
- Two tokens are added at each time slot and may be used in the same time slot they are added. Token bucket can hold max 8
- Show the head of the buffer and the token count in each bucket at the start of each time step

Size: 2 Flow: 2	Size: 3 Flow: 1	Size: 4 Flow: 1	Size: 4 Flow: 2	Size: 6 Flow: 2	Size: 3 Flow: 1	
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Handling Multiple Flows

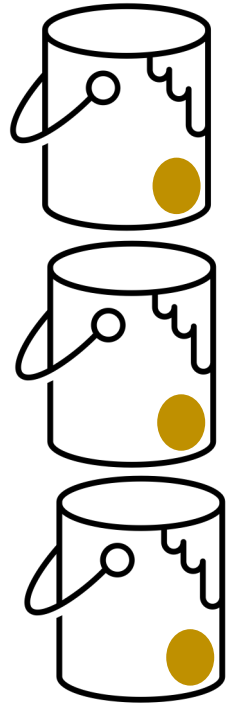
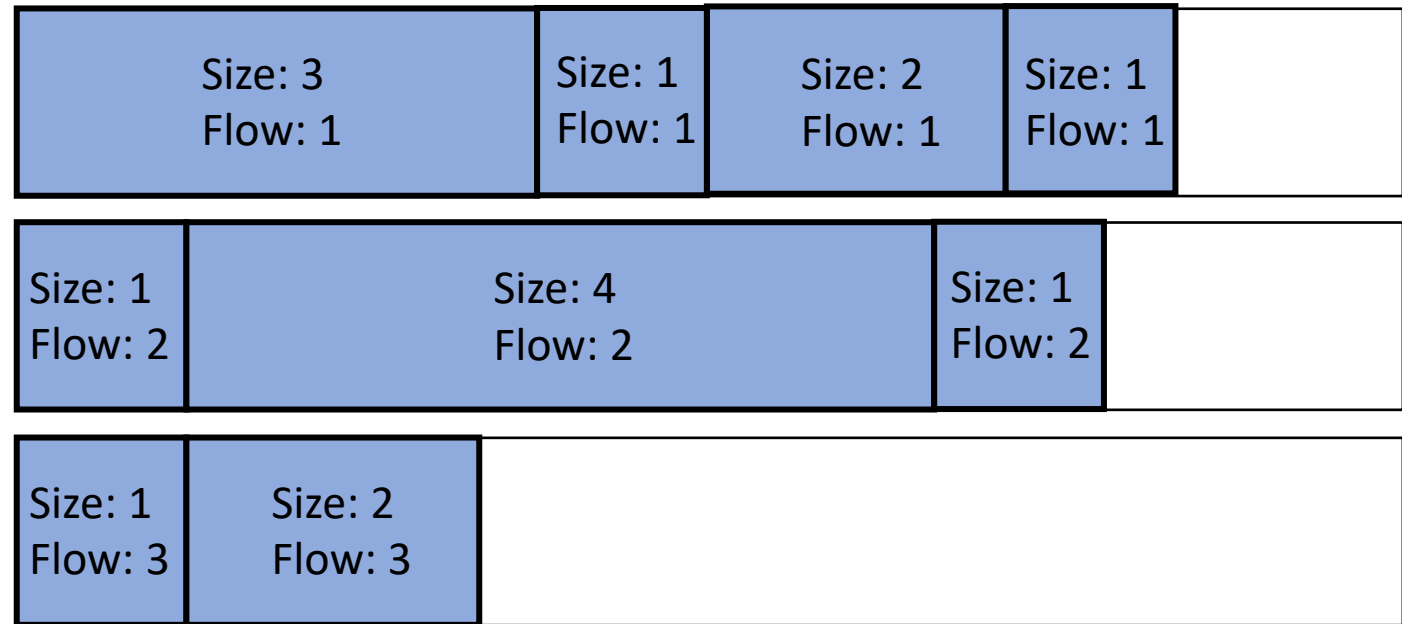
- In the case where we have multiple flows and we wait, how do we decide which flow to send from?



Deficit Round Robin

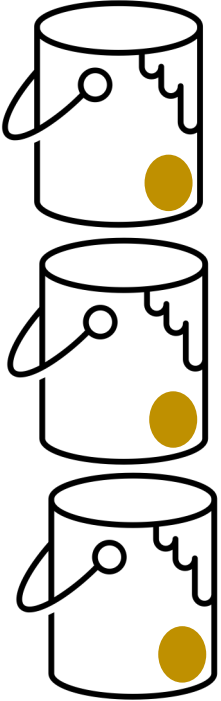
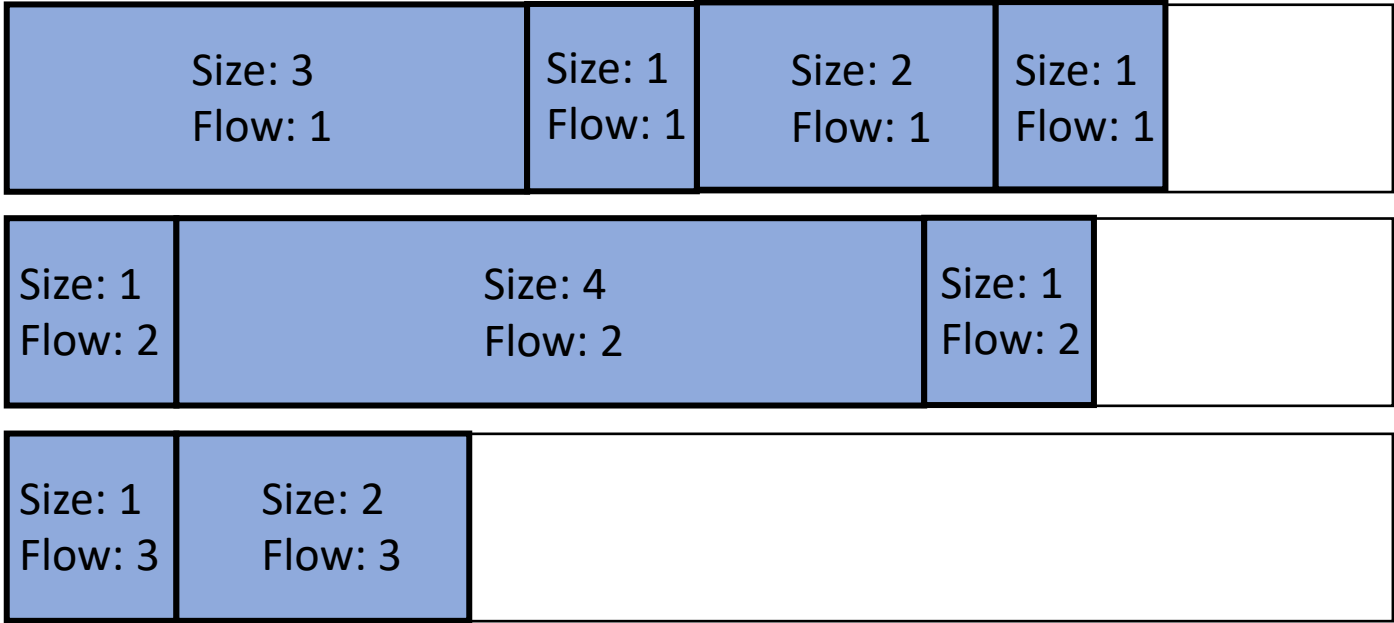
Salary = 1 token

Cycle 1



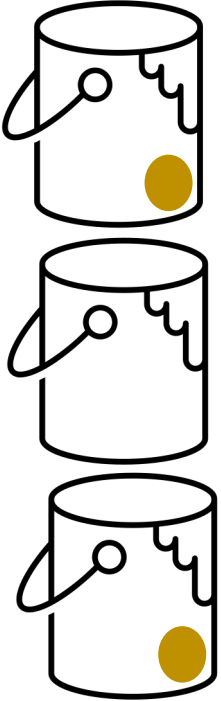
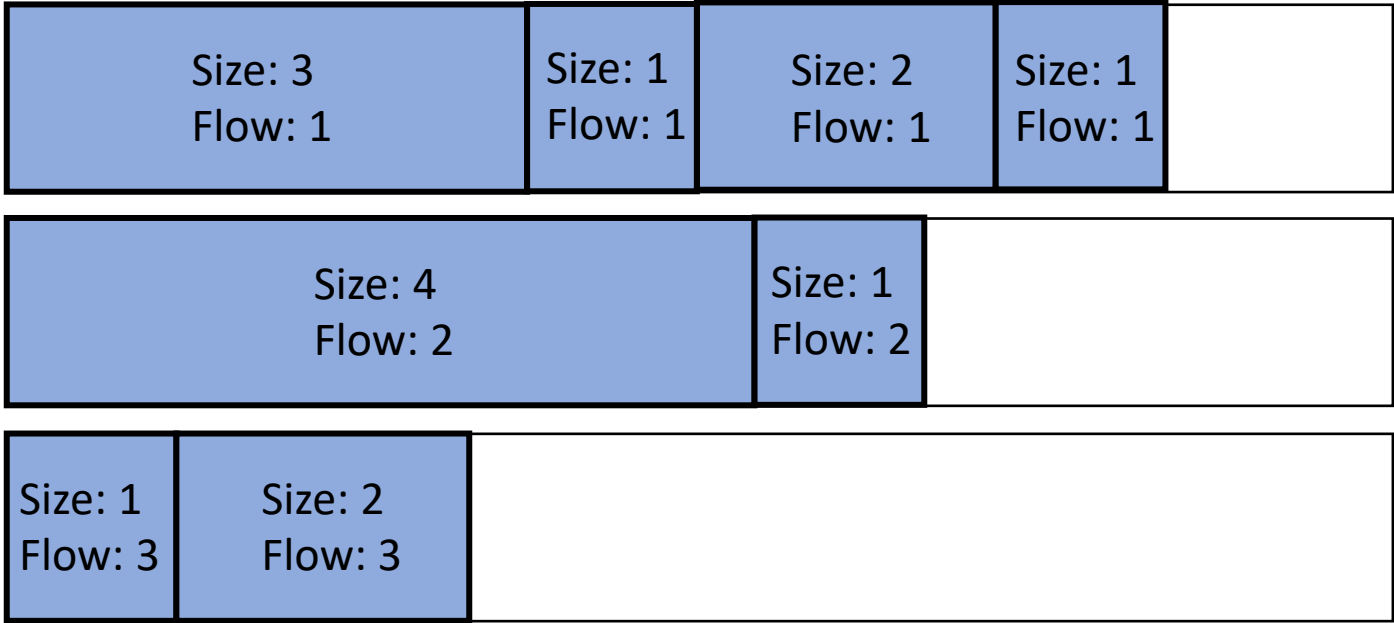
Deficit Round Robin

Cycle 1



Deficit Round Robin

Cycle 1

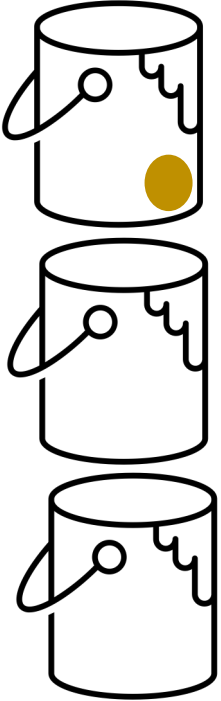
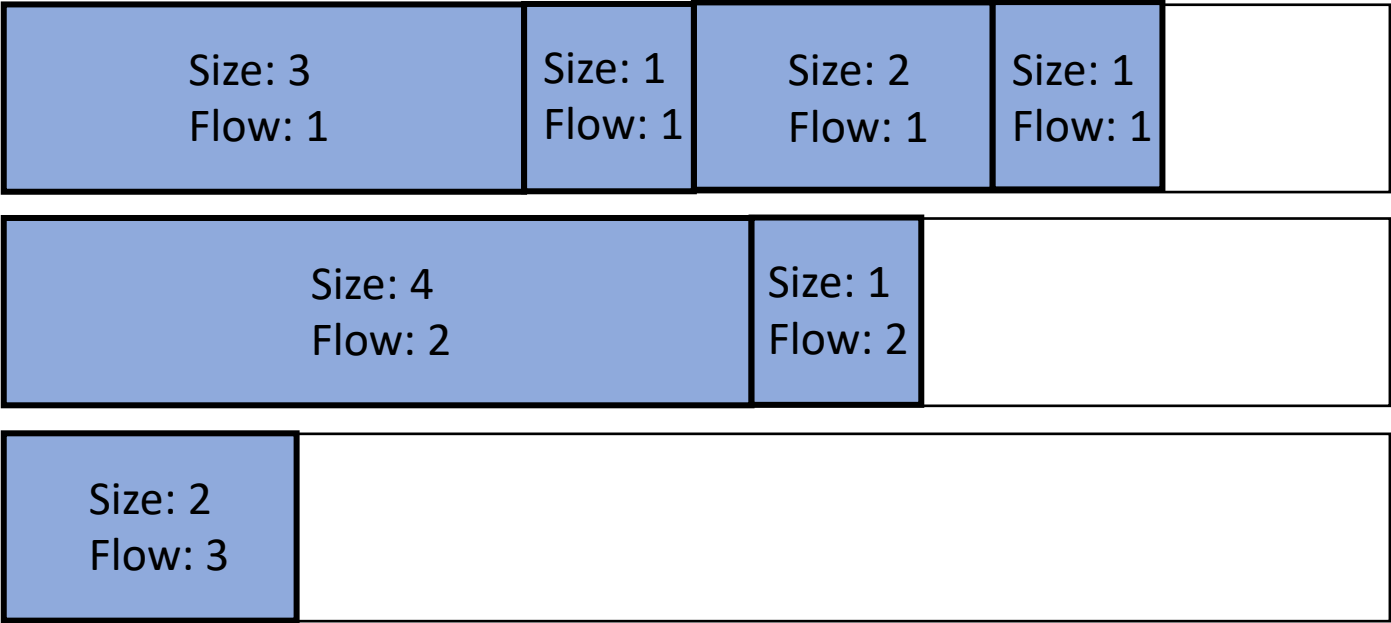


Deficit Round Robin

End of
Cycle 1

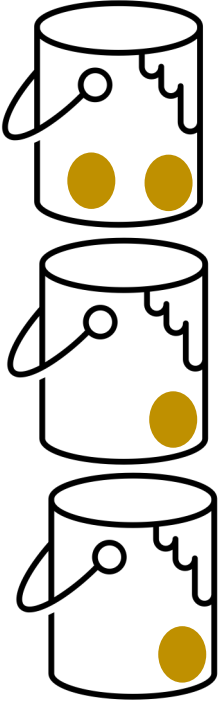
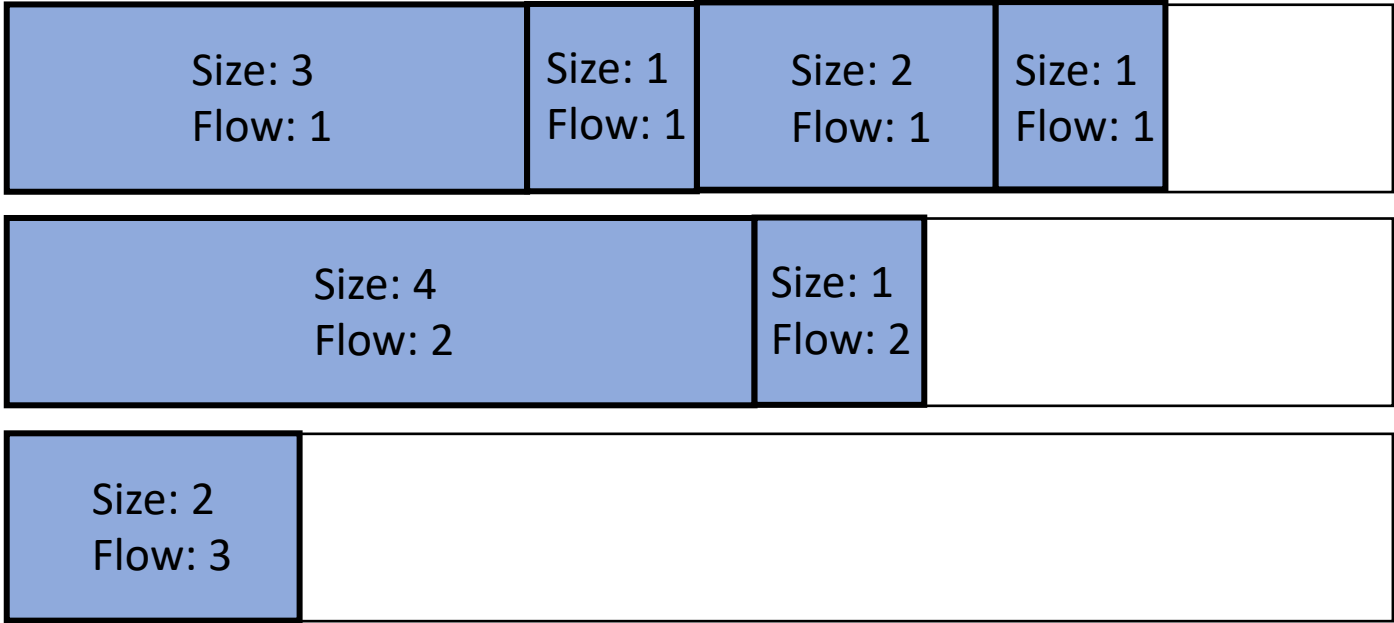
1

2



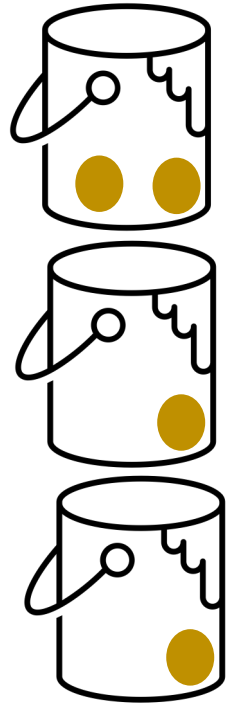
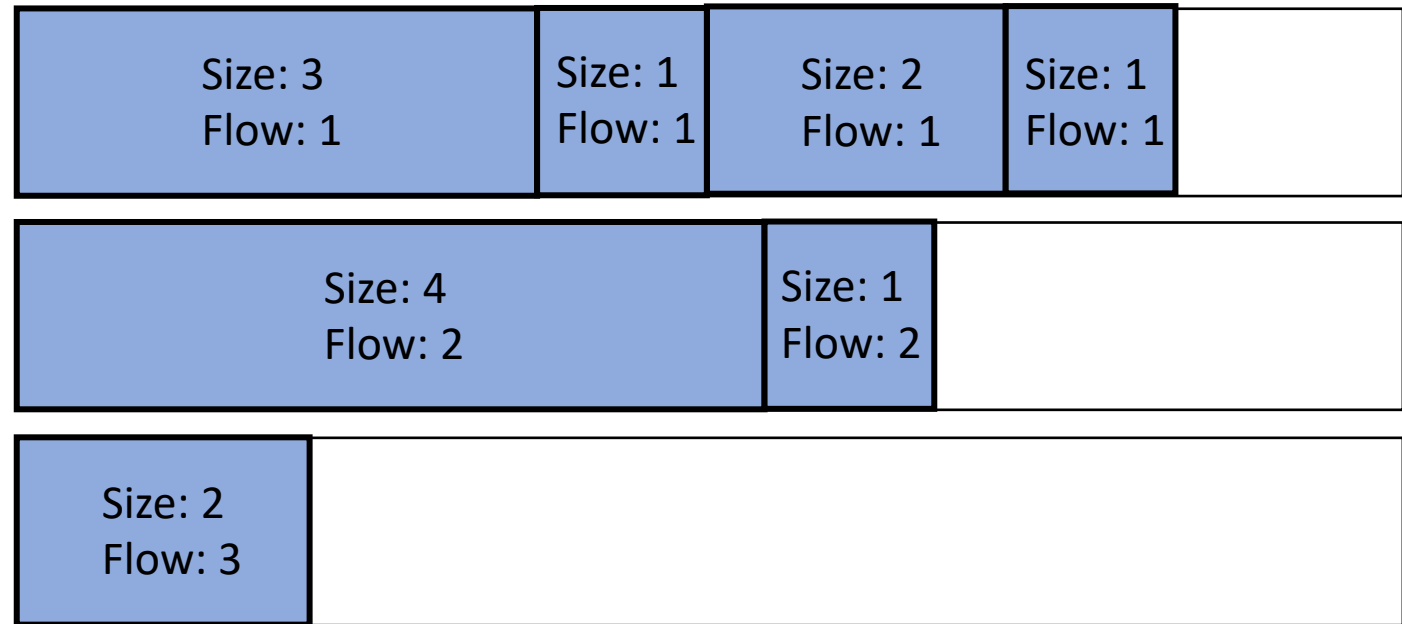
Deficit Round Robin

Cycle 2



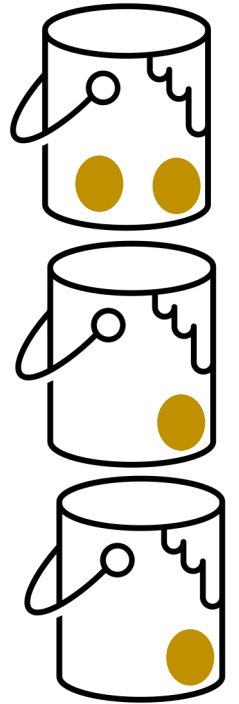
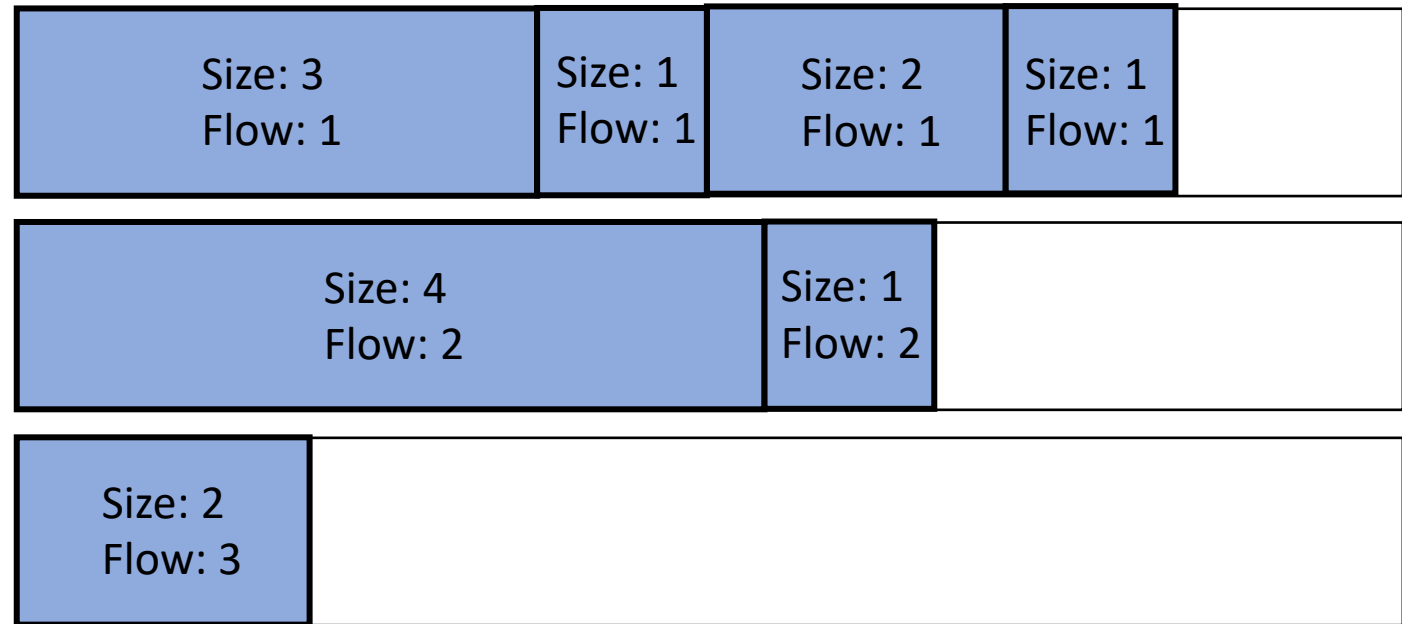
Deficit Round Robin

Cycle 2



Deficit Round Robin

Cycle 2

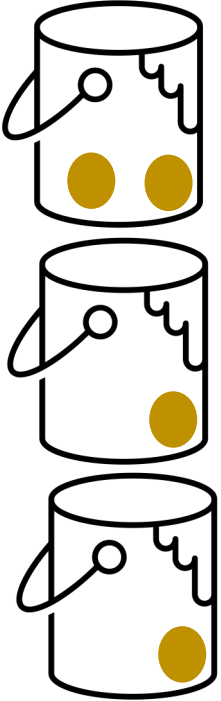
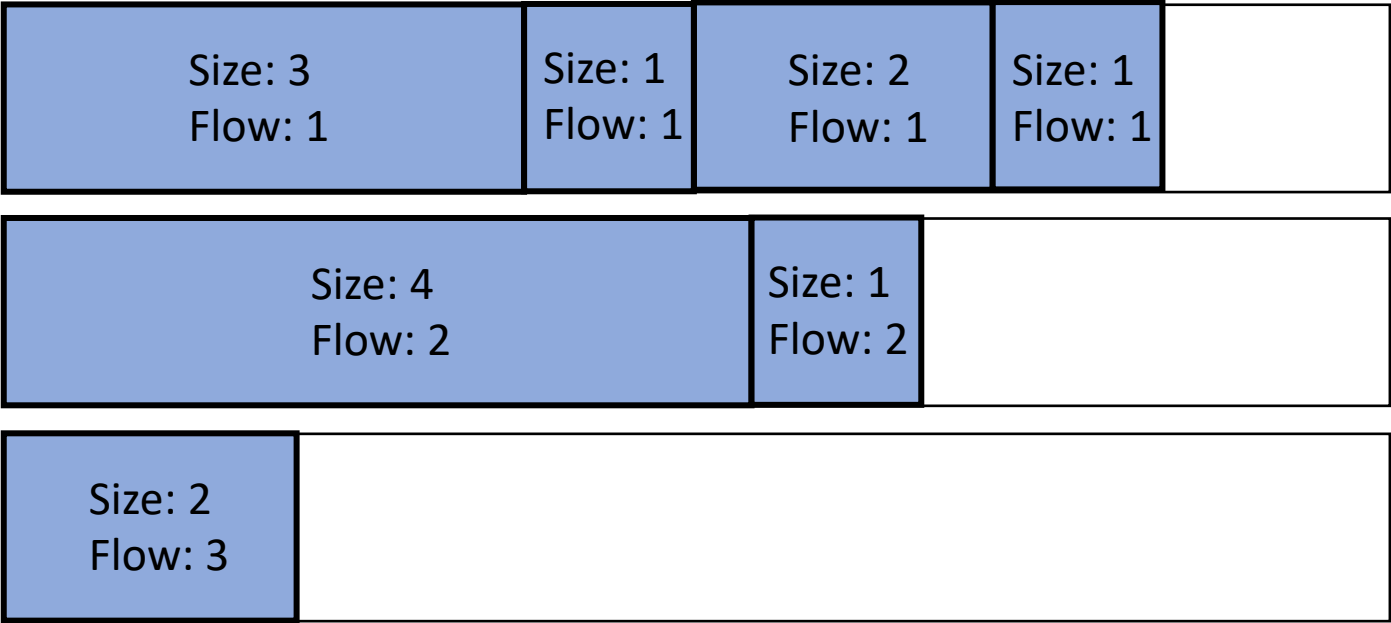


Deficit Round Robin

End of
Cycle 2

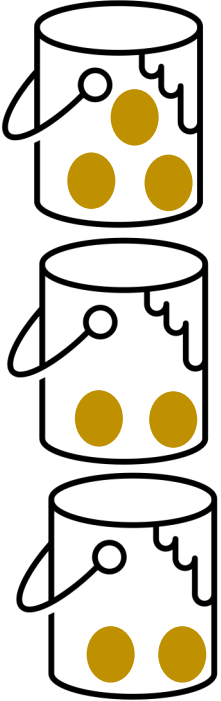
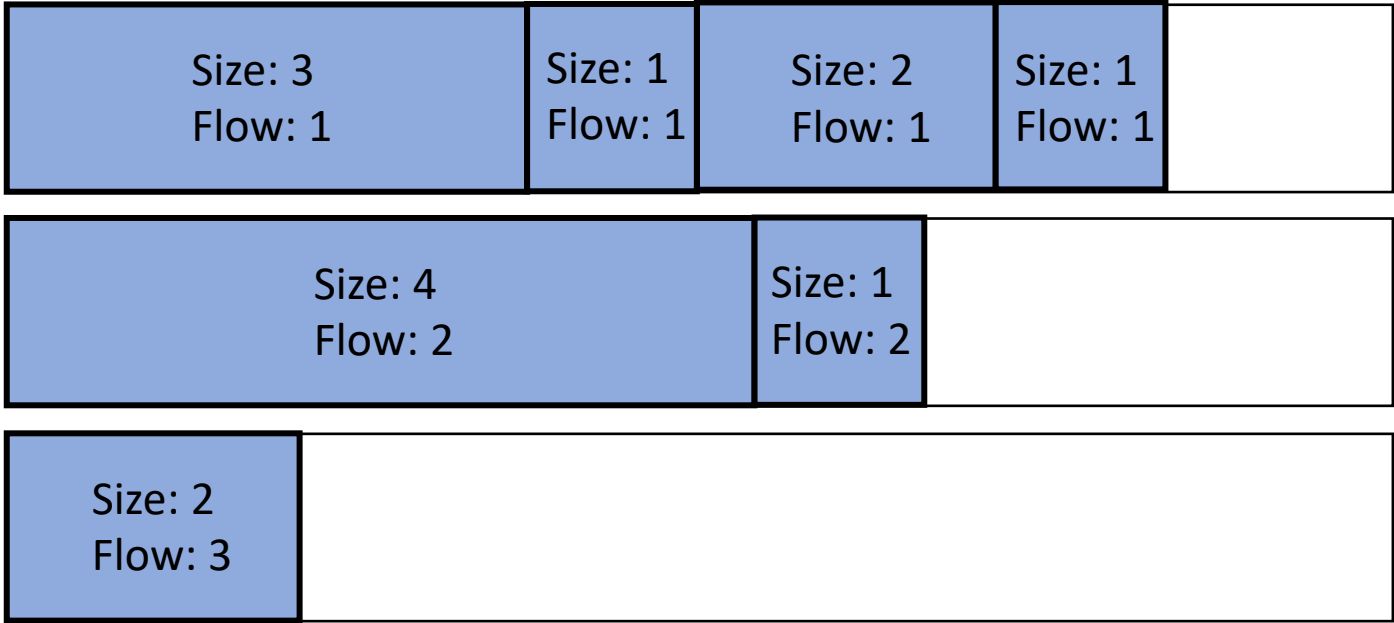
1

2



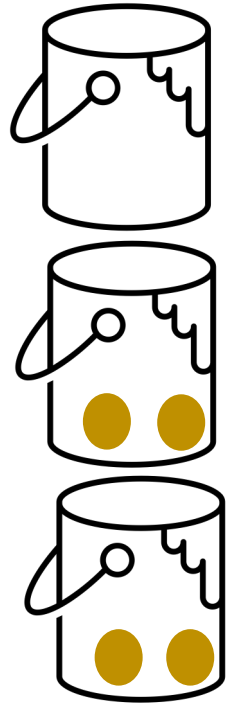
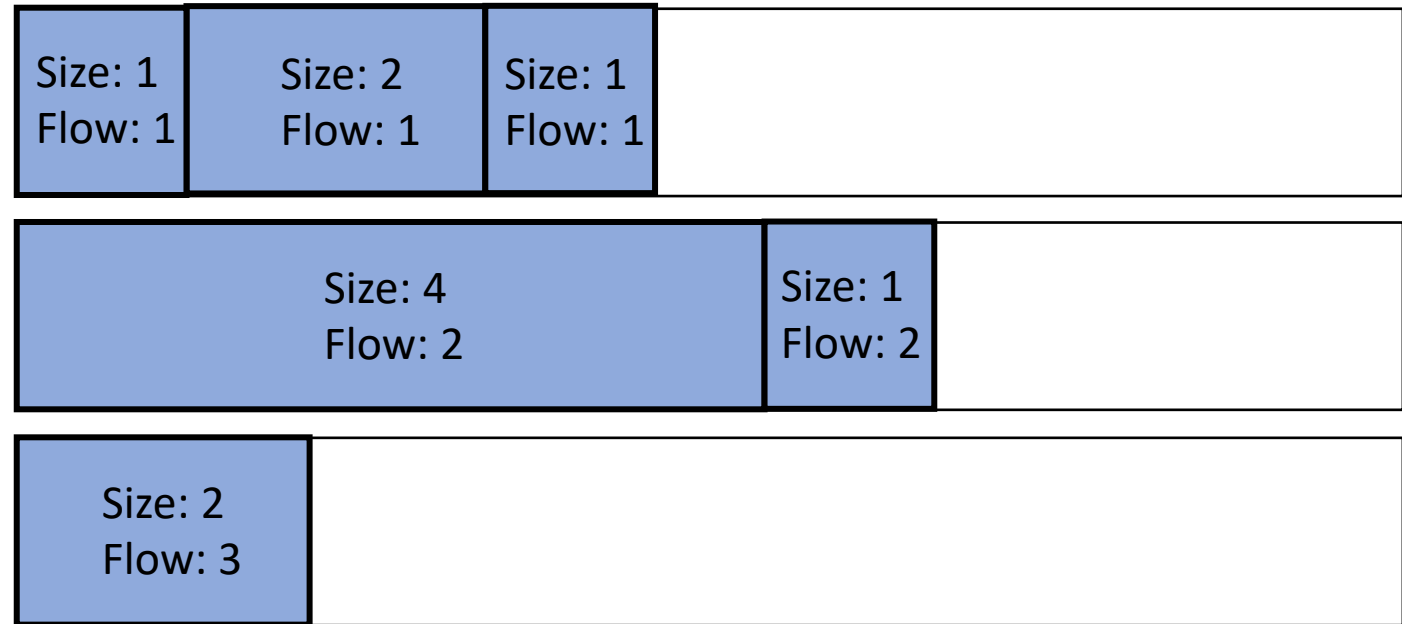
Deficit Round Robin

Cycle 3



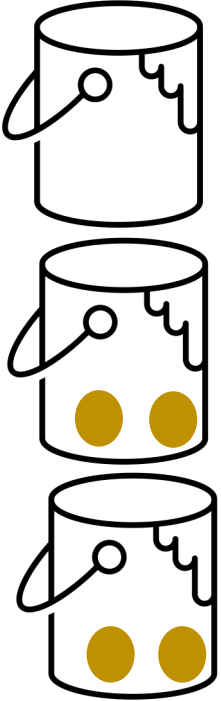
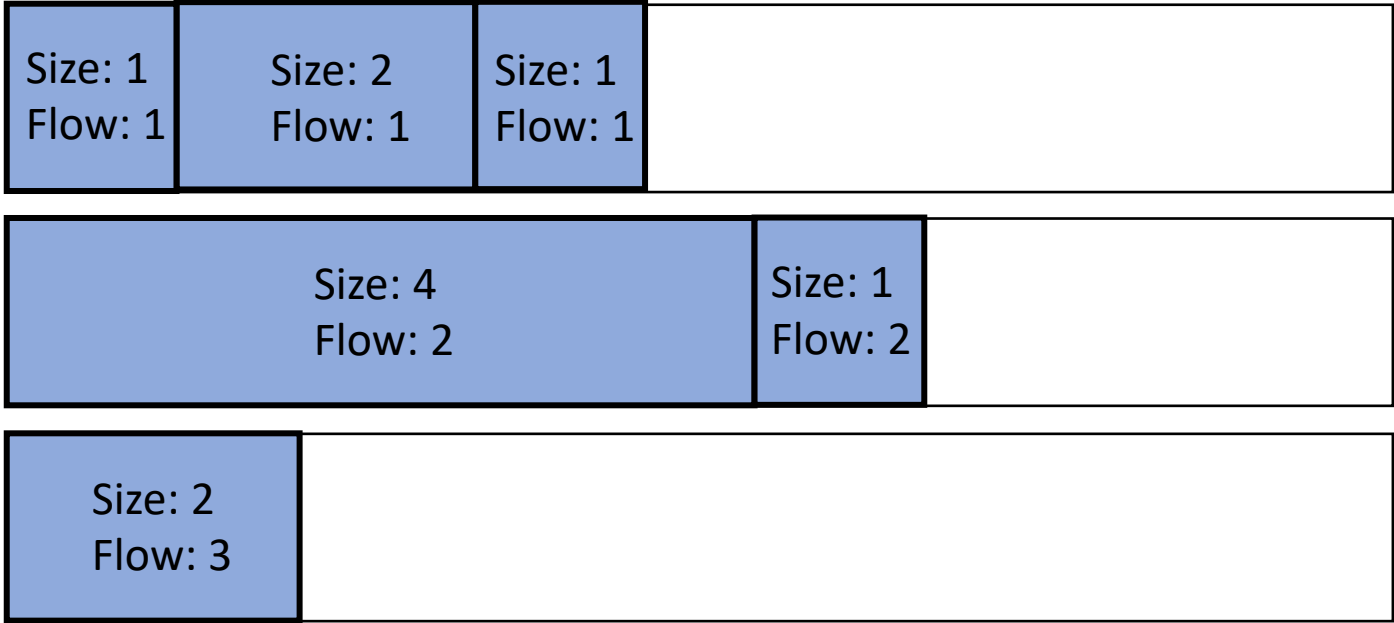
Deficit Round Robin

Cycle 3



Deficit Round Robin

Cycle 3

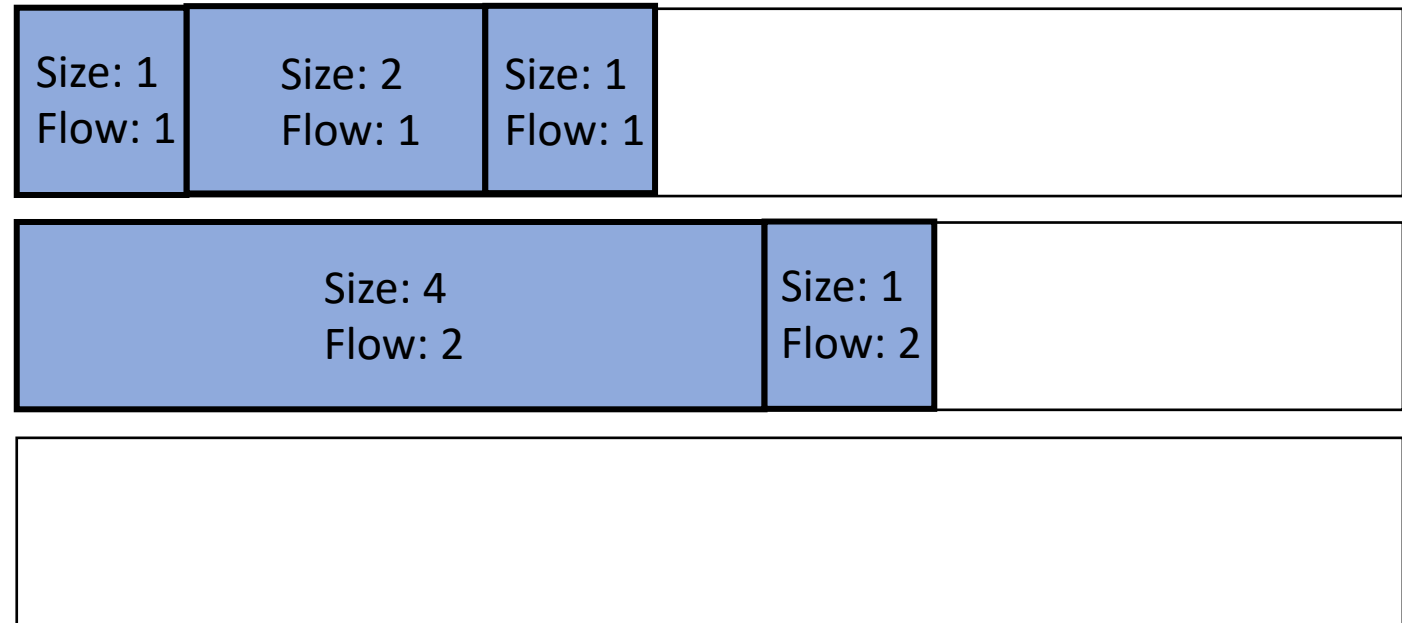


Deficit Round Robin

End of
Cycle 3

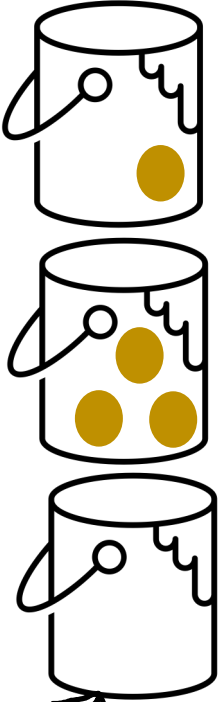
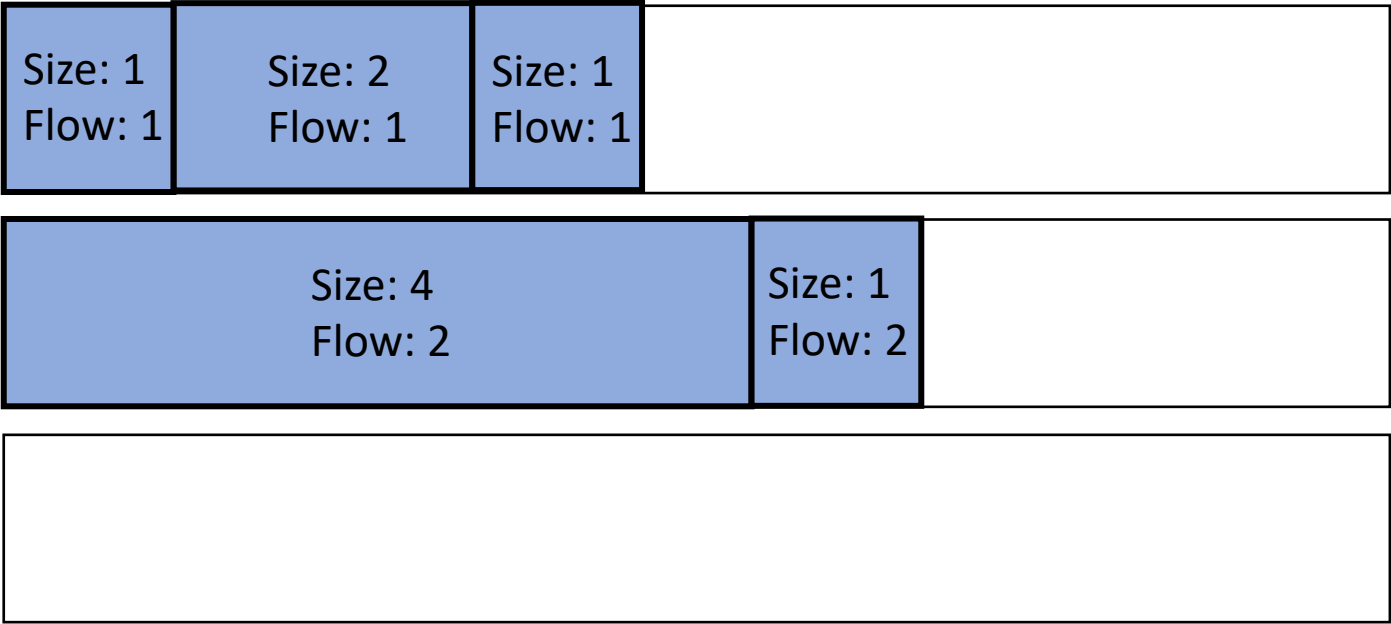
1

2



Deficit Round Robin

Cycle 4



Do not give salary to flow whose buffer is empty



Deficit Round Robin

- Keep buffer and token bucket per flow and round robin, sending as many packets as possible each turn
- Each bucket gets a “salary,” or an amount put in on each cycle of the round robin